# Quantum-Darwinism-encoding Transitions 

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## Classical objectivity from quantum mechanics

Consider a system $S$ and its environment $E$. $E$ has objective knowledge of $S$ only if many small fractions $F \subset E$ are correlated with $S$. [Zurek et al]

[illustration by a 4 y.o.]

## The Quantum Darwinism (QD) plateau

A strong indication of the emergence of objectivity is the "QD plateau" of the mutual information $I(S, F)$ between the system and a fraction of the environment. When $S$ is a single qubit, we have:


A perfect QD plateau (black) and some approximate ones (dashed green).

## Reminder: mutual information

The mutual information between $A$ and $B$ is defined as

$$
I(A, B):=H(A)+H(B)-H(A B)
$$

where $H(X):=-\operatorname{Tr}\left[\rho_{X} \log _{2} \rho_{X}\right]$ is the von Neumann entropy of the reduced density matrix to $X$ (in unit of bits).

- Without correlation $\left(\rho_{A B}=\rho_{A} \otimes \rho_{B}\right), I(A, B)=0$.
- If $A$ and $B$ are classically correlated bits, $I(A, B)=1$.
- If $A$ and $B$ are maximally entangled qubits, $I(A, B)=2$.



## Basic example

The GHZ state

$$
|\mathrm{GHZ}\rangle=\left(|0\rangle_{S}|00 \ldots 0\rangle_{E}+|1\rangle_{S}|11 \ldots 1\rangle_{E}\right) / \sqrt{2}
$$

has a perfect Darwinism plateau:

$$
I(S, F \subset E)=\left\{\begin{array}{ll}
0 & f=0 \\
1 & 0<f<1 \\
2 & f=1
\end{array} \quad(f=|F| /|E|)\right.
$$



## Opposite of objectivity: encoding

A random state in the Hilbert space $\mathcal{H}_{S E}$ does not have a plateau. Instead, it exhibits encoding:


A small subsystem $(f<1 / 2)$ knows nothing of $S$ ! (This is a consequence of the Page curve.)

## Question

Are QD and encoding "phases of information", separated by some sharp transition(s)?



Since generic chaotic unitary dynamics leads to encoding, such a transition would have a similar flavor to MBL and measurement-induced transitions.

## In this talk

We consider two models with Quantum Darwinism-encoding transitions (QDETs).

- Both are "mean-field" models, defined on an expanding tree (equivalent to all-to-all interaction).
- Model 1 is a simple random Clifford model [2305.03694].
- Model 2 is deterministic and non-Clifford [to appear].



## Models for QDET: general setup

- Start with an entangled pair $R A$ and $N-1$ environment qubits ("recruits") in a factorized state.
- $A$ and the recruits undergo a unitary evolution $U$ depending on a parameter $p$.
- Pick a subset $F$ of the output qubits; $f=|F| / N$.

Figure of merit: $I(F, R)$ as a function of $f$ and $p$


## Models on a expanding tree

We now further specify the geometry:


Here, a branching duplicates the in-going bit (in the $Z$ basis)

$$
Y=\emptyset \underset{|0\rangle}{\not \oint_{0}}=\sum_{s= \pm 1}|s\rangle|s\rangle\langle s|
$$

and $\circ$ is a one-site unitary depending on the parameter $p$.

## The QD limit $(p=0)$

In both models, $\boldsymbol{\bullet}=I$ when $p=0$. So

generates a GHZ state on $R$ and the $N$ output bits. Then

$$
I(F, R)=1,(0<|F|<N)
$$

The information on $R$ is perfectly broadcast. When $p>0$, the $o$ 's will tend to "scramble " the information, and drive transitions towards encoding.

## Model 1: definition

- Each $u=0$ is a random one-site Clifford ${ }^{1}$ with probability $p$, and identity otherwise.
- The fraction $F$ is chosen randomly: every leaf $\in F$ with probability $f$.

${ }^{1}$ By definition, $O \mapsto u^{\dagger} O u$ permutes randomly the Pauli operators.


## Model 1: exact phase diagram




Two transitions and a mixed phase (stochastic mixture).

## Model 2: non-Clifford and deterministic

- The scramblers are deterministic and identical

$$
u=\boldsymbol{o}=\exp (i \theta Y)=\left(\begin{array}{cc}
\cos \theta & -\sin \theta \\
\sin \theta & \cos \theta
\end{array}\right), \theta=\frac{p \pi}{4}
$$

$$
u_{p=0}=I, u_{p=1} \text { swaps } Z \text { and } X
$$

- $F$ is deterministic, with $f=1 / 2^{k}(k>1)$.

(Example of $f=1 / 4$ )


## Model 2: phase diagram



- An encoding phase at $p>p_{c}=\frac{1}{2}$ (exact prediction).
- A mixed phase at $p_{d}<p<p_{c}$ and a QD phase at $p<p_{d}$ for some $p_{d} \approx 0.33$ (inset).


## Coarse-grained measure

What can we learn about $R$ by measuring only the total spin $M=\sum_{j \in F} Z_{j}$ ? Or the total spin in the two half trees $\left(M_{L}, M_{R}\right)$, etc? Almost everything!


Useful for observing the transitions in practice.

## Summary

Unitary dynamics can exhibit both Quantum Darwinism and encoding phases, separated by transitions. The phases are characterized by how information dissipates. Questions for future work include:

- Universal critical properties in finite dimensions (in an expanding spacetime?).
- Relevance of statistical physics for foundational questions, e.g., Wigner's friend paradox.



## Quantities of interest

- Mutual information $I(F, R)$.
- Conditional entropy: does $F$ know about $R$ ? Or, can measuring $F$ (in a basis of your choice) disentangle $R$ ?

$$
J(R ; F)=H(R)-\min _{\text {basis }}\langle H(R)\rangle_{\text {measure }}
$$

In general $J(R ; F) \leq I(F, R) . I(F, R)-J(R ; F)$ is called the "quantum discord'.


## Defining the order parameter (Model 2)

Consider measuring $Z$ 's on $F$. To each outcome $\vec{m}=\left(m_{j}\right)_{j \in F}$ is associated an operator on $A$ :

$$
Q_{\vec{m}}=V^{\dagger}\left[\prod_{j \in F} \frac{1+m_{j} Z_{j}}{2}\right] V=p_{\vec{m}}\left(I+u_{\vec{m}} Z+v_{\vec{m}} X\right)
$$

where $V|\psi\rangle=U|\psi\rangle\left|0^{N-1}\right\rangle . p_{\vec{m}}$ is the probability of the outcome. $\left(u_{\vec{m}}, v_{\vec{m}}\right)$ is the information on $R$ revealed by the measure. The order parameter is a distribution of $(u, v)$ (in the unit disk)

$$
P(u, v)=\sum_{\vec{m}} p_{\vec{m}} \delta\left(u_{\vec{m}}-u\right) \delta\left(v_{\vec{m}}-v\right)
$$

## Order parameter: properties

$$
P(u, v)=\sum_{\vec{m}} p_{\vec{m}} \delta\left(u_{\vec{m}}-u\right) \delta\left(v_{\vec{m}}-v\right)
$$

"Distribution of knowledge of $F$ on $R$ " (no longer discrete).

- It determines both conditional entropy and mutual information $(f \leq 1 / 2)$

$$
J(F ; R)=I(F, R)=1+\int\left(\lambda_{+} \log _{2} \lambda_{+}+\lambda_{-} \log _{2} \lambda_{-}\right) P(u, v)
$$

where $\lambda_{ \pm}=\left(1 \pm \sqrt{u^{2}+v^{2}}\right) / 2$ (no discord).

- It satisfies a backward recursion relation (harder to solve).

$$
P^{\prime}(u, v)=\int\left(1+u_{1} u_{2}\right) \delta\left(u-\frac{u_{1}+u_{2}}{1+u_{1} u_{2}}\right) \delta\left(v-\frac{v_{1} v_{2}}{1+u_{1} u_{2}}\right) P\left(u_{1}, v_{1}\right) P\left(u_{2}, v_{2}\right)
$$

## Order parameter distribution

Plots of $\log _{10} P(u, v)(f=1 / 8, N=64)$ :


- In the QD phase, $P(u, v)$ is concentrated on the unit circle.
- In the encoding phase, $P(u, v)$ is concentrated on $(0,0)$.
- Nontrivial in the mixed phase.


## The order parameter (Model 1)

Consider a product of Pauli's (Pauli string) $P$ on $F$. $U^{\dagger} P U$ then a acting on $A$ and $N-1$ recruits. Contracting the recruits gives an operator on $A$ :

$$
Q_{A}[P]:=\left\langle 0^{N-1}\right| U^{\dagger} P U\left|0^{N-1}\right\rangle \in\{I, X, Y, Z, 0\}
$$

It is not hard to show that

$$
\mathbf{s}:=\left\{Q_{A}[P]: P \text { Pauli string on } F\right\} \cap\{I, X, Y, Z\}
$$

is a subgroup of $\{I, X, Y, Z\} \simeq \mathbb{Z}_{2}^{2}$, i.e. $\mathbf{s} \in\{\mathbf{n}, \mathbf{z}, \mathbf{x}, \mathbf{y}, \mathbf{a}\}$ where

$$
\mathbf{n}=\{I\}, \mathbf{z}=\{I, Z\}, \mathbf{x}=\{I, X\}, \mathbf{y}=\{I, Y\}, \mathbf{a}=\{I, Z, X, Y\}
$$

## The order parameter

The subgroup

$$
\mathbf{s}:=\left\{Q_{A}[P]: P \text { Pauli string on } F\right\} \cap\{I, X, Y, Z\}
$$

tells us what $F$ knows about $R$ (since $A R$ form an EPR pair):

- $\mathbf{s}=\mathbf{n}: F$ and $R$ are uncorrelated, $I(F, R)=0$.
- $\mathbf{s}=\mathbf{z}$ : some operator on $F$ is perfectly correlated with $Z_{R}$ : $I(F, R)=J(R ; F)=1$.
- $\mathbf{s}=\mathbf{a}$ : we can distill from $F$ a qubit that is maximally entangled with $R: I(F, R)=2(J=1)$.
The order parameter is the probability distribution of $\mathbf{s}$,

$$
\pi:=\left(\pi_{\mathbf{n}}, \pi_{\mathbf{z}}, \pi_{\mathbf{x}}, \pi_{\mathbf{y}}, \pi_{\mathbf{a}}\right)
$$

where $\pi_{\mathbf{n}}=\operatorname{Prob}(\mathbf{s}=\mathbf{n})$ etc.

## Calculation on a tree by example

Let's compute $Q_{A}[P]$ where $P=\prod_{i \in F} Z_{i}$ in a fixed realization.
We contract with the recruit at each branching. The calculation is a "backward recursion" (from top to bottom) and follows the table at each branching:


## Calculation on a tree by example

We can compute $\mathbf{s}$ in a fixed realization by a similar backward recursion ( $\mathbf{n}$ is omitted in the figure):


|  | $\mathbf{n}$ | $\mathbf{z}$ | $\mathbf{x}$ | $\mathbf{y}$ | $\mathbf{a}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{n}$ | $\mathbf{n}$ | $\mathbf{Z}$ | $\mathbf{n}$ | $\mathbf{n}$ | $\mathbf{z}$ |
| $\mathbf{Z}$ | $\mathbf{Z}$ | $\mathbf{z}$ | $\mathbf{z}$ | $\mathbf{z}$ | $\mathbf{z}$ |
| $\mathbf{x}$ | $\mathbf{n}$ | $\mathbf{z}$ | $\mathbf{x}$ | $\mathbf{y}$ | $\mathbf{a}$ |
| $\mathbf{y}$ | $\mathbf{n}$ | $\mathbf{z}$ | $\mathbf{y}$ | $\mathbf{x}$ | $\mathbf{a}$ |
| $\mathbf{a}$ | $\mathbf{z}$ | $\mathbf{z}$ | $\mathbf{a}$ | $\mathbf{a}$ | $\mathbf{a}$ |

Note: a 1-site Clifford permutes $\mathbf{z}, \mathbf{x}, \mathbf{y}$ and leaves $\mathbf{a}$ and $\mathbf{n}$ invariant.

## Recursion of order parameter

We can easily turn these rules to a recursion relation between $\pi$ of a $t+1$-generation tree and that of a $t$-generation one:

$$
\pi(t+1)=M(\pi(t))
$$

$M: \mathbb{R}^{5} \rightarrow \mathbb{R}^{5}$ is a nonlinear (quadratic) map that depends on $p$. (See paper for explicit expression.) Initial condition depends on $f=|F| / N$ :

$$
\pi(t=0)=(1-f, 0,0,0, f)
$$

At this point, solving the model amounts to analyzing the asymptotic $(t \rightarrow \infty)$ behavior of the "flow" generated by $M$.

## Phase diagram


$\pi(t \rightarrow \infty)$ is explicitly predicted everywhere (see paper).

## More plots with finite size numerics



## Flow and symmetries



- $\mathbb{Z}_{2}$ : swapping $\mathbf{n} \leftrightarrow \mathbf{a}$, or $f \rightarrow 1-f$. Explicitly broken by the initial condition $(f \neq 1 / 2)$. Spontaneously broken in the encoding and mixed phase, restored in the QD phase.
- $\mathcal{S}_{3}$ : permuting $\mathbf{z}, \mathbf{x}, \mathbf{y}$. Explicitly broken by the branching isometry. Broken in the QD and mixed phase ( $\pi_{\mathrm{z}}>\pi_{\mathrm{x}}=\pi_{\mathbf{y}}$ ), restored in the encoding phase.


## Annealed mutual information (replica trick)

$$
\begin{aligned}
& I^{(2)}(F, R):=\log _{2} \operatorname{Tr}\left[\overline{\rho_{F R}^{2}}\right]-\log _{2} \operatorname{Tr}\left[\overline{\rho_{F}^{2}}\right]+1,
\end{aligned}
$$

## Mixed phase: joint order parameter

Two fractions $F \subset G,|F|=f N,|G|=g N$. Joint distribution of order parameter $\Pi_{s t}$ :


## QDET vs measurement-induced phase transitions (MIPT)


(a,b) QDET (Model 1). (c,d) MIPT: a sharp transition requires full access to the environment (measurement ancillas).

