

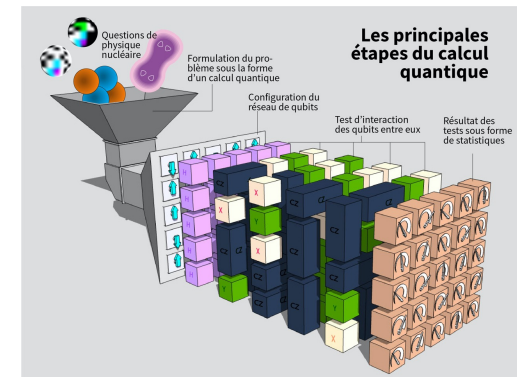
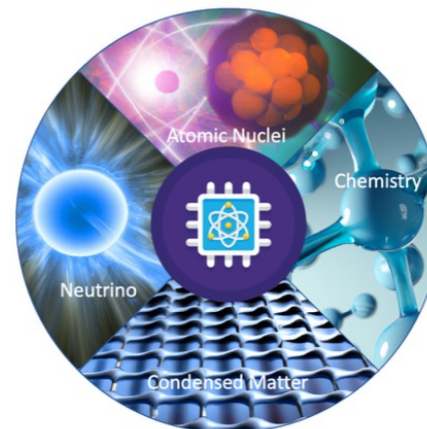
Quantum computing

Denis Lacroix (IJCLab, Paris-Saclay University)

General introduction on
Quantum computing today



Current Challenges and applications



(First part)

General aspects of quantum computers and quantum advantage



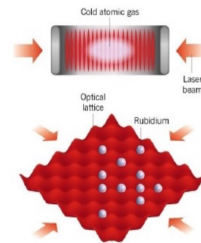


What means quantum devices today

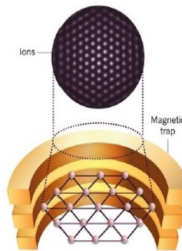
There are many types of quantum computers: ***analog versus digital quantum computers***

There are now many quantum objects one can manipulate

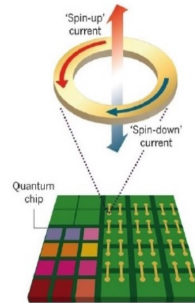
Cold atoms



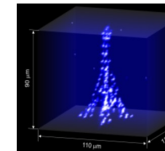
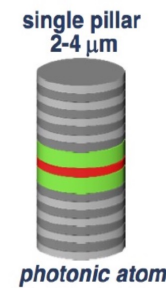
Trapped ions



Superc. loops



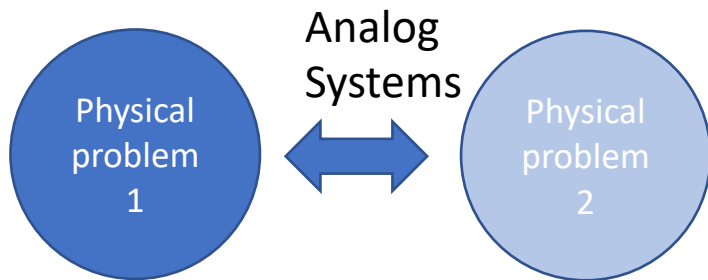
Polaritons



Atomes de Rydberg dans des pinces optiques

Analog quantum simulator

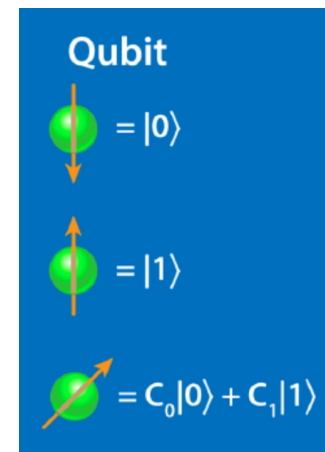
Digital quantum simulator



Complex problem that cannot or hardly be simulated on classical computers

Analog problem to 1 that could be tested in laboratory

➔ Non-universal



➔ Universal Quantum simulation

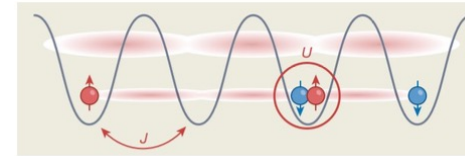


Quantum analogue simulation

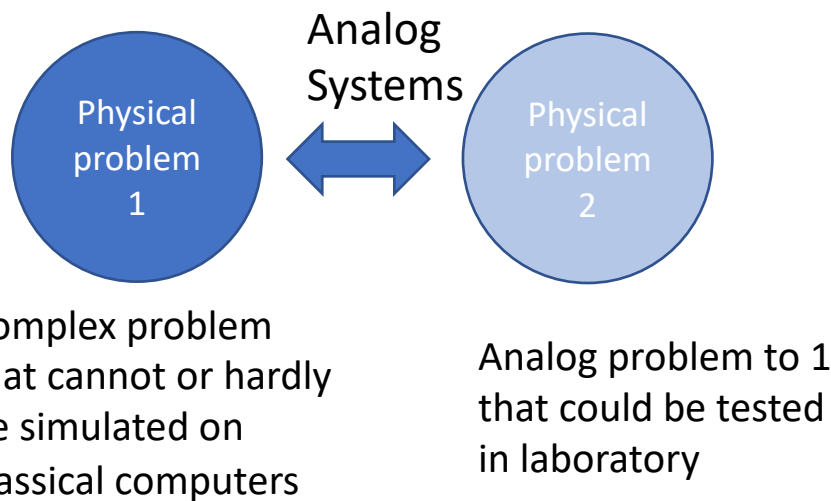
There are many types of quantum computers: ***analog versus digital quantum computers***

A few examples

Analog systems on lattice (Fermi-Hubbard, Schwinger model, ...)

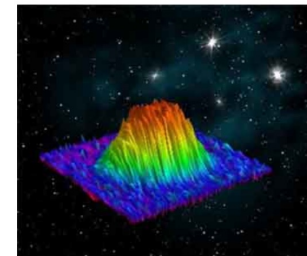


Analog quantum simulator

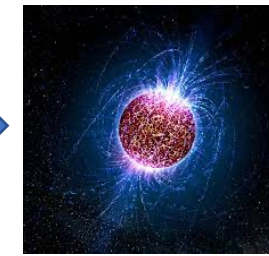


Analog simulation of astrophysics/cosmology

Ultracold Fermi gas

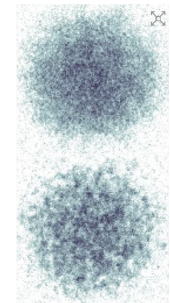


Neutron stars

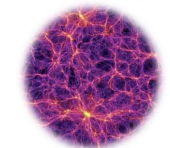


An expanding universe is simulated in a quantum droplet
22 Mar 2023 Campbell McLauchlan

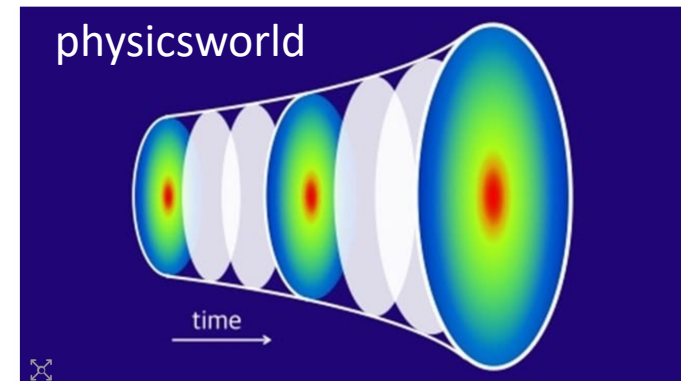
Viermann et al.
Nature



Ultracold atoms



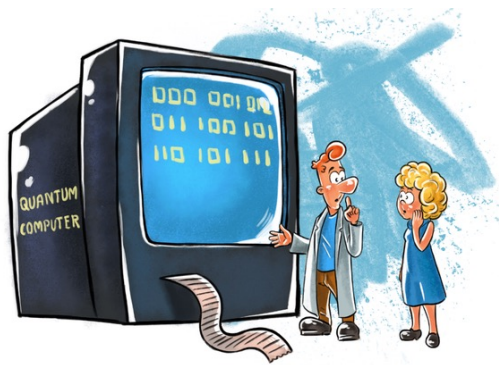
Dark matter



- ➔ The mapping from one physical problem to another physical problem is a delicate issue
- ➔ It strongly depends on the problem itself (non-universality)

Digital quantum devices are supposed to be universal

Classical computers
Works with bits



Bits are only 0 or 1

Obvious advantage

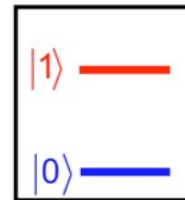
And with many
Qubits

Quantum computers with
Quantum bits

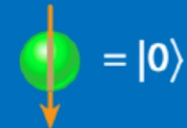
Qubits can be seen
As two-level systems

qubit

2 level system



Qubit

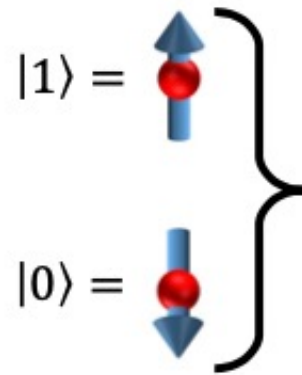
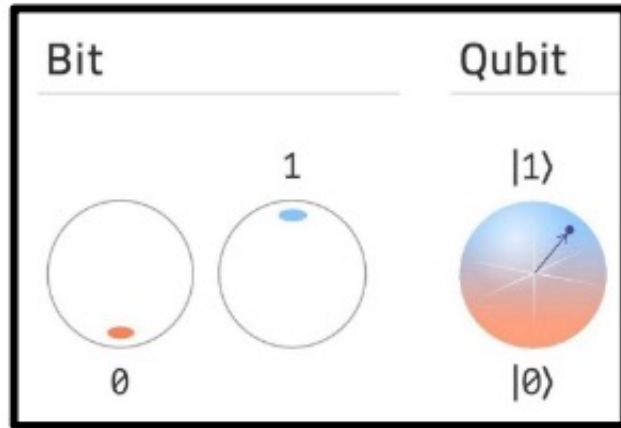


A single Qubit can be any superposition of 0 and 1

New aspects can be used like quantum interference and entanglement

Where do we expect quantum advantage?

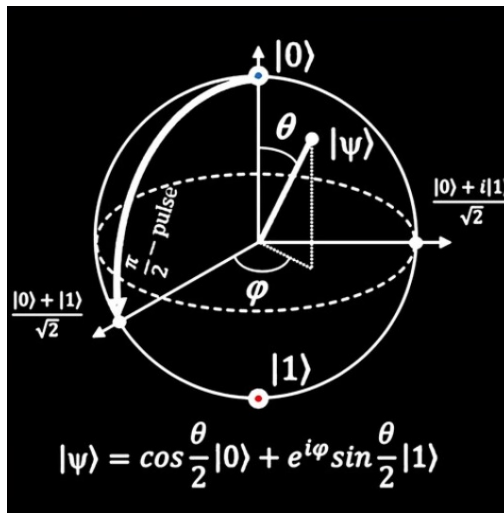
A few examples



Quantum RAM advantage

Storing with
-1 bit: 2 integers 0,1
-2 bits: 4 integers 0,1,2,3
...

Bloch sphere picture



Storing with one qubit
 $|\Phi\rangle = f_1(\theta)|0\rangle + f_2(\theta, \phi)|1\rangle$
Two-qubits
 $|\Phi\rangle = f_1(\theta_1, \theta_2)|00\rangle + \dots$
...

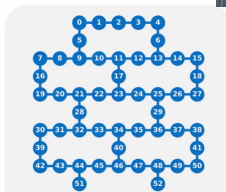


Direct multi-parameters function encoding
Continuous function programming

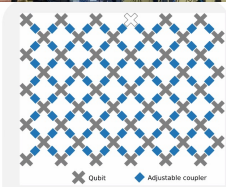
(+quantum speedup of runtime – see later)

Quantum Computation and Quantum Information

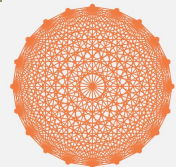
MICHAEL A. NIELSEN and ISAAC L. CHUANG



IBM
Rochester 53 qubits, October 2019
65 qubits, October 2020

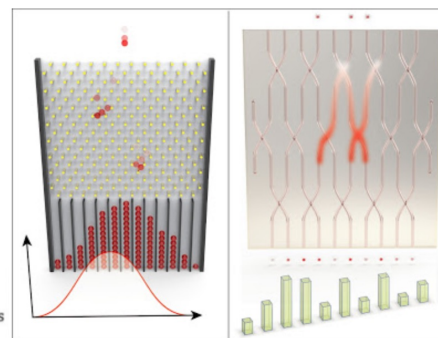


Google
Sycamore 53 qubits, October 2019

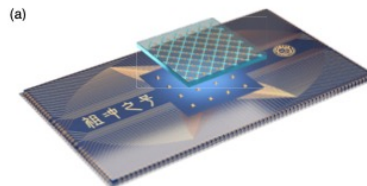


IonQ
11 qubits, 2018
32 qubits, 2020

Quantum computational advantage using photons, Science 370 (2020)



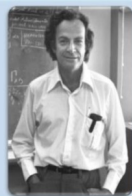
Zuchongzhi (66 qubits)
(PRL October 2021)



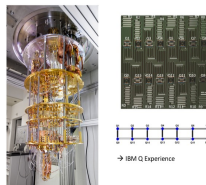
Simulating physics with computers-1982

Richard P. Feynman (Nobel Prize in Physics 1965)

"Nature isn't classical, dammit, and if you want to make a simulation of nature, you'd better make it quantum mechanical, and by golly it's a wonderful problem, because it doesn't look so easy."



IBM QX5 (16 qubits)



128 qubits
Rigetti

72 qubits
Google

1152 qubits
DWave

2048 qubits
DWave

512 qubits
DWave

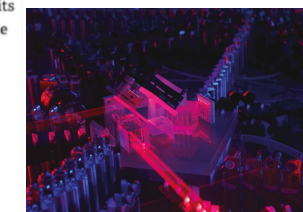
12 qubits
MIT

7 qubits
Los Alamos

128 qubits
DWave

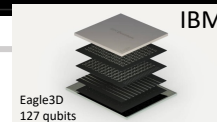
50 qubits
IBM

17 qubits
IBM



This photonic computer did a task that would take a classical computer 2.5 billion years.

Scaling IBM Quantum technology					
IBM Q System One (Current)		IBM Q System Two (Development)		Next family of IBM Quantum systems	
Year	Qubits	Year	Qubits	Year	Qubits
2019	27 qubits Falcon	2020	65 qubits Hummingbird	2021	127 qubits Eagle
				2022	433 qubits Osprey
				2023	1,121 qubits Condor
				and beyond	Path to 1 million qubits and beyond Long-stack systems



Eagle3D
127 qubits

(2020) (2021)

IonQ Gemini desk computer

Quantum supremacy using a programmable superconducting processor

Nature | Vol 574 | 24 OCTOBER 2019 | 505

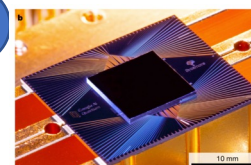
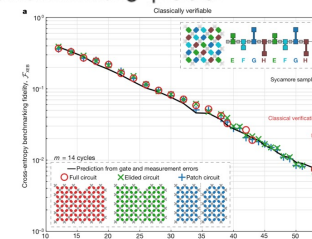


Fig. 1 | The Sycamore processor. a, Layout of processor, showing a rectangular array of 54 qubits (grey), each connected to its four nearest neighbours with couplers (blue). The inoperable qubit is outlined. b, Photograph of the Sycamore chip.



Quantum Theory

1927

Quantum Computer

1982

55 YEARS

18 YEARS

6 YEARS

1 YEAR

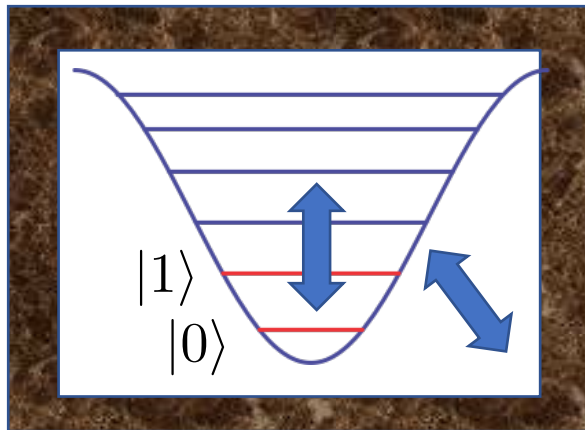


?

More on quantum computers
and quantum hardwares

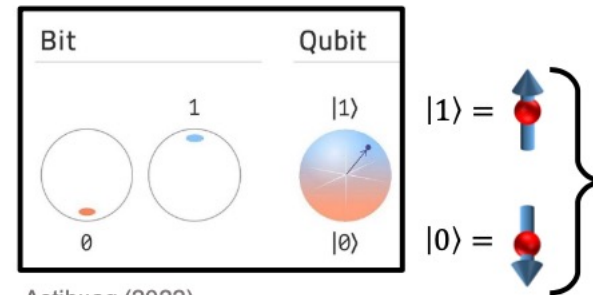
Quantum computing today is firstly a technological challenge

Simplified view of a general quantum system



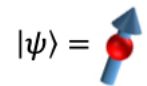
Challenges

➔ Be able to isolate and manipulate accurately a qubit



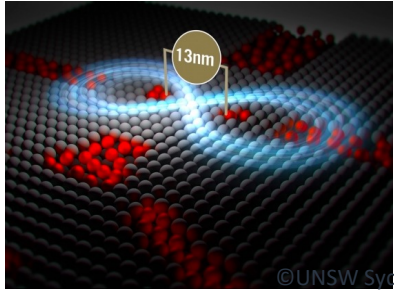
Astibuag (2022)

➔ Be able to manipulate accurately a qubit (fidelity)



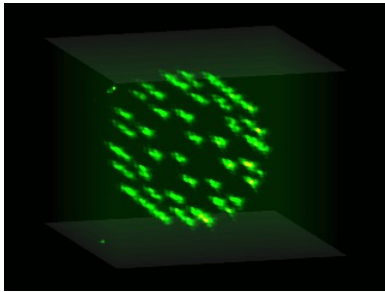
➔ Be able to increase the number of “connected” qubits (scalability, connectivity)

Many technologies are explored

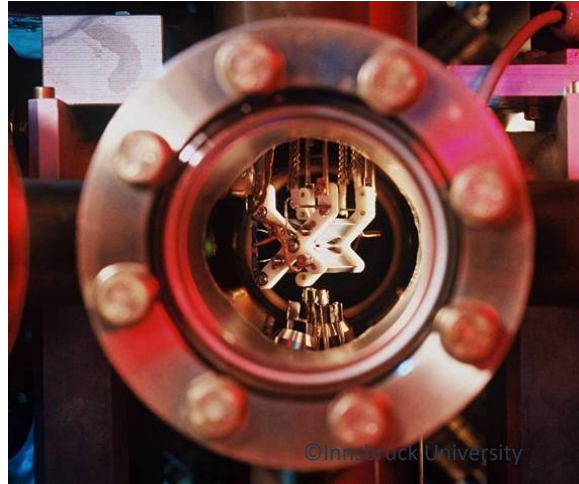


©UNSW Sydney

Silicon qubits

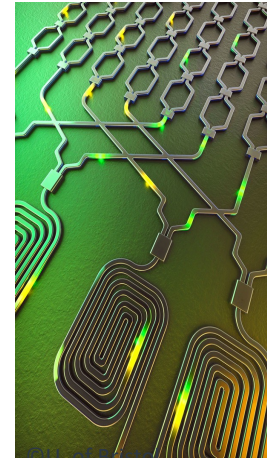


Neutral atoms



©Innsbruck University

Trapped ions



©U. of Bristol

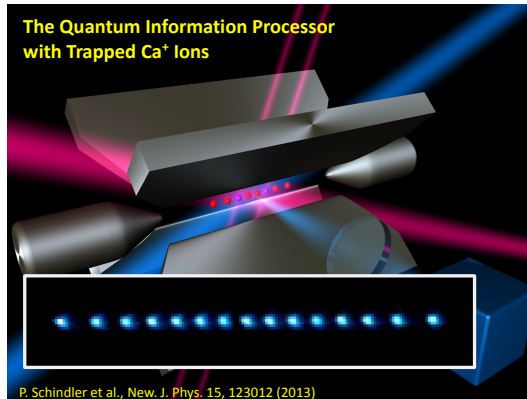
Photons



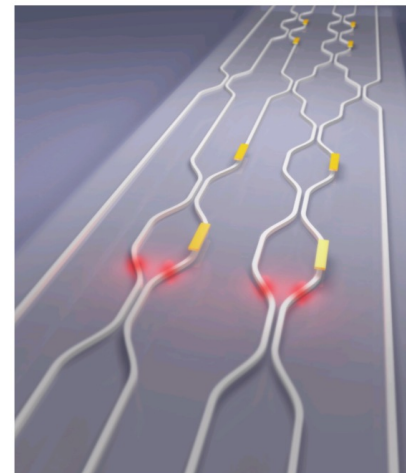
©Google

Superconducting qubits

NMR



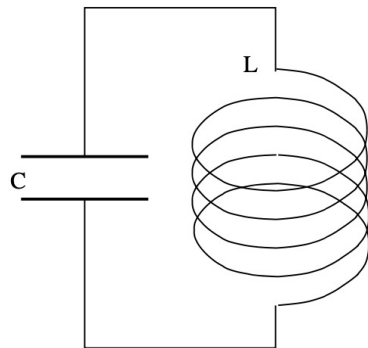
P. Schindler et al., New J. Phys. 15, 123012 (2013)



Example 1: Superconducting qubits



Simple oscillator by LC circuit



$$H = \frac{1}{2L}\phi^2 + \frac{1}{2C}n^2$$

ϕ : flux in the inductor

n : charge in the capacitor

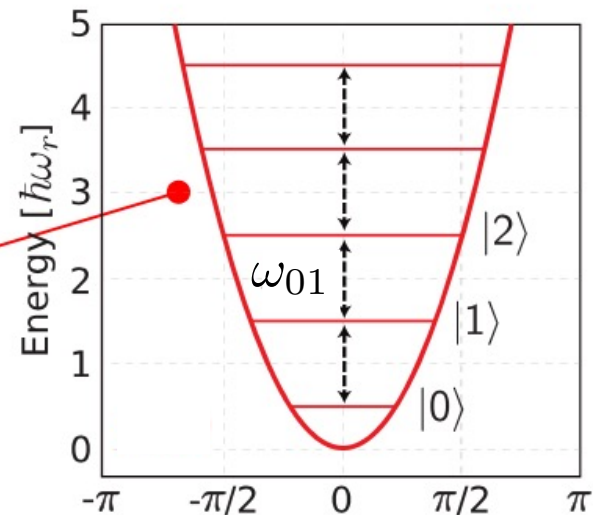
Methodology

➔ We consider the limit where quantum effects are important

$$k_B T \ll \hbar\omega_{01}$$

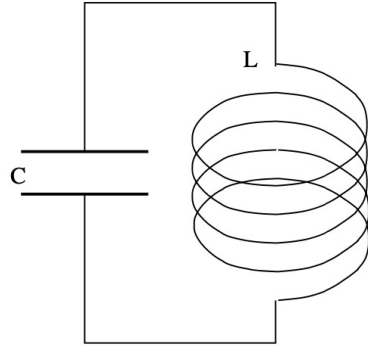
➔ n, ϕ becomes equivalent to canonical conjugated variables like (x, p) for the harmonic oscillator (HO)

Quantum
HO





Simple oscillator by LC circuit



$$H = \frac{1}{2L}\phi^2 + \frac{1}{2C}n^2$$

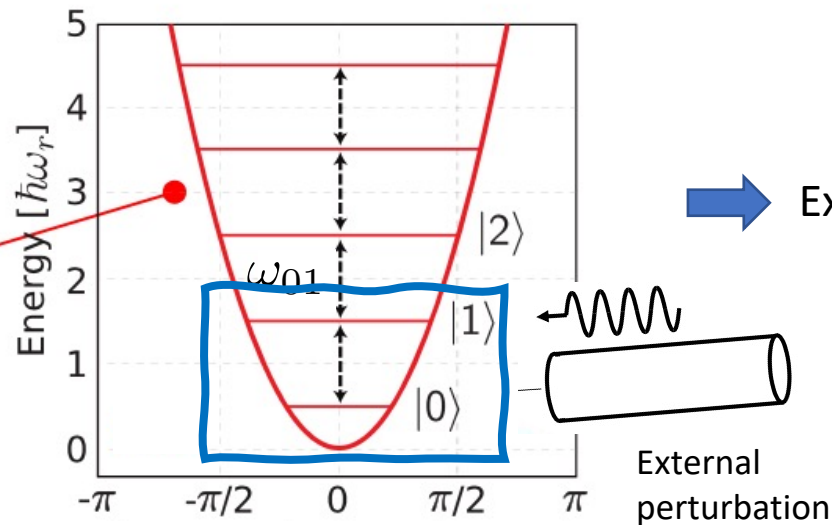
ϕ : flux in the inductor

n : charge in the capacitor

➔ We are interested in isolating $|0\rangle$ and $|1\rangle$

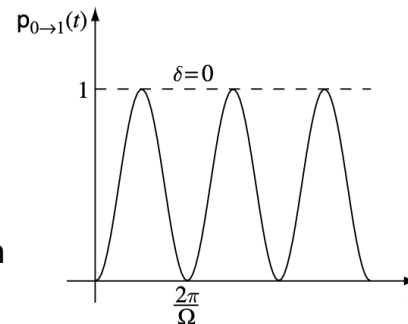
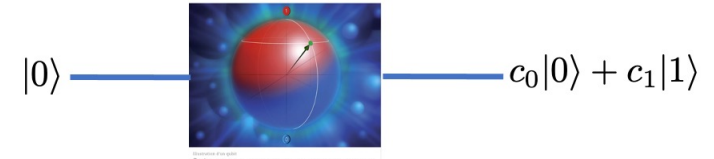
➔ We should prepare the system in any combination of them

Qubit identification and manipulation



Quantum HO

➔ Example: induce Rabi-like oscillation



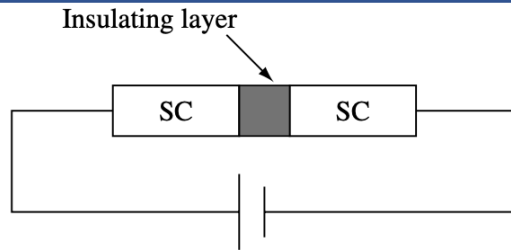
Two difficulties
Approximate operations

$|0\rangle \longrightarrow |1\rangle$

Unwanted transitions

$|1\rangle \longrightarrow |2\rangle$

The Josephson junction technology



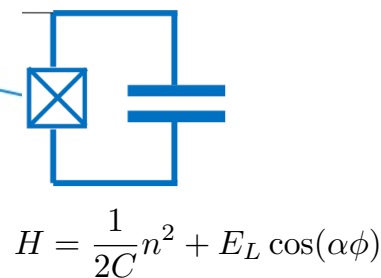
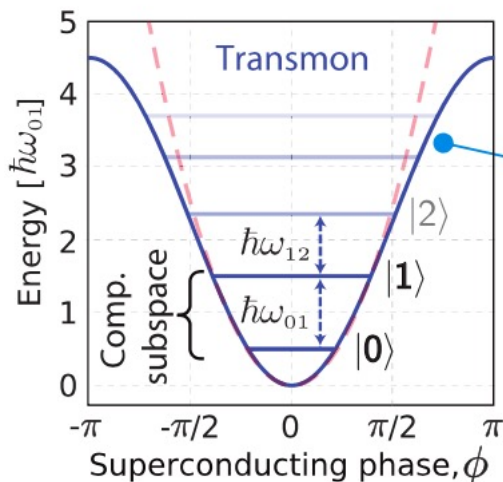
ϕ : relative phase of the two superconductors
 n : number of cooper pairs

$$H = \frac{1}{2C}n^2 + E_L \cos(\alpha\phi) + \frac{1}{2L}(\phi - \phi_{\text{ex}})^2 - I_{\text{ex}}\phi$$

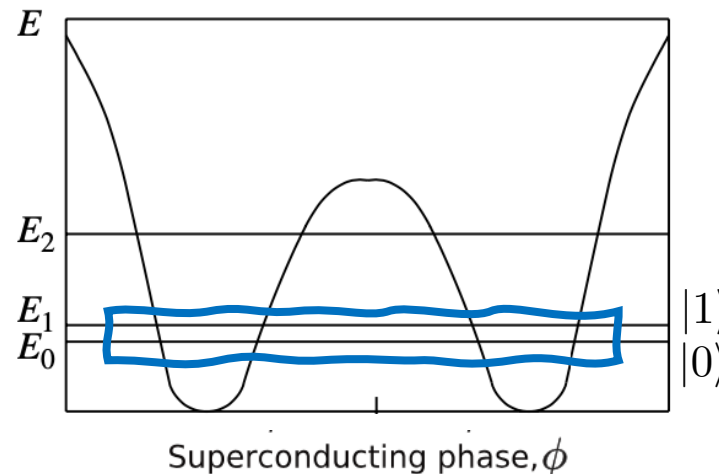
Equivalent to LC but anharmonic

with external current

Transmon



With an external current:

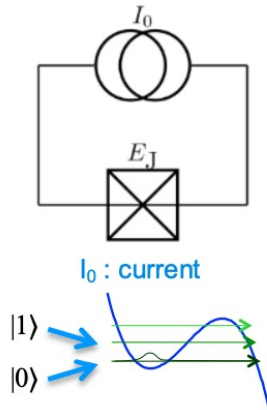


➔ Anharmonicity and gaps helps to reduce unwanted transitions

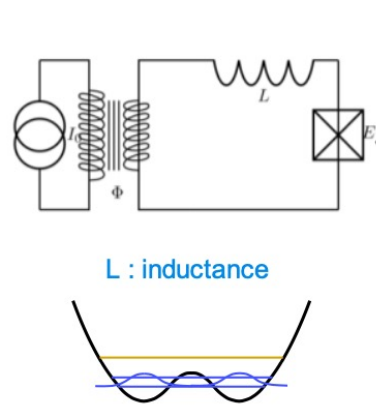
The Josephson junction technology



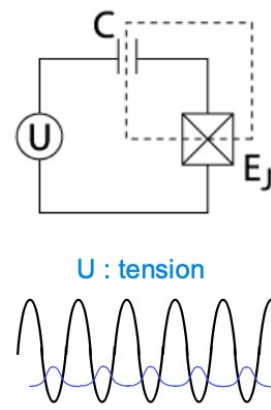
phase qubit



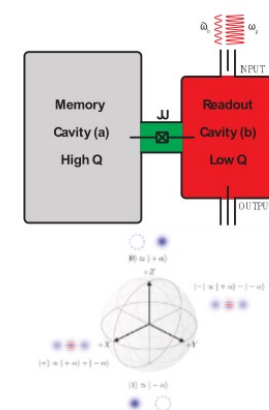
flux qubit



charge qubit - transmon



cat-qubits



Josephson junctions handle the qubit degree of liberty

Josephson junctions prepare, couple and correct the cat-qubits

|0> and |1> qubits two energy levels in a potential well

quantum gates micro-waves

qubits readout resonator and micro-waves

commercial vendors **abandoned**

two superconducting current directions

magnetic field

magnetometer (SQUID)

two levels of charge of Cooper pairs

micro-waves

resonator and micro-waves

pairs of entangled microwave photons in a cavity

micro-waves

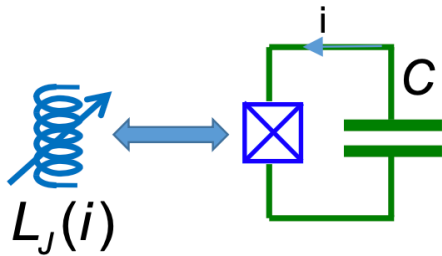
resonator and micro-waves



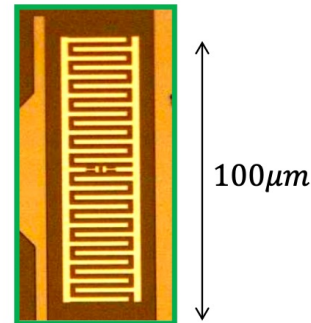
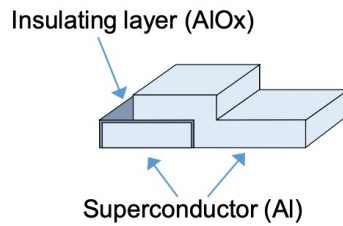
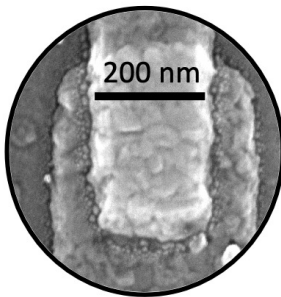
(cc) Olivier Ezratty, 2022

The Josephson junction technology

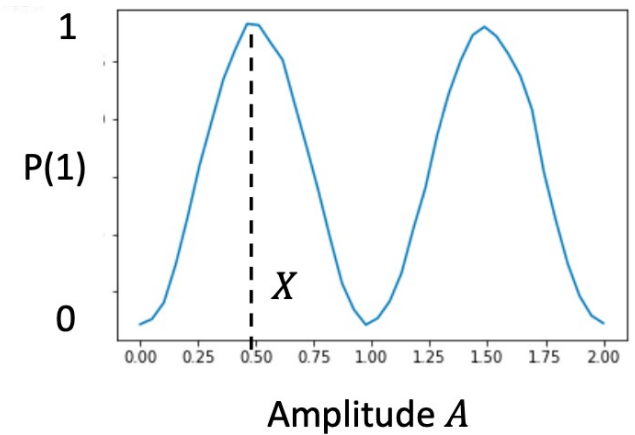
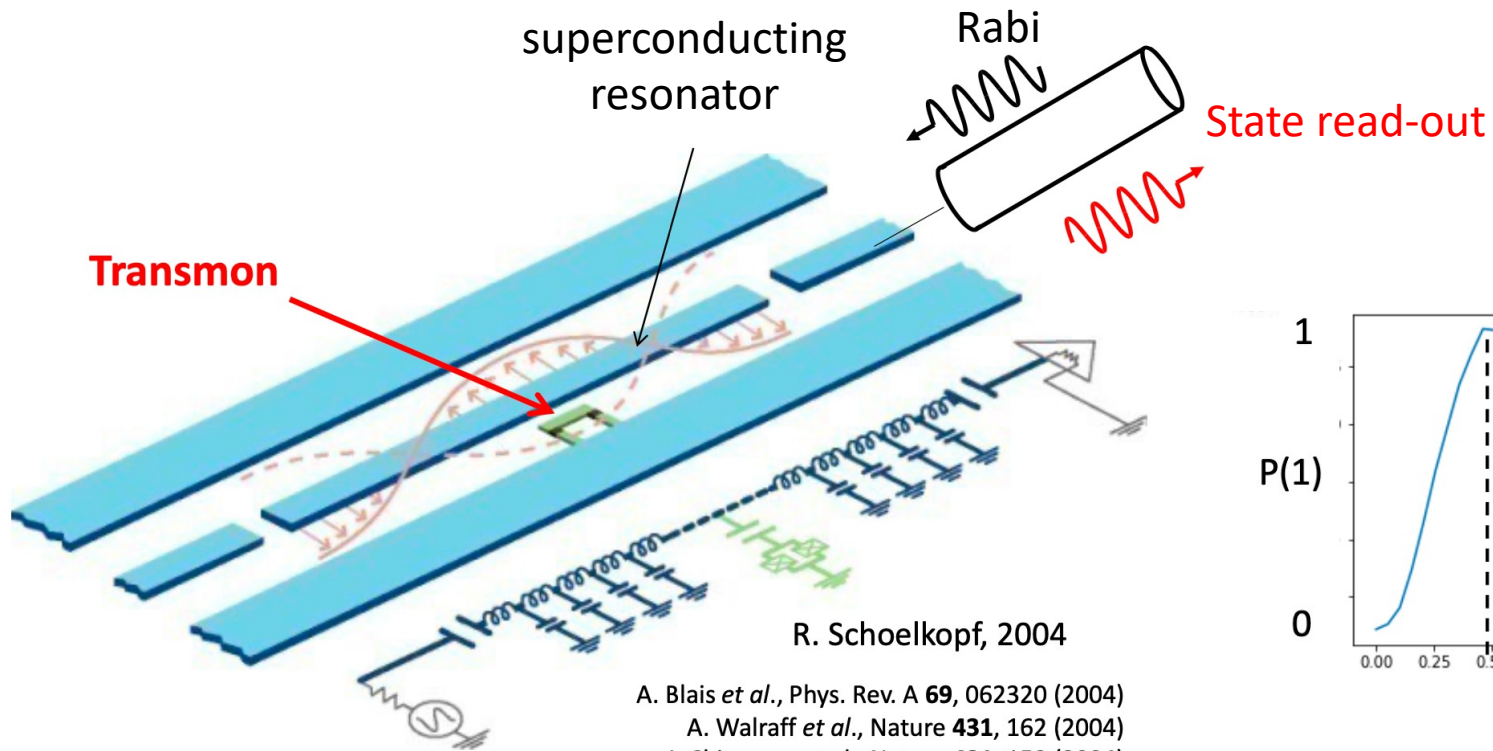
What it looks like



$$\nu_{01} \sim 5 - 10 \text{ GHz}$$
$$T = 10 \text{ mK}$$

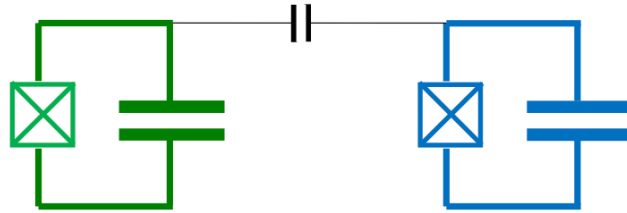


The Josephson junction technology



Next step: putting several qubits together

2 qubits can be coupled through electrostatic interactions



With this one can manipulate/entangle qubits



Some specific
Operations
(see next lecture)

When

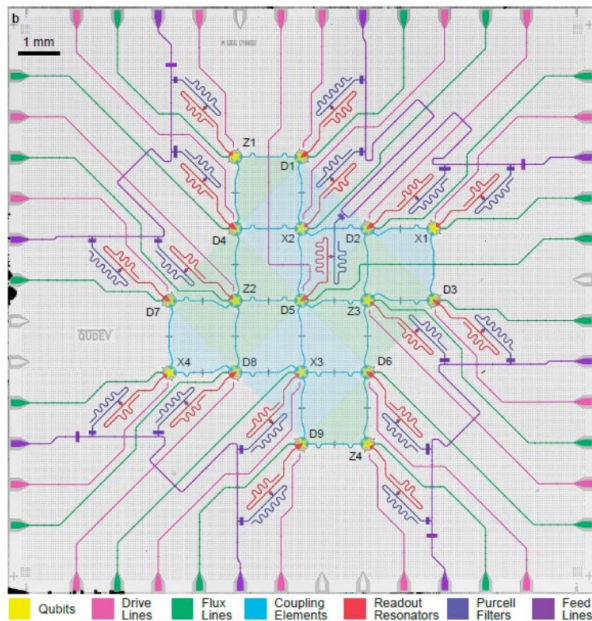
Fidelity

Acronym ^b	Layout ^c	First demonstration [Year]	Highest fidelity [Year]	Gate time
CZ (ad.)	T-T	DiCarlo et al. (72) [2009]	99.4% ^e Barends et al. (3) [2014]	40 ns
			99.7% ^e Kjaergaard et al. (73) [2020]	60 ns
\sqrt{i} SWAP	T-T	Neeley et al. (81) ^d [2010]	90% ^g Dewes et al. (74) [2014]	31 ns
CR	F-F	Chow et al. (75) ^h [2011]	99.1% ^e Sheldon et al. (5) [2016]	160 ns
\sqrt{b} SWAP	F-F	Poletto et al. (76) [2012]	86% ^g Poletto et al. (76) [2012]	800 ns
MAP	F-F	Chow et al. (77) [2013]	87.2% ^g Chow et al. (75) [2011]	510 ns
CZ (ad.)	T-(T)-T	Chen et al. (55) [2014]	99.0% ^e Chen et al. (55) [2014]	30 ns
RIP	3D F	Paik et al. (78) [2016]	98.5% ^e Paik et al. (78) [2016]	413 ns
\sqrt{i} SWAP	F-(T)-F	McKay et al. (79) [2016]	98.2% ^e McKay et al. (79) [2016]	183 ns
CZ (ad.)	T-F	Caldwell et al. (80) [2018]	99.2% ^e Hong et al. (6) [2019]	176 ns
CNOT _L	BEQ-BEQ	Rosenblum et al. (13) [2018]	~99% ^f Rosenblum et al. (13) [2018]	190 ns
CNOT _{T-L}	BEQ-BEQ	Chou et al. (82) [2018]	79% ^g Chou et al. (82) [2018]	4.6 μ s

Where we are now ?

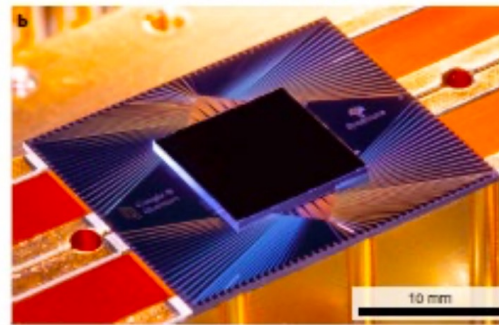
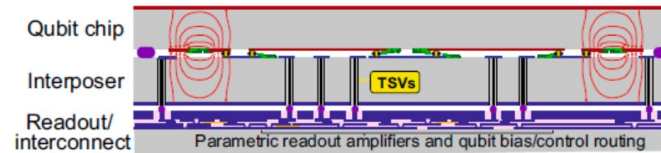


2D chip : 17 qubits



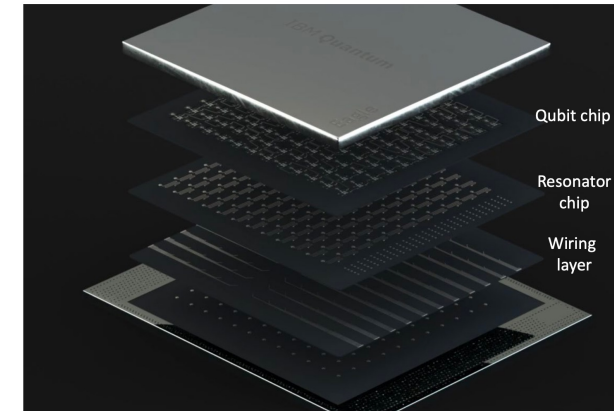
S. Krinner et al., arxiv (2021)
Wallraff group, ETHZ

3D integration ~100 qubits demonstrated

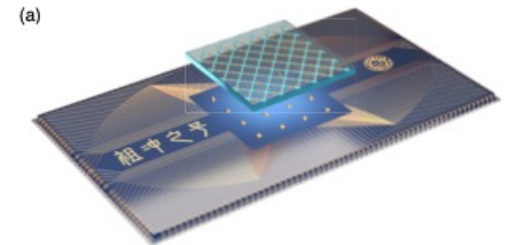


Rosenberg et al., NPJ (2017)
Arute et al., Nature (2019)

IBM Eagle
127 qubits



Zuchongzhi (66 qubits)
(PRL October 2021)

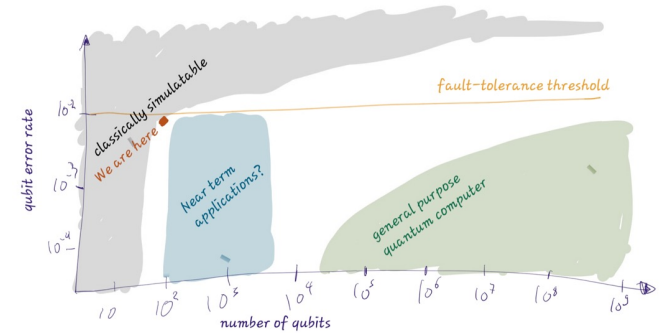


Quantum advantage

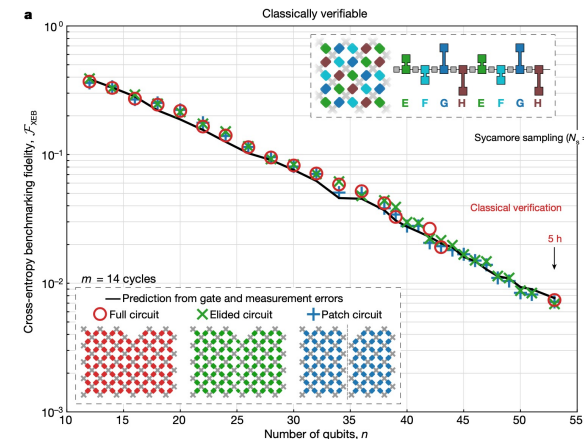
Article

Quantum supremacy using a programmable superconducting processor

The promise of quantum computers is that certain computational tasks might be executed exponentially faster on a quantum processor than on a classical processor¹. A fundamental challenge is to build a high-fidelity processor capable of running quantum algorithms in an exponentially large computational space. Here we report the use of a processor with programmable superconducting qubits²⁻⁷ to create quantum states on 53 qubits, corresponding to a computational state-space of dimension 2^{53} (about 10^{16}). Measurements from repeated experiments sample the resulting probability distribution, which we verify using classical simulations. **Our Sycamore processor takes about 200 seconds to sample one instance of a quantum circuit a million times—our benchmarks currently indicate that the equivalent task for a state-of-the-art classical supercomputer would take approximately 10,000 years.** This dramatic increase in speed compared to all known classical algorithms is an experimental realization of quantum supremacy⁸⁻¹⁴ for this specific computational task, heralding a much-anticipated computing paradigm.



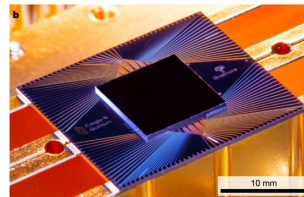
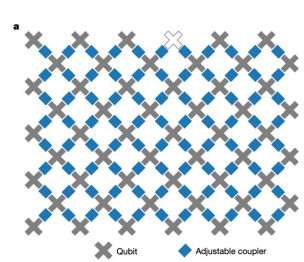
Proofs of fidelity



Task

Compute output probabilities of a “random” circuit, i.e. a set of random operation.

➡ Usefulness of the task ?

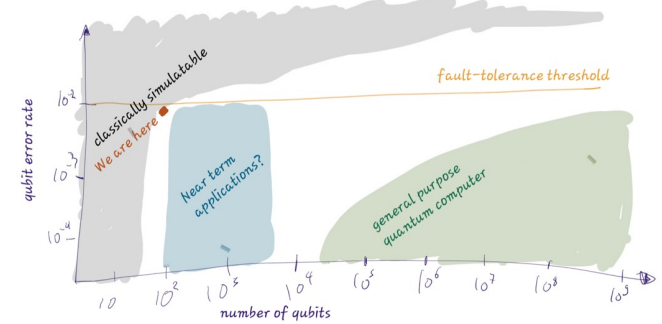


Quantum advantage

Article

Quantum supremacy using a programmable superconducting processor

The promise of quantum computers is that certain computational tasks might be executed exponentially faster on a quantum processor than on a classical processor¹. A fundamental challenge is to build a high-fidelity processor capable of running quantum algorithms in an exponentially large computational space. Here we report the use of a processor with programmable superconducting qubits²⁻⁷ to create quantum states on 53 qubits, corresponding to a computational state-space of dimension 2^{53} (about 10^{16}). Measurements from repeated experiments sample the resulting probability distribution, which we verify using classical simulation about 200 seconds to sample one instance of a quantum circuit. Current benchmarks currently indicate that the equivalent classical supercomputer would take approximately 10,000 years to simulate at the same speed compared to all known classical algorithms. This work demonstrates quantum supremacy⁸⁻¹⁴ for this specific computational task, marking the first time a quantum computer has performed a task beyond the reach of anticipated computing paradigm.



PHYSICAL REVIEW LETTERS **129**, 090502 (2022)

Editors' Suggestion

Solving the Sampling Problem of the Sycamore Quantum Circuits

Feng Pan^{1,2}, Keyang Chen^{1,3} and Pan Zhang^{1,4,5,*}

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²School of Physical Sciences, University of Chinese Academy of Sciences, Beijing 100049, China

³Yuanpei College, Peking University, Beijing 100871, China

⁴School of Fundamental Physics and Mathematical Sciences, Hangzhou Institute for Advanced Study, UCAS, Hangzhou 310024, China

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(Received 20 November 2021; accepted 20 July 2022; published 22 August 2022)

We study the problem of generating independent samples from the output distribution of Google's Sycamore quantum circuits with a target fidelity, which is believed to be beyond the reach of classical supercomputers and has been used to demonstrate quantum supremacy. We propose a method to classically solve this problem by contracting the corresponding tensor network just once, and is massively more efficient than existing methods in generating a large number of *uncorrelated* samples with a target fidelity. For the Sycamore quantum supremacy circuit with 53 qubits and 20 cycles, we have generated 1×10^6 *uncorrelated* bitstrings s which are sampled from a distribution $\hat{P}(s) = |\hat{\psi}(s)|^2$, where the approximate state $\hat{\psi}$ has fidelity $F \approx 0.0037$. The whole computation has cost about 15 h on a computational cluster with 512 GPUs. The obtained 1×10^6 samples, the contraction code and contraction order are made public. If our algorithm could be implemented with high efficiency on a modern supercomputer with ExaFLOPS performance, we estimate that ideally, the simulation would cost a few dozens of seconds, which is faster than Google's quantum hardware.

DOI: 10.1103/PhysRevLett.129.090502

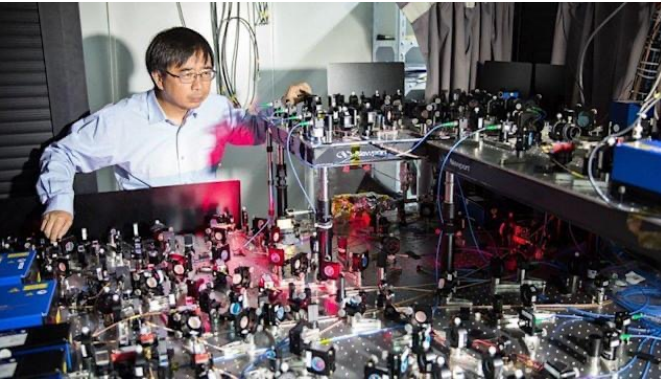
Task

Compute output probabilities of a "circuit, i.e. a set of random operations"

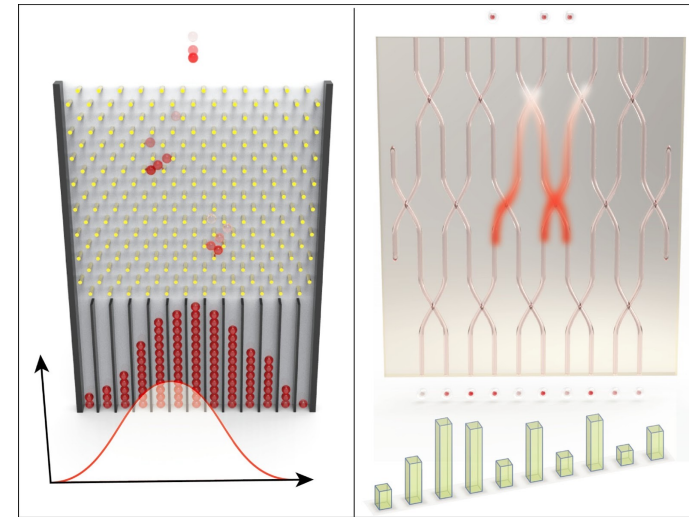
➡ Usefulness of the task ?

➡ Some controversial aspects

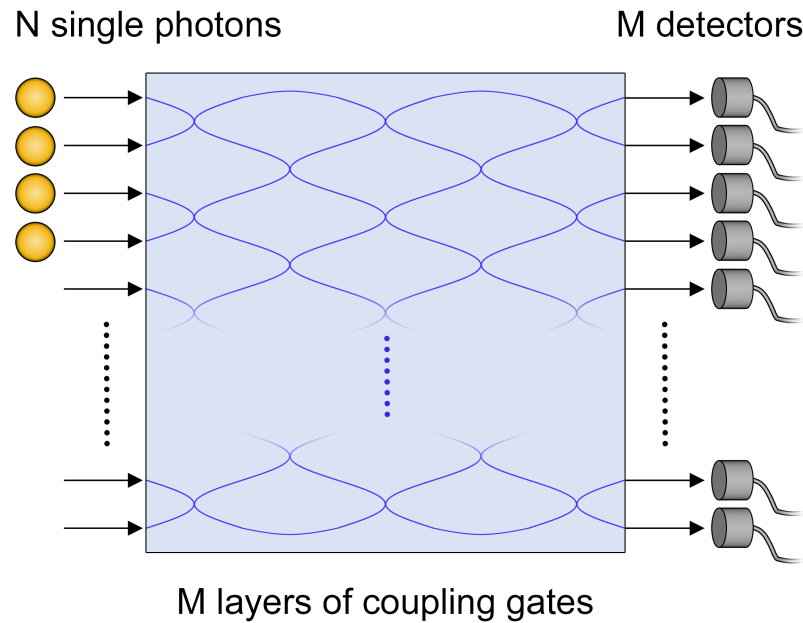
Quantum advantage



Boson sampling problem



Goal : predict the final distribution



Example of operations: beam splitter, beam shifter





















QUANTUM COMPUTING

Quantum computational advantage using photons

Han-Sen Zhong^{1,2*}, Hui Wang^{1,2*}, Yu-Hao Deng^{1,2*}, Ming-Cheng Chen^{1,2*}, Li-Chao Peng^{1,2}, Yi-Han Luo^{1,2}, Jian Qin^{1,2}, Dian Wu^{1,2}, Xing Ding^{1,2}, Yi Hu^{1,2}, Peng Hu³, Xiao-Yan Yang³, Wei-Jun Zhang³, Hao Li³, Yuxuan Li⁴, Xiao Jiang^{1,2}, Lin Gan⁴, Guangwen Yang⁴, Lixing You³, Zhen Wang³, Li Li^{1,2}, Nai-Le Liu^{1,2}, Chao-Yang Lu^{1,2,†}, Jian-Wei Pan^{1,2,†}

Quantum computers promise to perform certain tasks that are believed to be intractable to classical computers. Boson sampling is such a task and is considered a strong candidate to demonstrate the quantum computational advantage. We performed Gaussian boson sampling by sending 50 indistinguishable single-mode squeezed states into a 100-mode ultralow-loss interferometer with full connectivity and random matrix—the whole optical setup is phase-locked—and sampling the output using 100 high-efficiency single-photon detectors. The obtained samples were validated against plausible hypotheses exploiting thermal states, distinguishable photons, and uniform distribution. The photonic quantum computer, *Jiuzhang*, generates up to 76 output photon clicks, which yields an output state-space dimension of 10^{30} and a sampling rate that is faster than using the state-of-the-art simulation strategy and supercomputers by a factor of $\sim 10^{14}$.

(I) Quantum Hardware and
the NISQ (Noisy Intermediate Scale
quantum) period

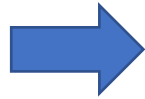
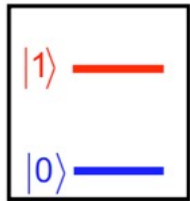
	Leading technologies in NISQ era ¹		Candidate technologies beyond NISQ		
	Superconducting ²	Trapped ion	Photonic	Silicon-based ³	Topological ⁵
 Qubit type or technology					
 Description of qubit encoding	Two-level system of a superconducting circuit	Electron spin direction of ionized atoms in vacuum	Occupation of a waveguide pair of single photons	Nuclear or electron spin or charge of doped P atoms in Si	Majorana particles in a nanowire
 Physical qubits ^{4,5}	IBM: 20, Rigetti: 19, Alibaba: 11, Google: 9	Lab environment: AQT ⁶ : 20, IonQ: 14	6×3⁹	2	target: 1 in 2018
 Qubit lifetime	~50–100 μs	~50 s	~150 μs	~1–10 s	target ~100 s
 Gate fidelity ⁷	~99.4%	~99.9%	~98%	~90%	target ~99.9999%
 Gate operation time	~10–50 ns	~3–50 μs	~1 ns	~1–10 ns	–
 Connectivity	Nearest neighbors	All-to-all	To be demonstrated	Nearest neighbor	–
 Scalability	 No major road-blocks near-term	 Scaling beyond one trap (>50 qb)	 Single photon sources and detection	 Novel technology potentially high scalability	
 Maturity or technology readiness level	 TRL ¹⁰ 5	 TRL 4	 TRL 3	 TRL 3	 TRL 1
 Key properties	Cryogenic operation Fast gating Silicon technology	Improves with cryogenic temperatures Long qubit lifetime Vacuum operation	Room temperature Fast gating Modular design	Cryogenic operation Fast gating Atomic-scale size	Estimated: Long lifetime High fidelities

NISQ (Noisy intermediate Scale Quantum) period

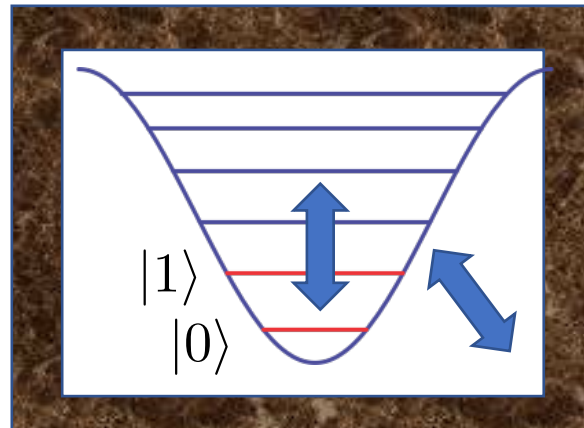
Many things tends to destroy the ideal qubits picture and the quantum coherence.

Ideal Qubits

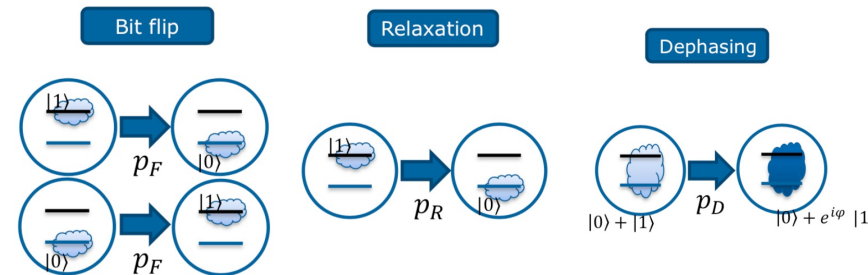
2 level system



External Exp. Setup



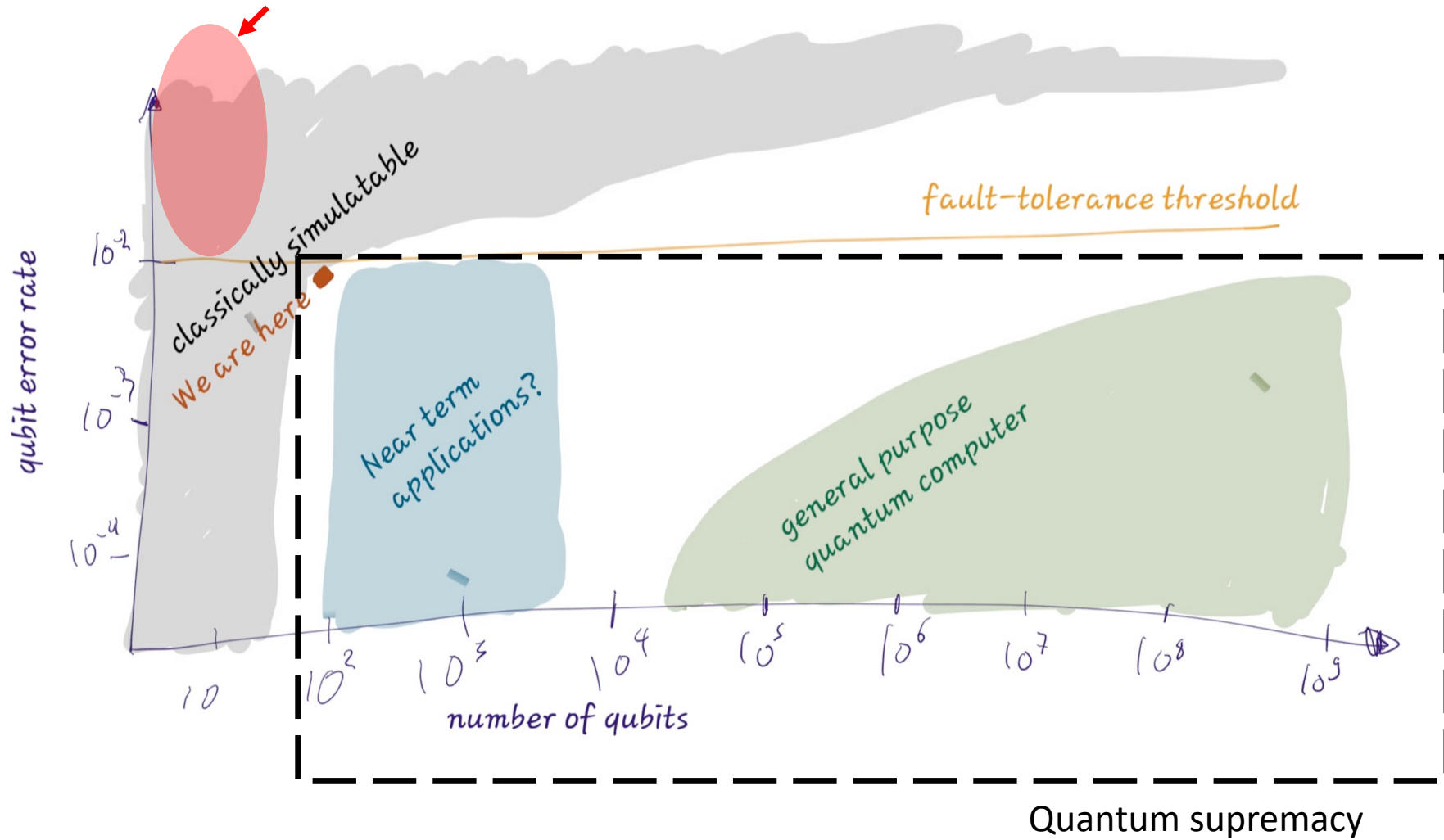
Leads to loss of information
And decoherence



For multi-qubits, also cross-talk
(making operations on one qubit
Impacts other qubits)

NISQ period: Working with quantum computers now means working in a noisy environment and “NISQ friendly” programs, i.e. only selected algorithm can be applied and error corrections should be made to get reasonable results.

Most applications are still tested here

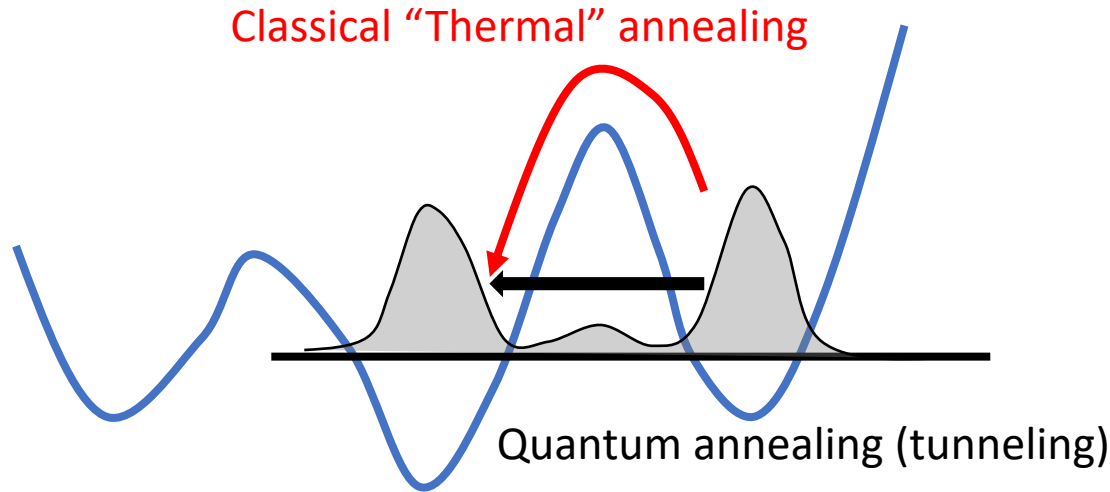


Where “quantum computing” is
expected to be disruptive?

Short introduction to bit versus Qubits

Illustration of quantum advantages

Quantum Tunneling and quantum annealing



Quantum entanglement

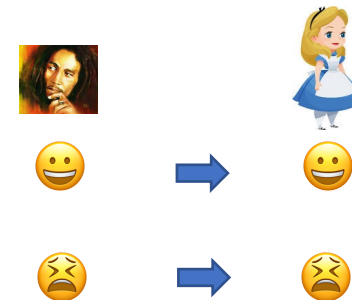
Assume two persons (Alice and Bob)



The humor of A&B are encoded in the wave-function

$$|\Phi\rangle = \alpha | \begin{matrix} \downarrow A \\ \text{😊} \end{matrix} \begin{matrix} \downarrow B \\ \text{😊} \end{matrix} \rangle + \beta | \begin{matrix} \downarrow A \\ \text{😞} \end{matrix} \begin{matrix} \downarrow B \\ \text{😞} \end{matrix} \rangle$$

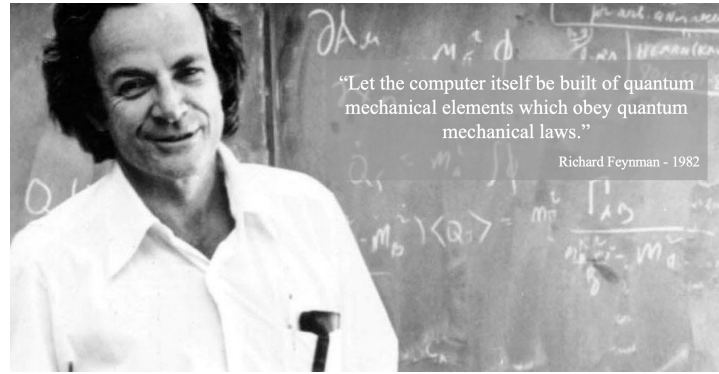
Suppose I measure Bob



I can measure partial info and get the full info
The info is destroyed after measurement

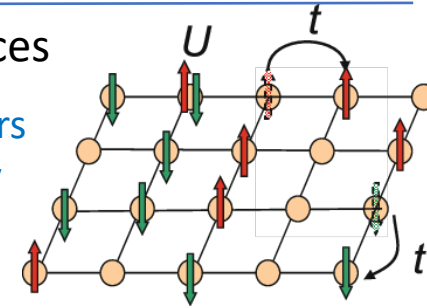
What are the anticipated applications ?

Simulation of Quantum complex systems



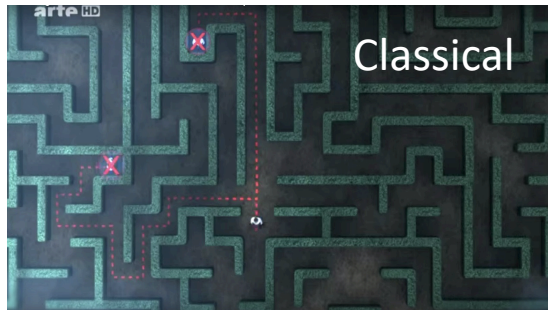
Ex: systems on lattices

On classical computers
Can be solved exactly
For max 20 particles.

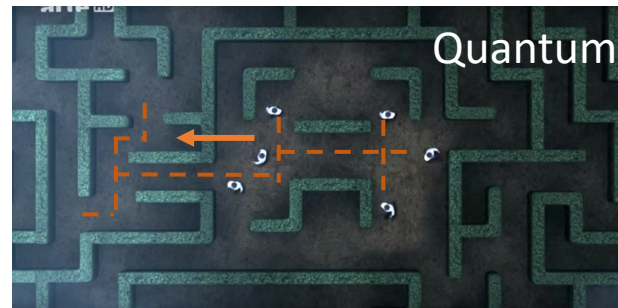


On quantum computers:
 N sites means only N qubits

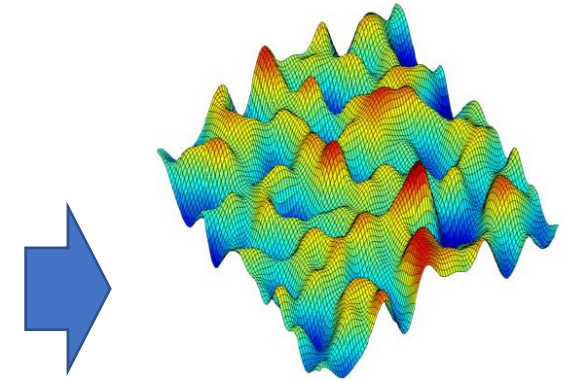
Quantum versus classical search



VS



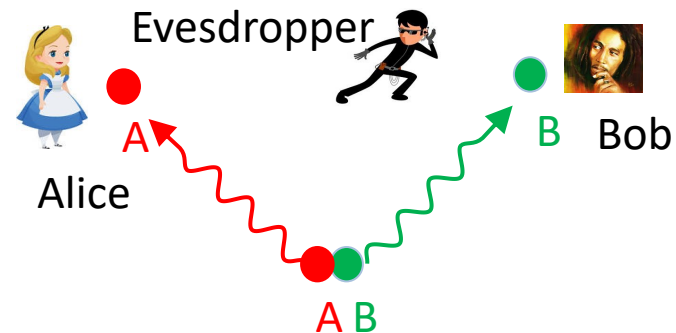
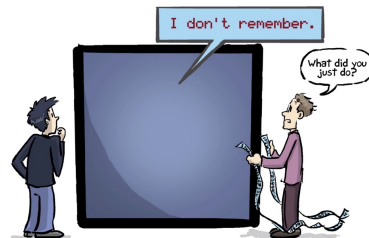
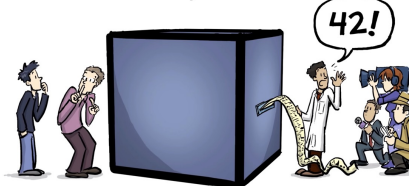
Credit: *The Fabric of The Cosmos: Quantum Leap*



Exploring complex landscape:
molecules,
customers preferences (amazon), ...

Quantum secrets (cryptography, quantum key, ...)

It's a secret computation...

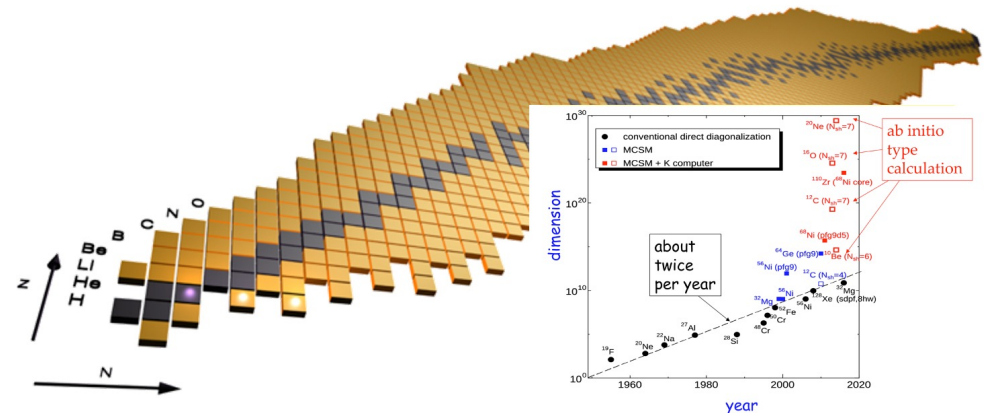
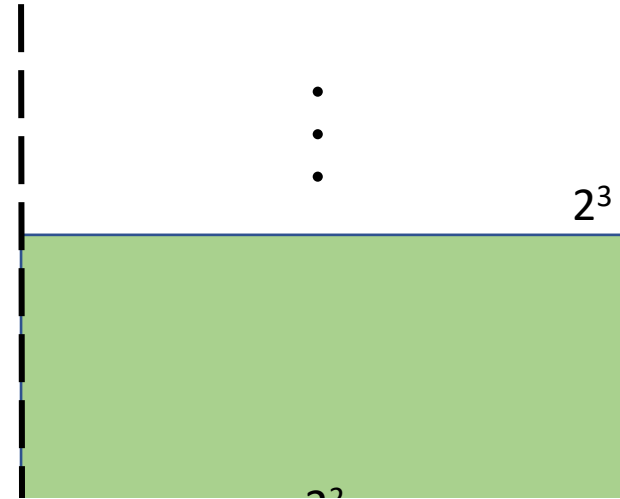
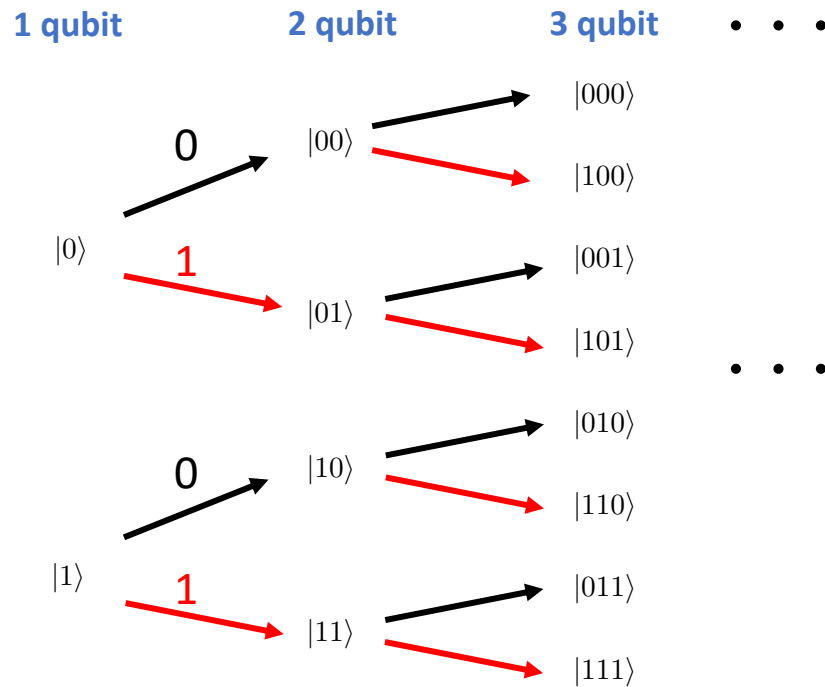


Hilbert Space dimension with qubits

Illustration of quantum advantages

Systems described on qubits

$$|\Psi\rangle = \sum_{s_i=0,1} \Psi_{s_1, \dots, s_N} |s_1, \dots, s_n\rangle$$



Quantum supremacy

$\sim 2^{50}$

With 2^{300} (i.e. 300 qubits) the size is more than the number of particles In the universe. J. Preskill

A more precise view of problem complexity and quantum advantage

Complexity of a problem

A fundamental question is how much “computational time” it takes to solve problem – this is linked to the complexity of the problem (Church-Turing thesis, Problem complexity classification, ...)

M.A. Nielsen and I. L. Chuang, “Quantum computation and quantum information”, Cambridge university press.

- ➔ What could be computed is independent on computers – only speed will change with hardware, resources progresses
- ➔ Quantitatively, we are interested in the “speedup” to perform a calculation

Speedup can be:

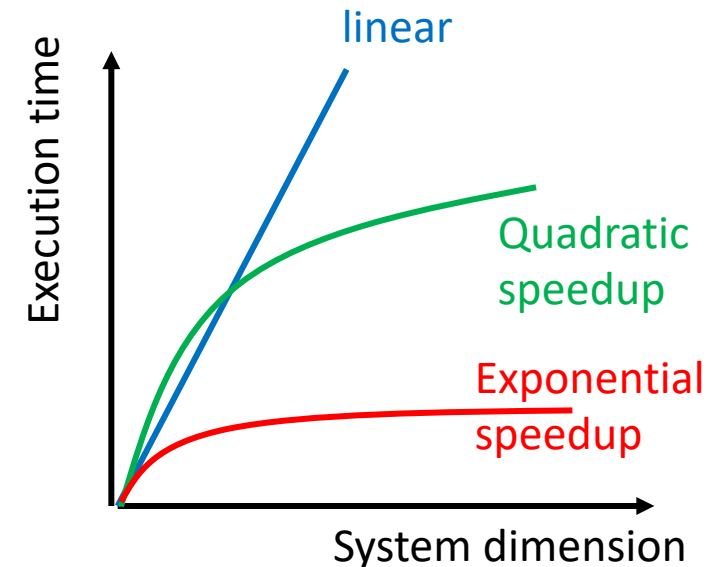
linear $t \rightarrow t/1000$

quadratic $t \rightarrow \sqrt{t} \rightarrow t^2 \rightarrow t$

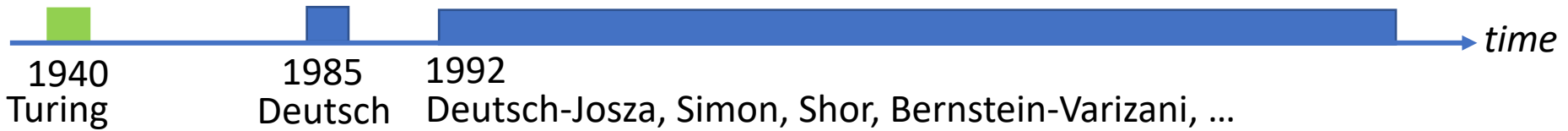
polynomial $t^x \rightarrow t$

But the speed up could not be exponential ! $2^t \rightarrow t$

Quantum computers contradict this thesis- some algorithms Promises exponential speedup for specific calculations



A calculation that's takes 1 year in “linear” case would takes 24 seconds with exponential speedups

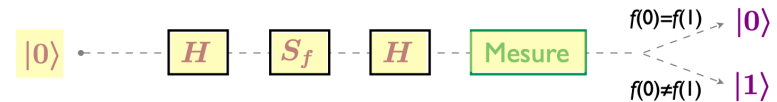
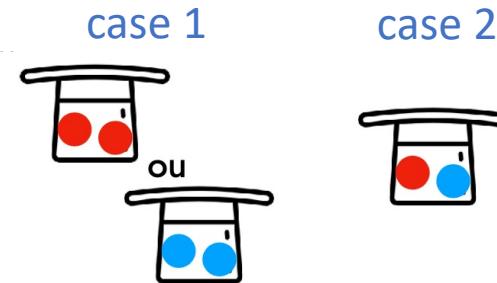


➔ Church Turing problem: analysis of a computational problem, complexity ...

➔ The “simple” Deutsch problem: $f : \{0,1\} \rightarrow \{0,1\}$ (Oracle)

Q: determine if $f(0)=f(1)$

- Classically requires to have 2 answers $f(0) ? f(1) ?$
- Quantum: one can directly ask $f(0)=f(1)$



➔ Deutsch-Josza problem (92)

$f : \{0,1\}^n \rightarrow \{0,1\}$ Q: Is f constant or an equal mixing of 0 and 1 ?

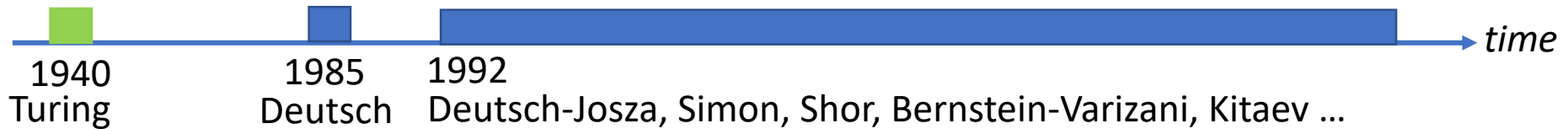
- Classically requires $1+2^{n-1}$ questions
- Quantum: requires only 1 question but n qubits

Example $n=4$

- $f(0000) \neq 0$
- $f(0001) \neq 0$
- $f(0010) \neq 0$
- $f(0011) \neq 0$
- $f(0100) \neq 0$
- $f(0101) \neq 0$
- $f(0110) \neq 0$
- $f(0111) \neq 0$
- $f(1000) \neq 0$

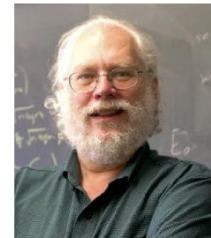
...

These problems are rather specific – usefulness?



Tracking applications and progress

- ➔ Simon (1994): find the period of a function $f: \{0,1\}^n \rightarrow \{0,1\}^n$
 - Classically requires exponential number of questions.
 - Quantum: linear number of questions
 - Ingredient: Quantum Fourier transform (QFT) in qubit representation
- ➔ Shor (1994): find the period of a function $f: \mathbb{Z} \rightarrow \mathbb{Z}_N$
 - Ingredient: QFT, quantum phase estimation (QPE), no Oracle
 - Applications: factorization, discrete log (quantum algo can potentially break the RSA public key encryption technique)
 - ➔ Birth of the post-quantum cryptography
- ➔ Kitaev 95: Circuits for QPE: eigenvalue problems, linear algebra, ...
- ➔ Grover search algorithm (1997): unsupervised search algorithm, data mining, ...



Peter Shor

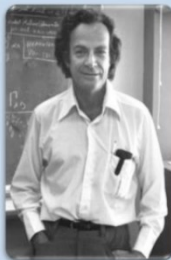
Why quantum computing is becoming mature now ?

Quantum technologies/devices

Simulating physics with computers-1982

Richard P. Feynman (Nobel Prize in Physics 1965)

"Nature isn't classical, dammit, and if you want to make a simulation of nature, you'd better make it quantum mechanical, and by golly it's a wonderful problem, because it doesn't look so easy."



Quantum Theory



1927

Quantum Computing

1982

7 qubits
Los Alamos

2000

12 qubits
MIT

2006

128 qubits
DWave

2011

50 qubits
IBM

2015

17 qubits
IBM

2017

128 qubits
Rigetti

2018

1152 qubits
DWave

2048 qubits
DWave

72 qubits
Google

Quantum Computing Cloud

General Quantum mechanics

1940
Turing

1985
Deutsch

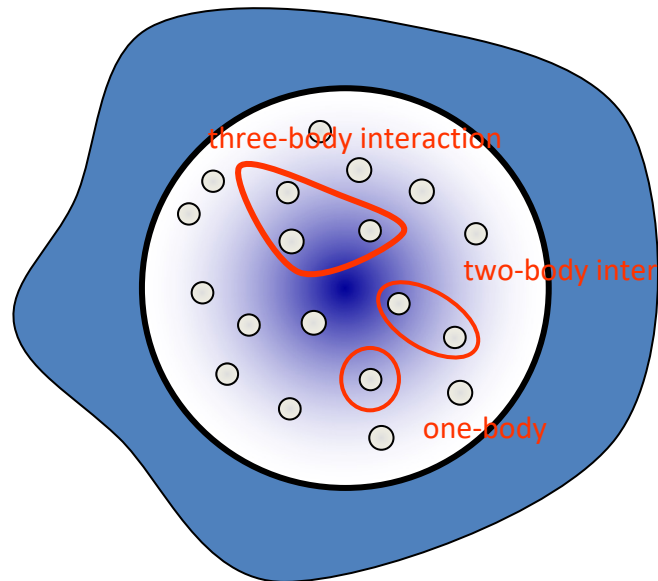
1992
Deutsch-Josza, Simon, Shor, Bernstein-Varizani, Kitaev ...

General Quantum algorithmic

Quantum computing is democratizing

time

One example: Simulation of complex quantum (interacting) systems

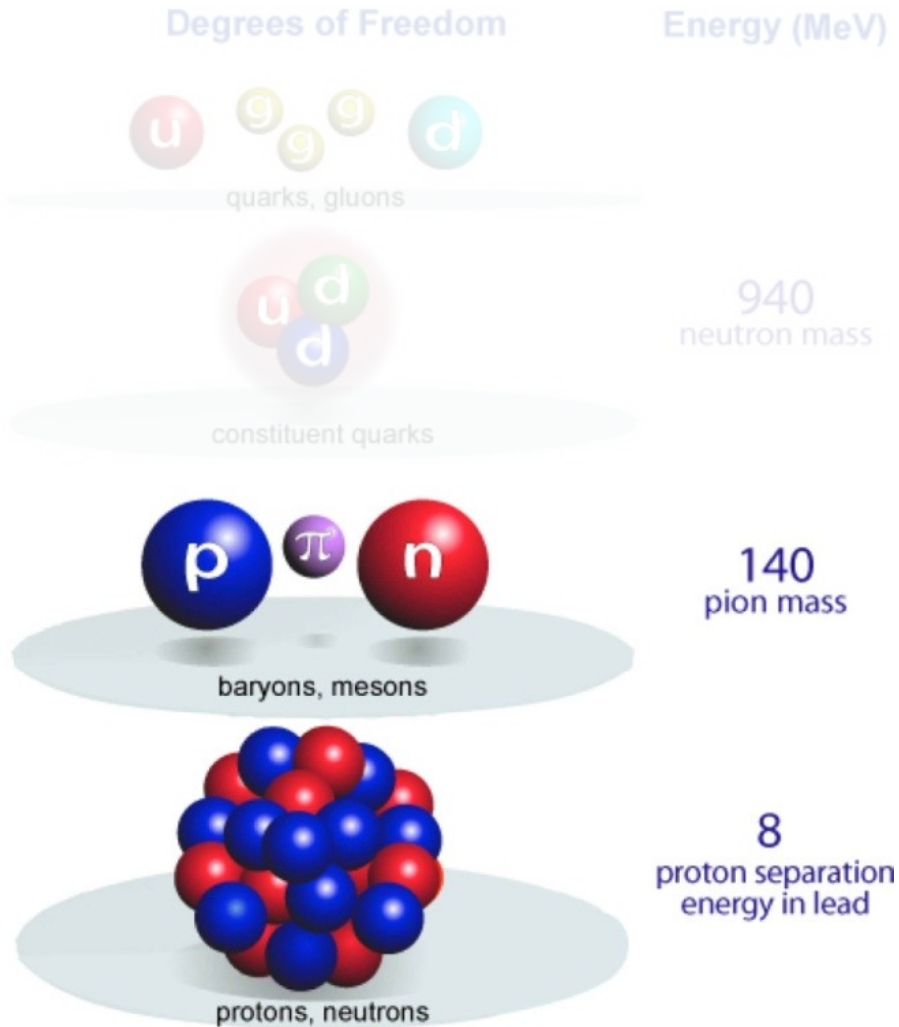


If you have N one-body degrees of freedoms
The Hilbert space has an exponential
Scaling ($\sim N!$)

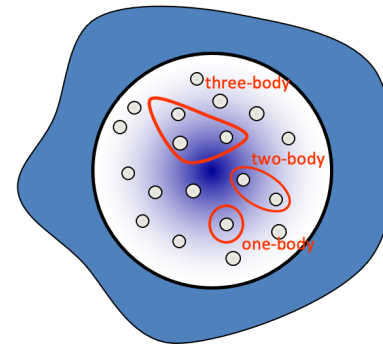


Even today, only a limited area (small systems- few %)
of the nuclear Chart can be calculated with most
powerful Supercomputers.

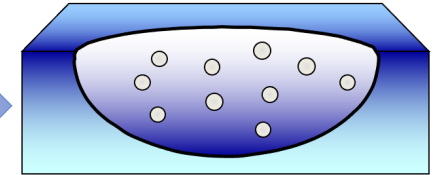
Some phenomenology



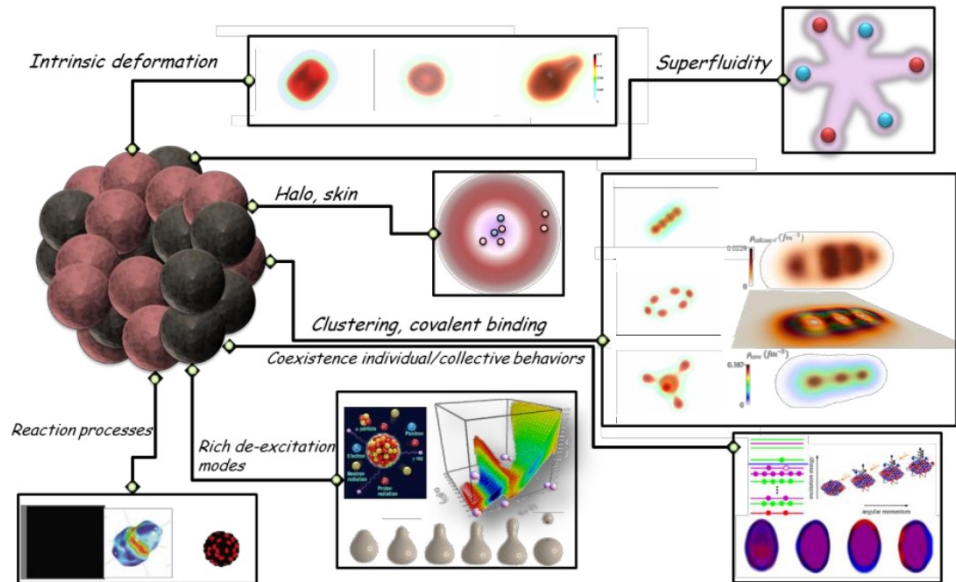
Strongly interacting fermions



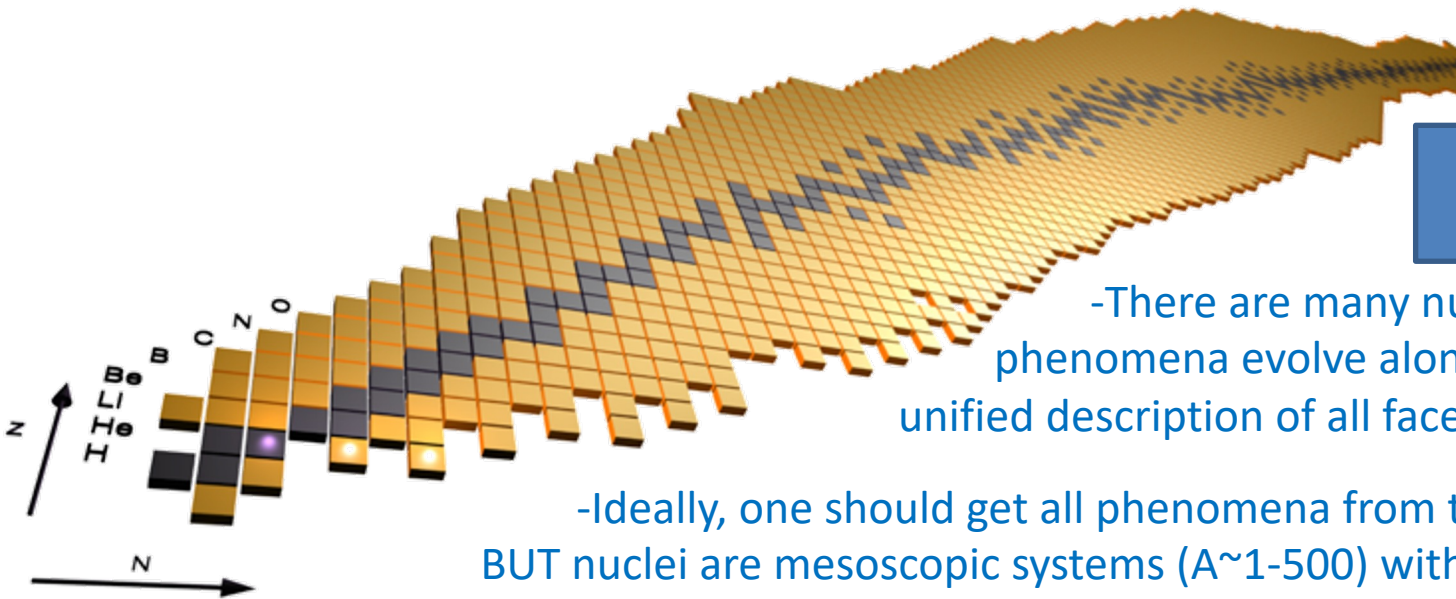
Quantum self-bound Fermi droplets



to atomic nuclei phenomena



Why quantum computing can help to describe nuclear physics ?

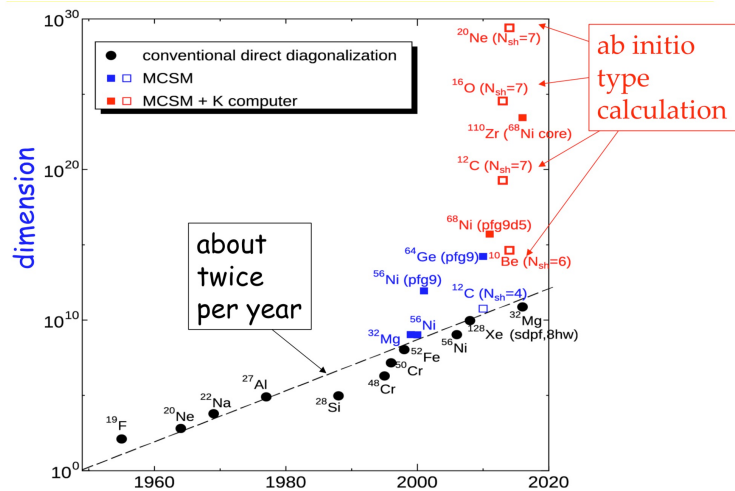


Some evident sources of complexity in nuclei

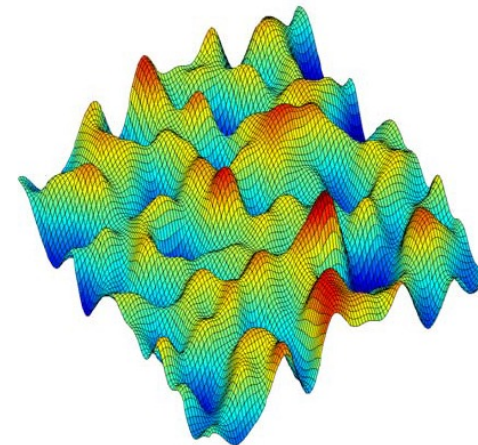
-There are many nuclei (>3000). Nuclear phenomena evolve along the nuclear chart. A unified description of all facets would be desirable.

-Ideally, one should get all phenomena from the bare interaction BUT nuclei are mesoscopic systems ($A \sim 1-500$) with bad numerical scaling.

-Each nucleus is a complex problem per se.

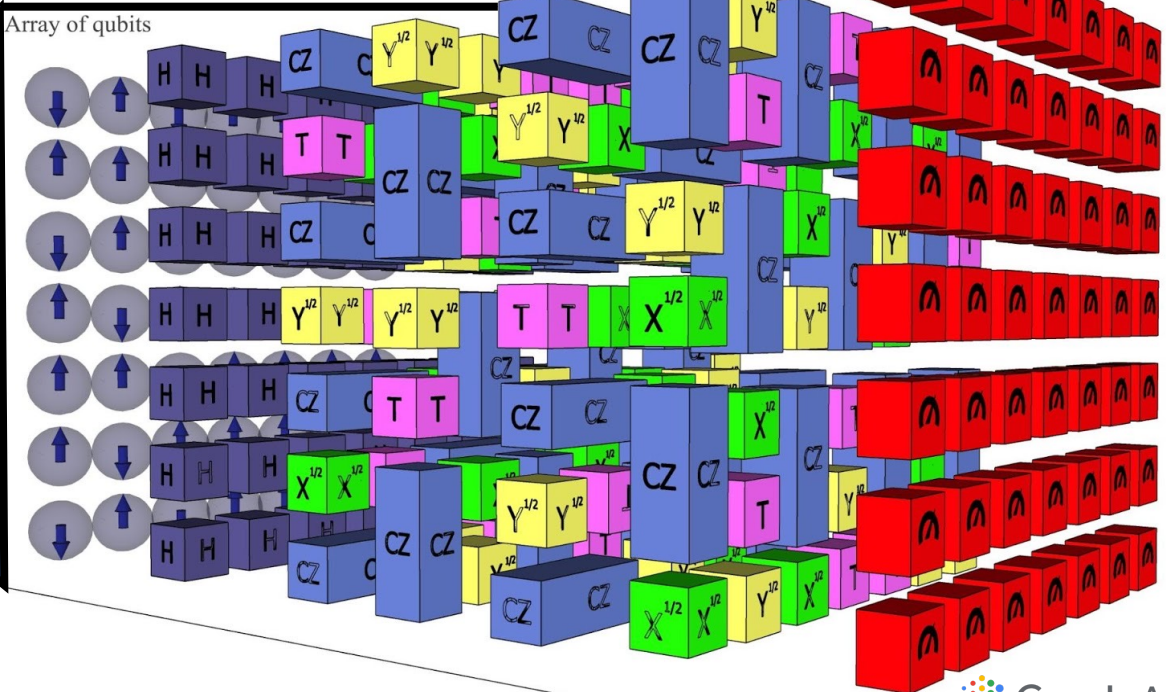
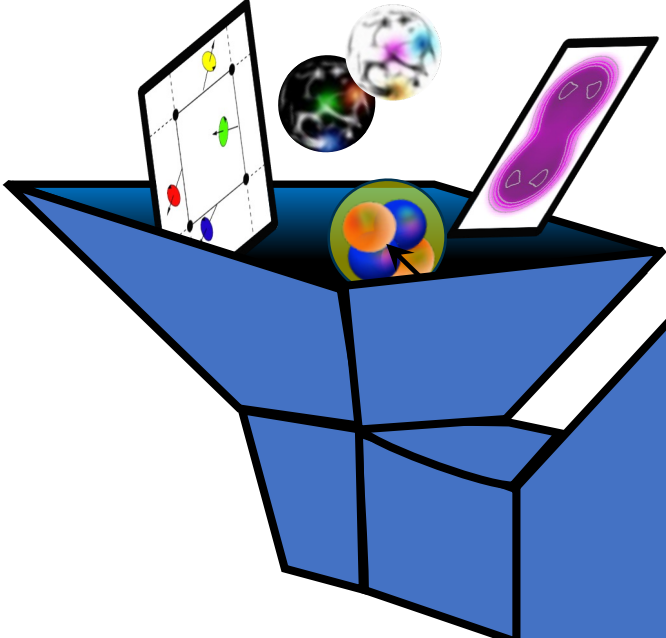


(from T. Otsuka)



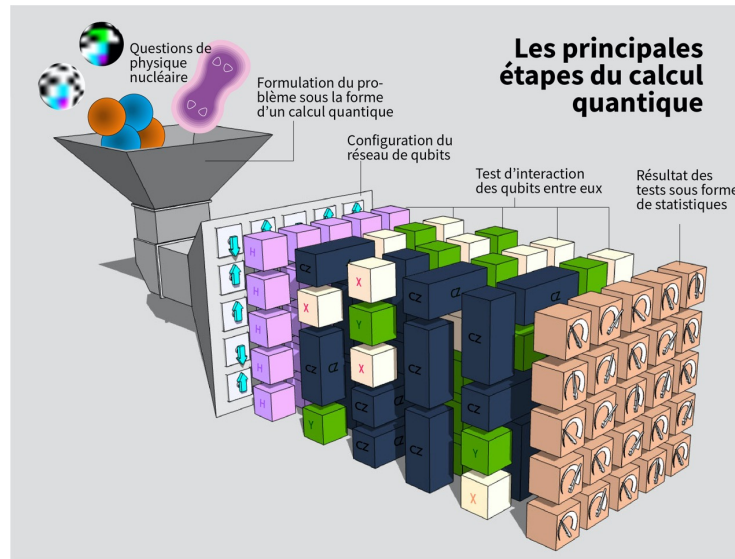
(Energy landscape of a molecule)

➔ This motivates the search of disruptive techniques (high risk/high potential benefits)



(Second part)

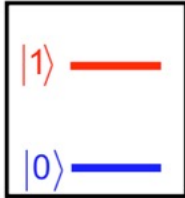
“PLAYING” WITH QUBITS



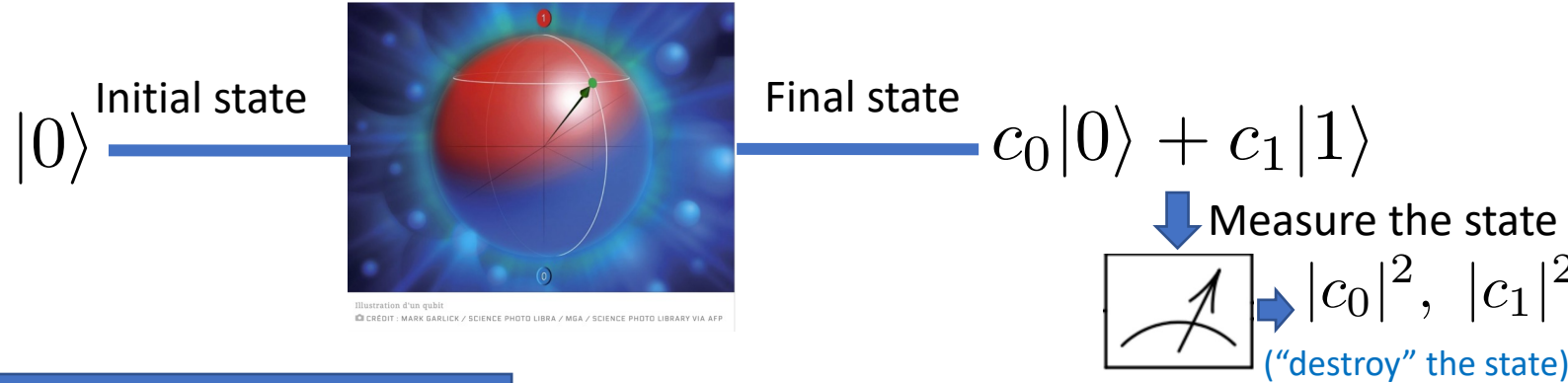
Introduction to digital quantum computing

qubit

2 level system

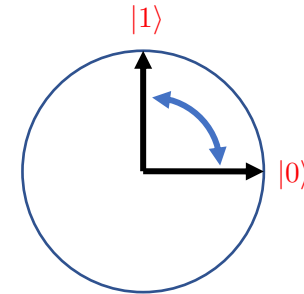
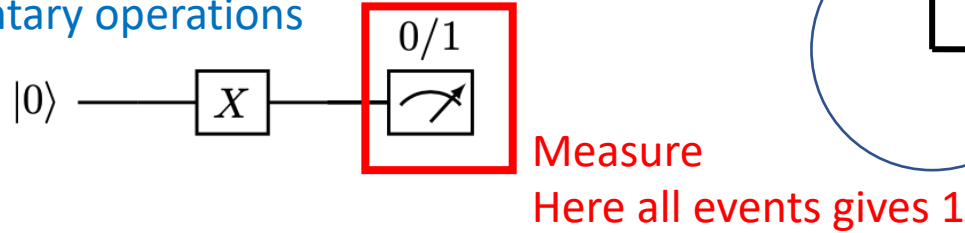
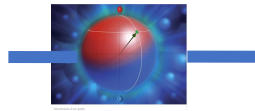


Manipulate the Qubits (Make rotations)



Example of simple unary operations (gates)

Elementary operations



Pauli matrices

Pauli-X (X) $\square \text{X}$ \oplus

Pauli-Y (Y) $\square \text{Y}$

Pauli-Z (Z) $\square \text{Z}$

$\begin{bmatrix} |0\rangle & |1\rangle \\ \langle 0| & \langle 1| \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \rightarrow$ NOT operation

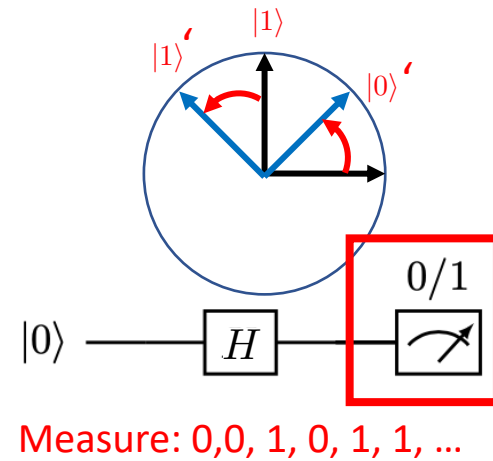
Other useful unary operations

$\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$ Hadamard (H) $\square \text{H}$ $\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & -1 \\ 1 & 1 \end{bmatrix}$

Phase (S, P) $\square \text{S}$ $\begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$

$\pi/8$ (T) $\square \text{T}$ $\begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{bmatrix}$

Hadamard



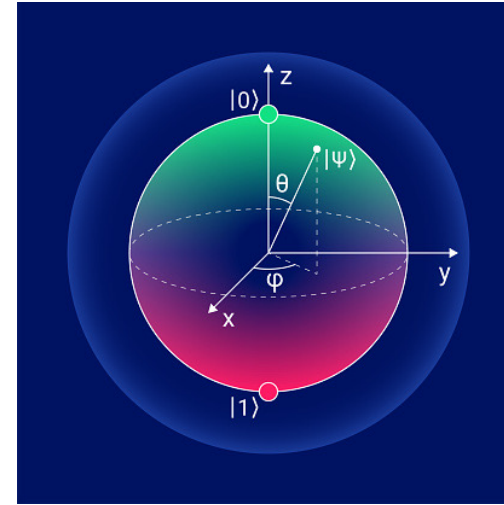
Another useful representation of operation (more familiar to physicists)

The Bloch sphere

From qubits to spins

$$\begin{aligned}\sigma_x &= X = |1\rangle\langle 0| + |0\rangle\langle 1| \\ \sigma_y &= Y = i|1\rangle\langle 0| - i|0\rangle\langle 1| \\ \sigma_z &= Z = |0\rangle\langle 0| - |1\rangle\langle 1|\end{aligned}$$

➔ Obey standard spin algebra



Rotations corresponds to gates

$$\text{---} \boxed{R_X(\varphi) = e^{-i\varphi X/2}} \text{---}$$

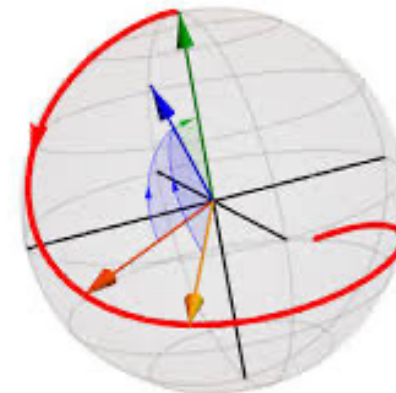
Any state can be written as

$$|\Psi\rangle = \cos(\theta)|0\rangle + e^{i\varphi} \sin(\theta)|1\rangle$$

Such state can always be obtained from 3 rotations

$$|\Psi\rangle = R_Y(\theta_1)R_Z(\theta_2)R_Y(\theta_3)|0\rangle$$

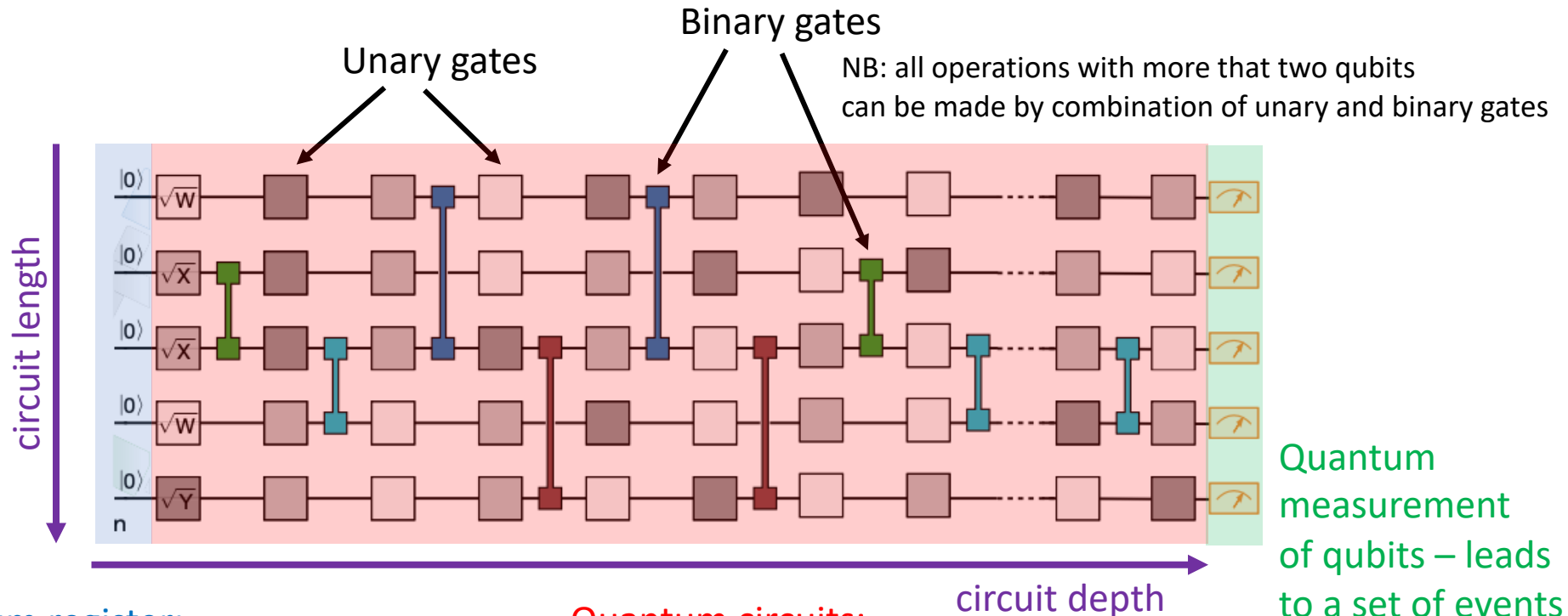
Unary operation can be represented as a path in the Bloch sphere



Quantum computing with more than one qubit

Some terminology – general aspects

Quantum circuits with more than one qubit



Quantum register:
Define the qubit computational

basis: $|0, 0, \dots, 0, 0\rangle$
 $|0, 0, \dots, 1, 1\rangle$
 $|0, 0, \dots, 0, 1\rangle$
 \dots

Hilbert space size 2^n

Quantum circuits:

Constraint: the circuit makes
Unitary transformation, i.e. no
loss of information

Advantage: one can imagine
to do e^{itH}

$$\sum_{i_k=0,1} a_{i_1 i_2 i_3 \dots i_{2N}} |i_1, i_2, i_3 \dots i_{2N}\rangle$$



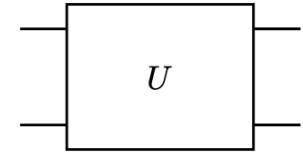
Gives the $|a|^2$

Quantum computing with more than one qubit

Some terminology – general aspects

Binary gates

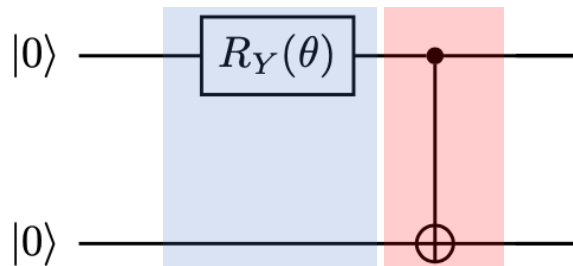
Binary gates are operations acting simultaneously on two qubits and induced entanglement



Remember the example:

$$|\Phi\rangle = \alpha|\😊😊\rangle + \beta|\😞😞\rangle \text{ [equivalent to Bell states]}$$

Quantum circuit



Rotation Control-X gate
Also known as CNOT

Step 1

$$|00\rangle \rightarrow \cos(\theta/2)|00\rangle + \sin(\theta/2)|01\rangle$$

Step 2

$$\rightarrow \cos(\theta/2)|00\rangle + \sin(\theta/2)|11\rangle$$



Indirect measurement and Born rule for measurement

If I measure 0 in first qubit ➡ After the measurement the state collapse to $|00\rangle$

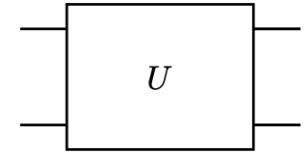
If I measure 1 in first qubit ➡ After the measurement the state collapse to $|11\rangle$

Quantum computing with more than one qubit

Some terminology – general aspects

Binary gates

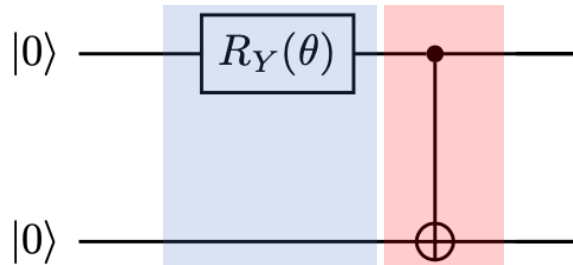
Binary gates are operations acting simultaneously on two qubits and induced entanglement



Remember the example:

$$|\Phi\rangle = \alpha |😊😊\rangle + \beta |😞😞\rangle \text{ [equivalent to Bell states]}$$

Quantum circuit



Rotation

Control-X gate

Also known as CNOT

Step 1

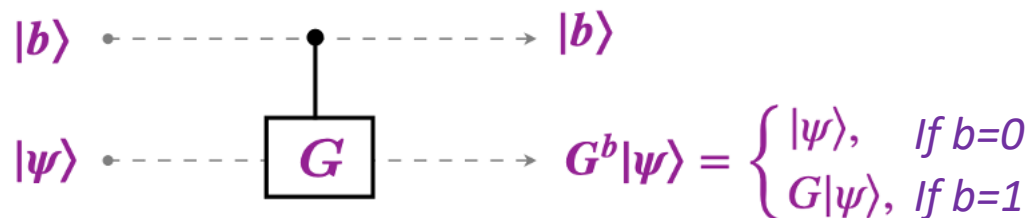
$$|00\rangle \rightarrow \cos(\theta/2)|00\rangle + \sin(\theta/2)|01\rangle$$

Step 2

$$\rightarrow \cos(\theta/2)|00\rangle + \sin(\theta/2)|11\rangle$$



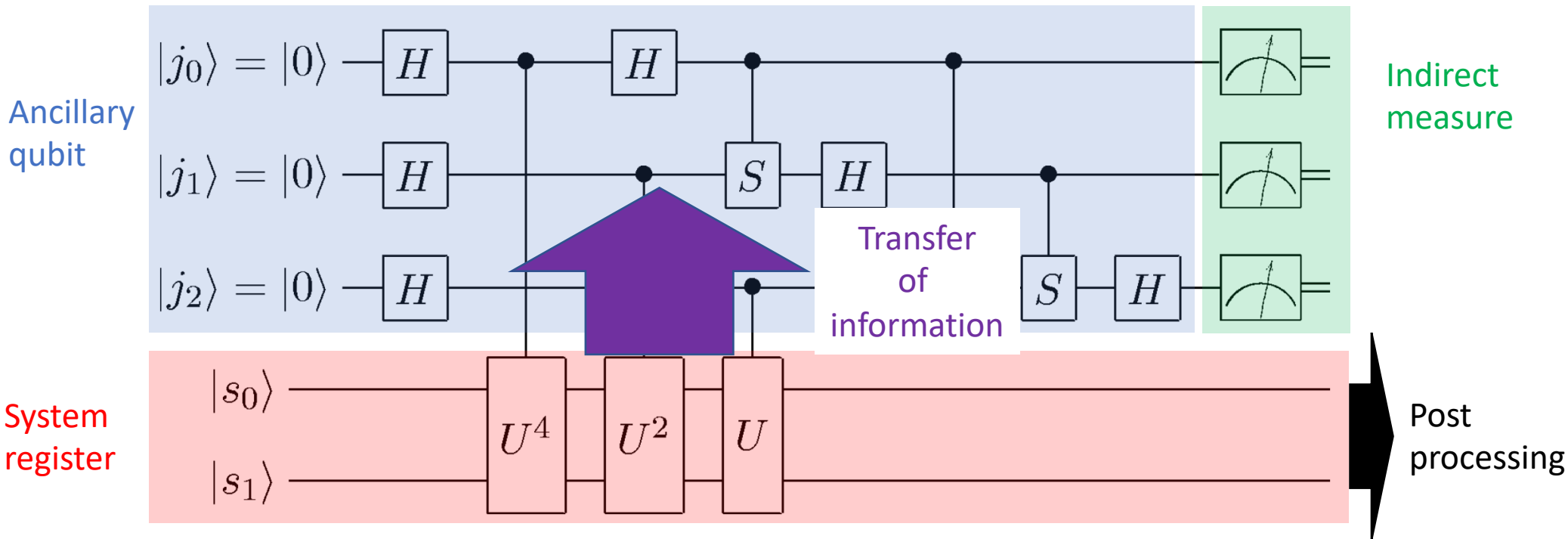
Controlled gates (IF-THEN-ELSE)



Quantum computing with more than one qubit

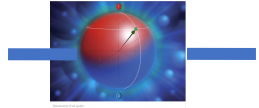
Some terminology – general aspects

A more complete view of quantum computing circuits

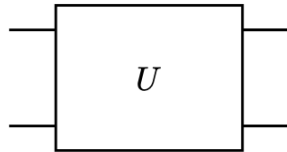


The quantum computing toolkit

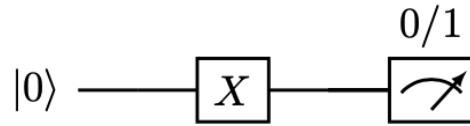
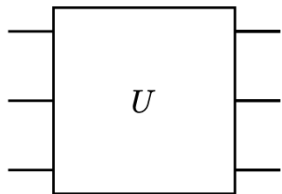
Unary operations



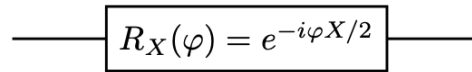
Binary operations



Ternary operations



Rotations



Standard examples

Controlled Not
(CNOT, CX)



$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

$$|11\rangle \leftrightarrow |10\rangle$$

Controlled Z (CZ)



$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix}$$

$$|11\rangle \leftrightarrow -|11\rangle$$

SWAP



$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{matrix} |00\rangle \\ |01\rangle \\ |10\rangle \\ |11\rangle \end{matrix}$$

$$|01\rangle \leftrightarrow |10\rangle$$

Standard example

Toffoli
(CCNOT,
CCX, TOFF)



$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

$$|110\rangle \leftrightarrow |111\rangle$$

Standard examples

Pauli-X (X)



$$\begin{matrix} |0\rangle & |1\rangle \\ \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} & \begin{bmatrix} 0 \\ 1 \end{bmatrix} \end{matrix}$$

Pauli-Y (Y)



$$\begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}$$

Pauli-Z (Z)



$$\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$$

Hadamard (H)

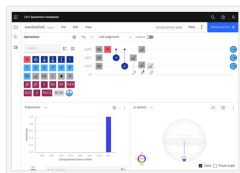


$$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

Real quantum computers. Right at your fingertips.

IBM offers cloud access to the most advanced quantum computers available.
Learn, develop, and run programs with our quantum applications and systems.

If you want to play with qubits Try the “IBM Quantum experience” online



Graphically build quantum circuits

Start building quantum circuits right away with IBM Quantum Composer. No sign in required.

[Explore Quantum Composer](#)

The screenshot displays the IBM Quantum Composer interface. The main workspace shows a quantum circuit with four qubits (q[0], q[1], q[2], q[3]) and a classical register (c4). The circuit includes a Z gate on q[0], an RX gate on q[2], and Hadamard (H) gates on q[0] and q[3]. CNOT gates connect q[0] to q[1], q[1] to q[2], and q[2] to q[3]. The interface includes a menu bar (File, Edit, View), an Operations panel with a search bar and a grid of quantum gates, and a visualization area at the bottom. The visualization area contains a bar chart for 'Probabilities' and a 'Q-sphere' visualization.

Probabilities

Computational basis states	Probability (%)
0000	25
0001	0
0010	25
0011	0
0100	0
0101	0
0110	0
0111	0
1000	25
1001	0
1010	25
1011	0
1100	0
1101	0
1110	0
1111	0

Q-sphere

The Q-sphere visualization shows the state of the four qubits. The state is a superposition of $|0000\rangle$, $|0010\rangle$, $|1000\rangle$, and $|1010\rangle$. The phase angle is 0 .

State Phase angle

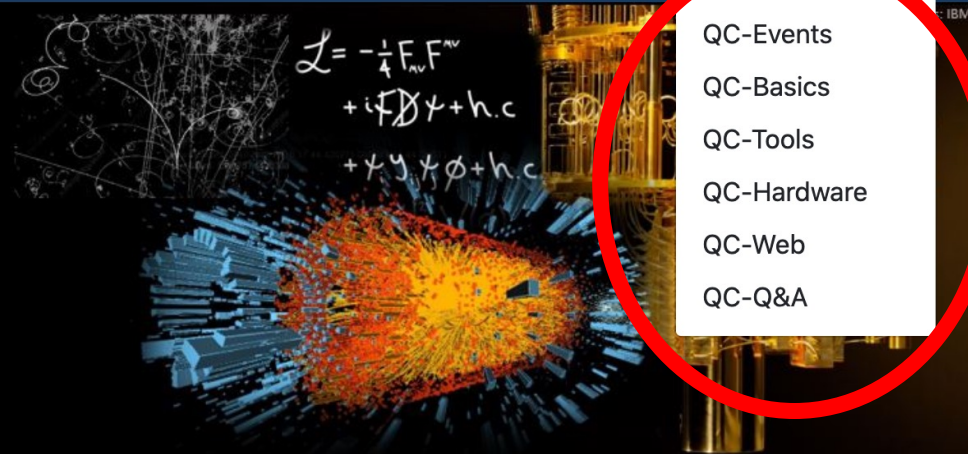
If you want to start QC programming

<https://qc.pages.in2p3.fr/web/>

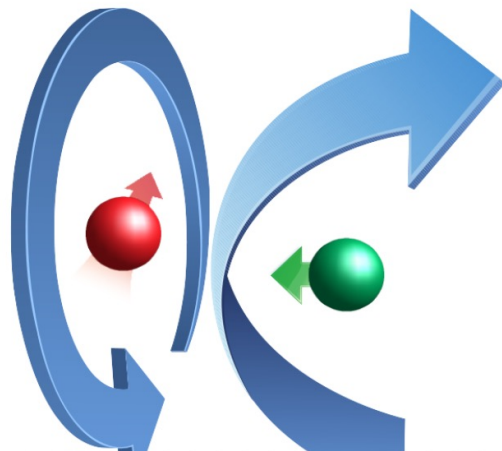
QC2I

About ▾ Members Meetings Talks Collab-Tools ▾ QC-Events

QC2I: Quantum Computing for the two Infinities



- QC-Events
- QC-Basics
- QC-Tools
- QC-Hardware
- QC-Web
- QC-Q&A

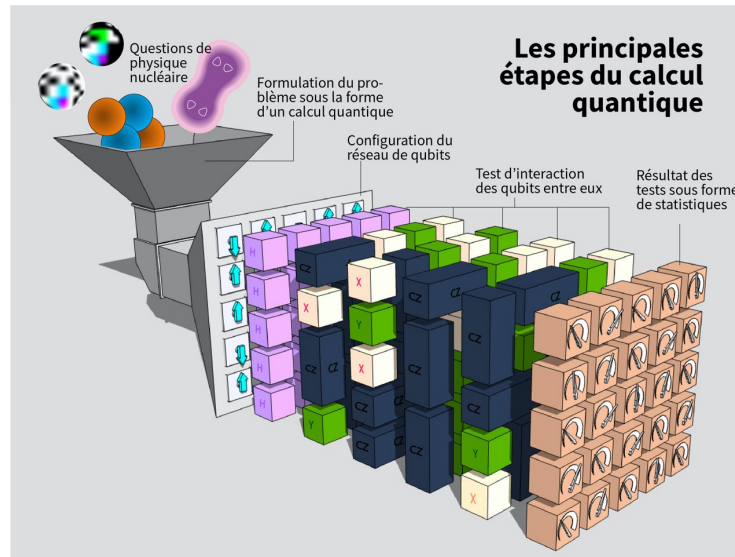


QC2I is a computing project supported by [IN2P3](#), the French national nuclear and particle physics institute. Its goal is to explore the possible applications of the emerging quantum computing technologies to particles and nuclear physics problems as well as astrophysics. The main tasks are:

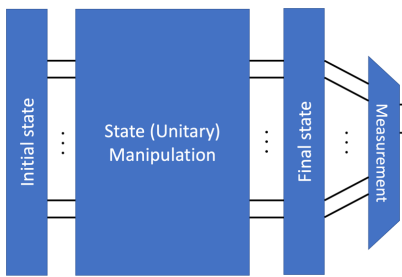
- to identify, within IN2P3, scientists/engineers/technicians who are interested in using quantum technologies,
- to facilitate the access and training on quantum computers,
- to identify milestones applications for nuclear/particle physics and astrophysics,
- to design dedicated algorithms and proof of principle applications.

The project action has three main directions: **Prepare the Quantum Computing Revolution (PQCR)**, **Quantum Machine Learning (QML)**, **Complex Quantum Systems Simulation**

Some illustrations of applications



General strategy for quantum computing



Take a problem
(classical or quantum)



Encode the problem
Onto a set of qubits



Make Quantum Computing
Program to solve the problem



Test on a
QC emulator



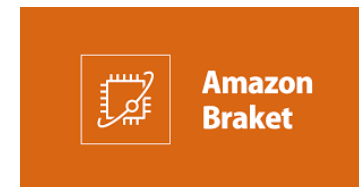
Test on a true
QC

There are many ways to encode a classical problem (phase, amplitude, time, ...)

Schuld and Petruccione,
Supervised learning with
quantum computers (2018)

For instance
Qiskit, cirq,
myQLM, ...

IBM Quantum





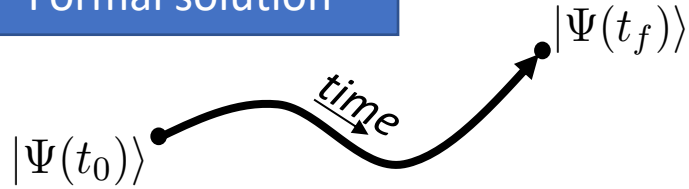
First Example:
Solution of a 1D
Time-Dependent Schroedinger Eq.

Example 1: solving a 1D Schrödinger Equation on a quantum computer

I want to solve $i\hbar \frac{d}{dt} \Psi(x) = \left\{ \frac{-\hbar^2}{2m} \Delta + V(x) \right\} \Psi(x)$ \rightarrow Discretization on a mesh

Formal solution

Formal solution



$$|\Psi(t)\rangle = e^{\frac{1}{i\hbar}(t-t_0)H} |\Psi(t_0)\rangle$$

H is usually a big matrix

Practical solution

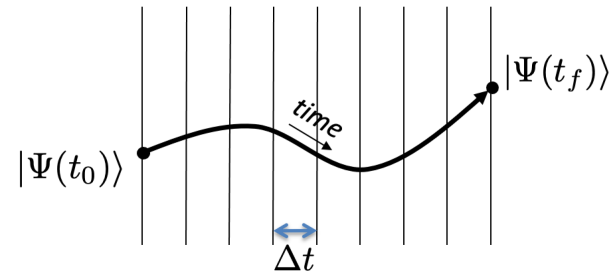
$$i\hbar \dot{\mathbf{F}}(t) = \mathbf{H} \times \mathbf{F}(t)$$



$$\mathbf{F}(t + \Delta t) = \exp\left(\frac{\Delta t}{i\hbar} \mathbf{H}\right) \times \mathbf{F}(t)$$

Time discretization

time: $\{t_i\}$ time-step: Δt



Direct

$$\exp\left(-\frac{\Delta t}{i\hbar} \mathbf{H}\right) \simeq 1 - \frac{\Delta t}{i\hbar} \mathbf{H} + \frac{1}{2!} \left(\frac{\Delta t}{i\hbar} \mathbf{H}\right)^2 + \dots$$

$(\Delta t)^n$, non-unitary, any dim.

Crank-Nicholson

$$\mathbf{F}(t + \Delta t) = \frac{1 - \frac{\Delta t}{2i\hbar} \mathbf{H}}{1 + \frac{\Delta t}{2i\hbar} \mathbf{H}} \mathbf{F}(t)$$

$(\Delta t)^2$, unitary, 1D only

Split-Operator

$$\mathbf{F}(t + \Delta t) \simeq e^{-i\Delta t \frac{\mathbf{P}^2}{4\hbar m}} e^{-\frac{i}{\hbar} \Delta t \mathbf{V}} e^{-i\Delta t \frac{\mathbf{P}^2}{4\hbar m}} \times \mathbf{F}(t)$$

$(\Delta t)^2$, unitary, any dim.

Non unitary

Requires to inverse a matrix

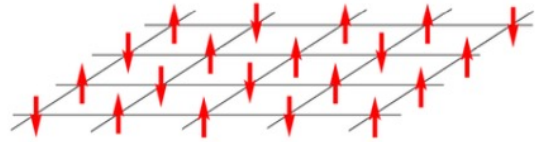
Might be used

Solving a 1D Schrödinger Equation on a *quantum computer*

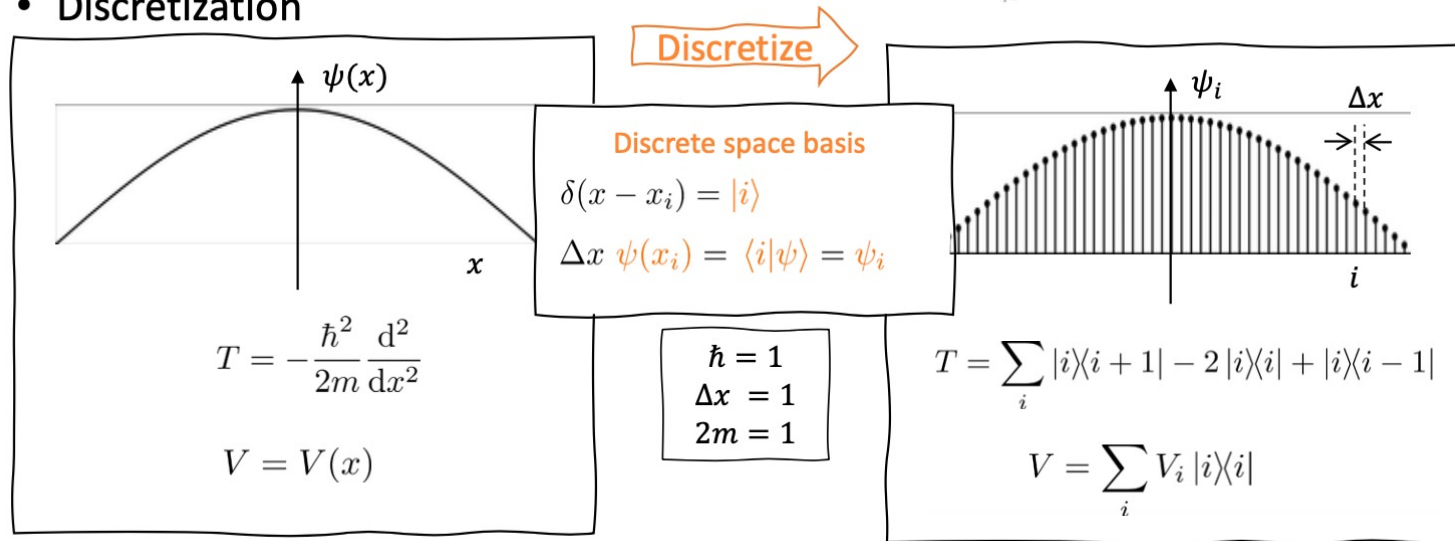
I want to solve $i\hbar \frac{d}{dt} \Psi(x) = \left\{ \frac{-\hbar^2}{2m} \Delta + V(x) \right\} \Psi(x)$ \rightarrow Discretization on a mesh



Jing Zhang
PhD IJCLab



• Discretization



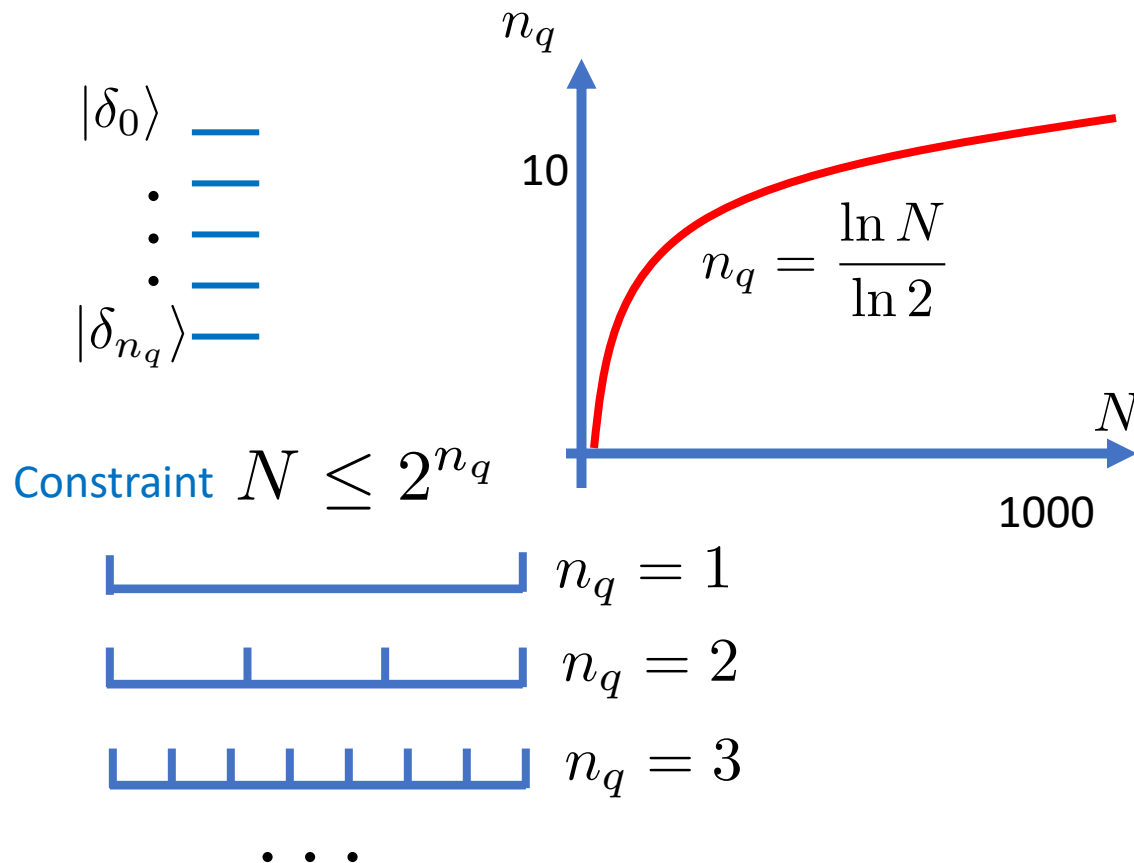
$$\begin{aligned}
 \text{e.g., } \langle \psi | T | \psi \rangle &= -\frac{\hbar^2}{2m} \int \psi^* \frac{d^2}{dx^2} \psi dx \\
 &\approx -\frac{\hbar^2}{2m} \sum_i \psi_i^* \frac{\psi_{i+1} - 2\psi_i + \psi_{i-1}}{(\Delta x)^2} \\
 &= \sum_i \langle \psi | i \rangle \frac{\langle i+1 | \psi \rangle - 2 \langle i | \psi \rangle + \langle i-1 | \psi \rangle}{(\Delta x)^2}
 \end{aligned}$$

Solving a 1D Schrödinger Equation on a *quantum computer*

Encoding into qubits

We have a discrete basis $|i = 0 \dots N - 1\rangle$

Natural encoding – Use the binary representation $|i\rangle = |\delta_{n_q} \dots \delta_0\rangle$ with $i = \sum_k \delta_k 2^k$



$ i\rangle$	Standard Binary (SB) Encoding	Gray code (GC) Encoding
$ 0\rangle$	$ 000\rangle$	$ 000\rangle$
$ 1\rangle$	$ 001\rangle$	$ 001\rangle$
$ 2\rangle$	$ 010\rangle$	$ 011\rangle$
$ 3\rangle$	$ 011\rangle$	$ 010\rangle$
$ 4\rangle$	$ 100\rangle$	$ 110\rangle$
$ 5\rangle$	$ 101\rangle$	$ 111\rangle$
$ 6\rangle$	$ 110\rangle$	$ 101\rangle$
$ 7\rangle$	$ 111\rangle$	$ 100\rangle$

Encoding into qubits with
Gray code

Gray code encoding of Hamiltonian

$$H = \sum_k (|k\rangle\langle k+1| - 2|k\rangle\langle k| + |k\rangle\langle k-1|)$$

➔ Goal : transform the Hamiltonian into gates products of (X, Y, Z)

How this works: $X = |1\rangle\langle 0| + |0\rangle\langle 1|$ We add the identity $I = |0\rangle\langle 0| + |1\rangle\langle 1|$
 $Y = i|1\rangle\langle 0| - i|0\rangle\langle 1|$
 $Z = |0\rangle\langle 0| - |1\rangle\langle 1|$ ➔ $\begin{cases} |1\rangle\langle 1| \rightarrow \frac{1}{2}(I - Z) \\ |0\rangle\langle 0| \rightarrow \frac{1}{2}(I + Z) \end{cases}$

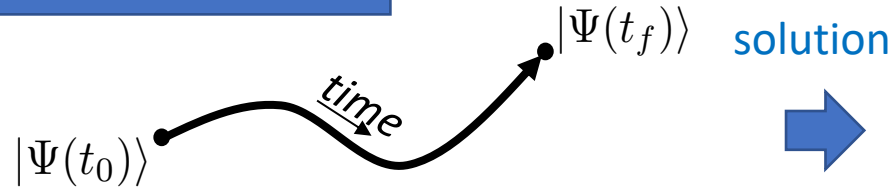
e.g., $|2\rangle\langle 3| + |3\rangle\langle 2| = |011\rangle\langle 010| + |010\rangle\langle 011|$
 $= |0\rangle\langle 0| \otimes |1\rangle\langle 1| \otimes (|1\rangle\langle 0| + |0\rangle\langle 1|)$
 $= \frac{1}{2}(I + Z_0) \otimes \frac{1}{2}(I - Z_1) \otimes X_2$

This can be made systematically

$$H = \text{fct}(X_\alpha, Y_\alpha, Z_\alpha)$$

Solving a 1D Schrödinger Equation on a *quantum computer*

Time-evolution



This can be a priori computed
Directly on a QC

$$|\Psi(t)\rangle = e^{\frac{1}{i\hbar}(t-t_0)H} |\Psi(t_0)\rangle$$

1. Decomposition of H into elementary blocks

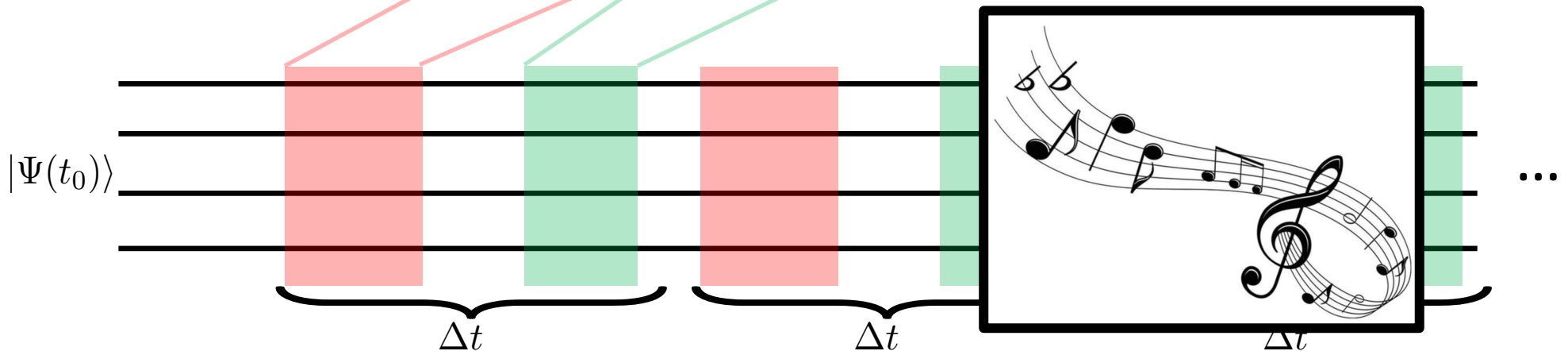
$$H = \sum_l H_l$$

2. Use a transformation (Trotter-Suzuki)

$$e^{-ix(A+B)} = \left(e^{-iAx/N} e^{-iBx/N} \right)^N + \mathcal{O}(t^2/N)$$

Example: $e^{i\Delta t H_1/\hbar} = e^{-i\Delta t H_1/\hbar} e^{-i\Delta t H_2/\hbar}$

3. Transforms to circuit



Solving a 1D Schrödinger Equation on a quantum computer

Deeper in technical details

The case of free propagation: $H = \frac{p^2}{2m}$

position $\xrightarrow{\text{Fourier transform}}$ momentum

Hamiltonian decomposition

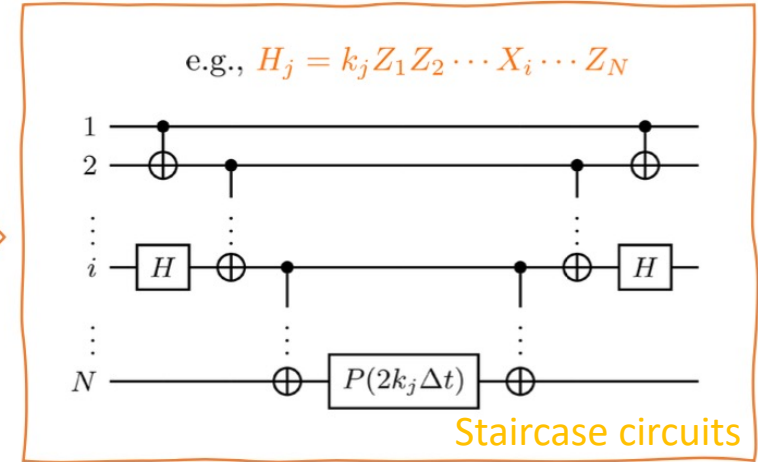
Time evolution:

$$\exp(-i\hat{H}t) = \exp(-i\sum_j \hat{H}_j t)$$

$$\approx \left[\prod_j \exp(-i\hat{H}_j \Delta t) \right]^{t/\Delta t}$$

Trotter decomposition

$$e^{i(A+B)t} = \lim_{\Delta t \rightarrow 0} (e^{iA\Delta t} e^{iB\Delta t})^{t/\Delta t}$$



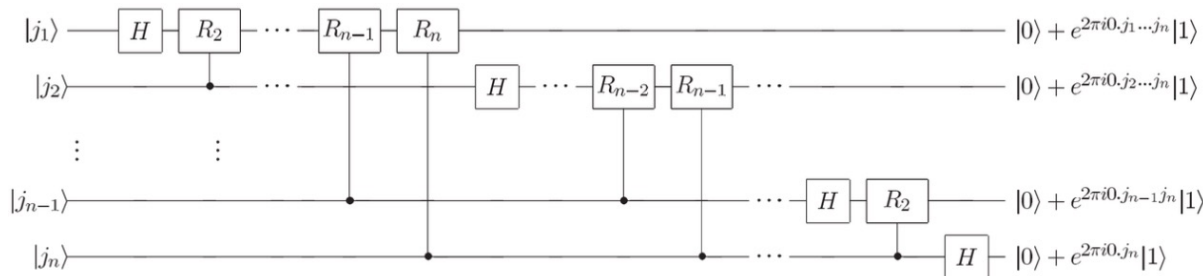
Digital or Quantum Fourier transformation

$$QFT: |j\rangle \rightarrow \frac{1}{2^{N/2}} \sum_{k=0}^{2^N-1} e^{2\pi i j (k-2^{N-1})/2^N} |k\rangle$$

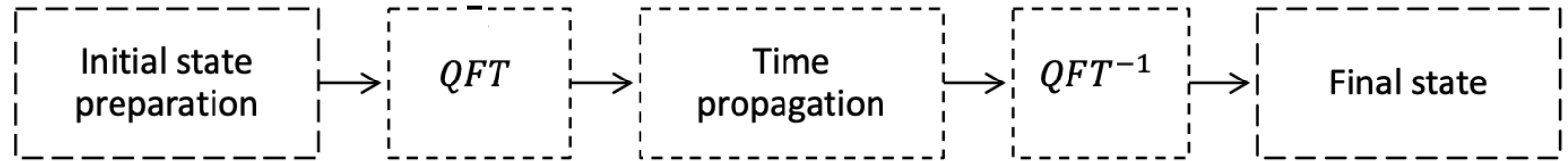
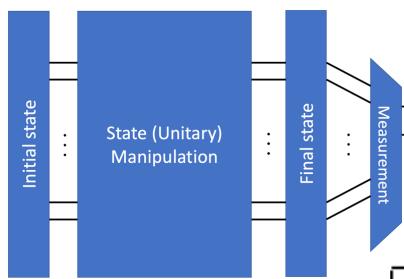
$$= \frac{e^{-\pi i j}}{2^{N/2}} (|0\rangle + e^{2\pi i 0 \cdot j_N} |1\rangle) \otimes (|0\rangle + e^{2\pi i 0 \cdot j_{N-1} j_N} |1\rangle)$$

$$\otimes \dots \otimes (|0\rangle + e^{2\pi i 0 \cdot j_1 j_2 \dots j_N} |1\rangle).$$

QFT circuit

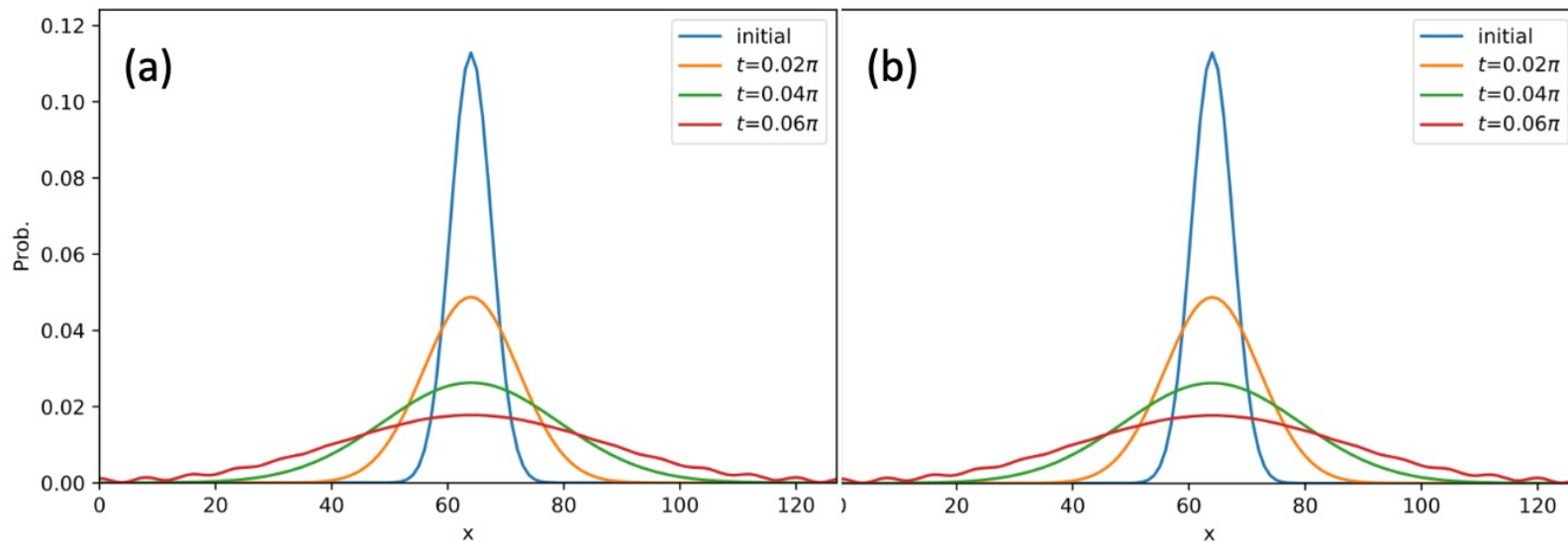


Solving a 1D Schrödinger Equation on a *quantum computer*



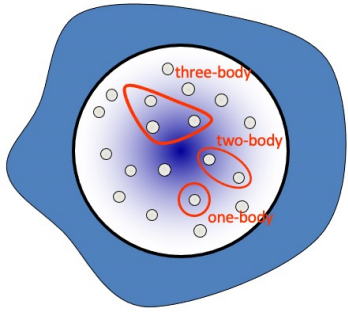
Simulation on classical computer

Simulation on quantum computer (emulator)

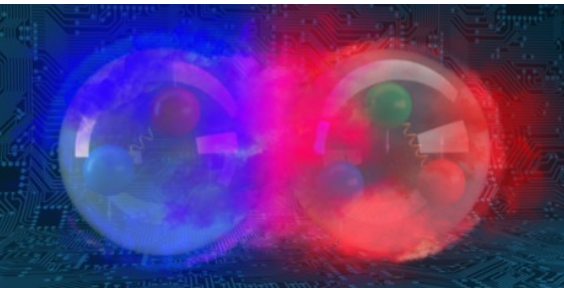


Some shortcomings:

- The number of qubits to get sufficient accuracy is already quite large
- Absorbing waves at the boundary?



Second Example: The deuteron problem in a schematic model



PHYSICAL REVIEW LETTERS **120**, 210501 (2018)

Editors' Suggestion

Featured in Physics

Cloud Quantum Computing of an Atomic Nucleus

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R. C. Pooser,^{1,4} D. J. Dean,³ and P. Lougovski^{1,†}

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²Computer Science and Mathematics Division, Oak Ridge National Laboratory, Oak Ridge, Tennessee 37831, USA

³Physics Division, Oak Ridge National Laboratory, Oak Ridge, Tennessee 37831, USA

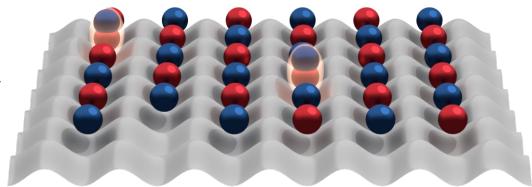
⁴Department of Physics and Astronomy, University of Tennessee, Knoxville, Tennessee 37996, USA

⁵National Center for Computational Sciences, Oak Ridge National Laboratory, Oak Ridge, Tennessee 37831, USA

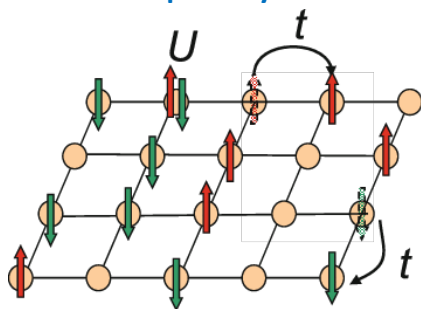
Solution of simple many-body problem with QC: a brief history

Jordan-Wigner (1928)

Fermions on [1D] lattice



Spin systems



Quantum chemistry
Condensed matter

1997- Abrams and Lloyd

A quantum algorithm
For eigenvalue problems

2001- Bravyi-Kitaev
Mapping fermions-Qubits

-2011-2012-Whitfield et al
-Seeley et al
The H_2 Hamiltonian

-2014 – Peruzzo et al,
The VQE algorithm
For classical-quantum calc.

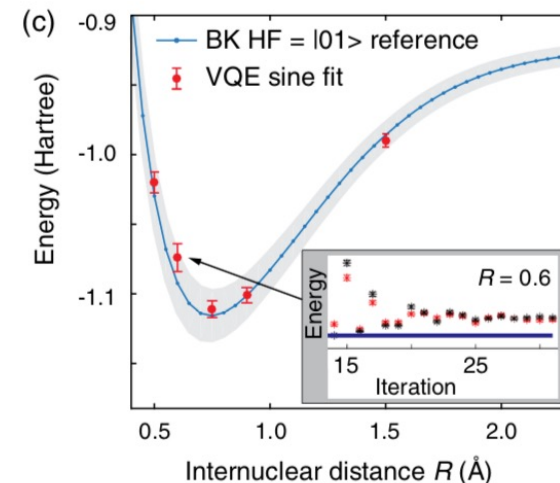
-2016 - O'Malley et al
First “real” calculations H_2

-2017 - Kandala et al
Calculations for H_2 , LiH, HBeH

-2018 - Hempel et al

...

Nuclear Physics



Creation of a US coll. To
Prepare QC

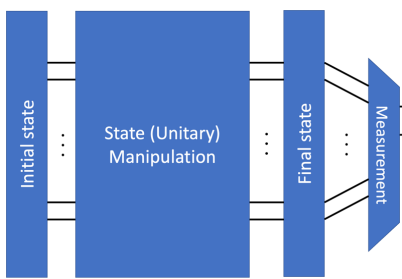
-2018 - Dimitrescu et al

First “real” calculations
For deuteron

-2019 - Lu et al

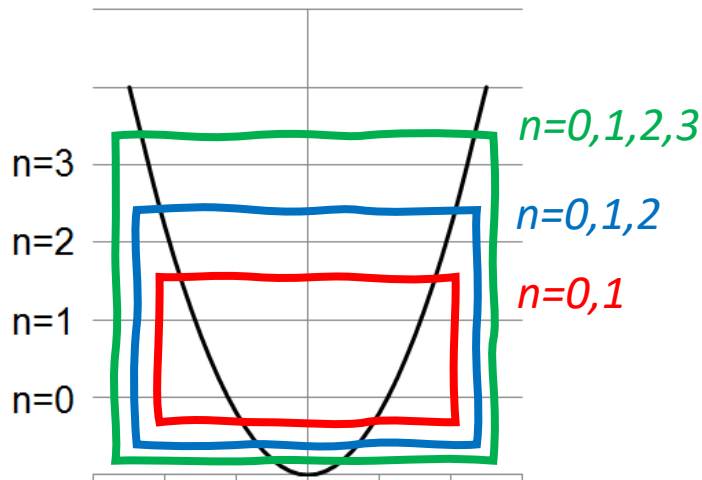
3H , 3He , α

Solving the deuteron problem on a quantum computer



Goal : map the deuteron problem into a simple Hamiltonian Usable in current quantum devices

Schematic Hamiltonian in a truncated HO basis



$$H_N = \sum_{n,n'=0}^{N-1} \langle n' | (T + V) | n \rangle a_{n'}^\dagger a_n$$

$$\langle n' | T | n \rangle = \frac{\hbar\omega}{2} \left[(2n + 3/2) \delta_n^{n'} - \sqrt{n(n + 1/2)} \delta_n^{n'+1} - \sqrt{(n + 1)(n + 3/2)} \delta_n^{n'-1} \right],$$

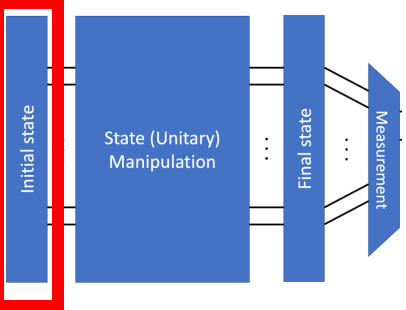
$$\langle n' | V | n \rangle = V_0 \delta_n^0 \delta_n^{n'}$$

The parameters are adjusted to reproduce the ground state and excited deuteron energy

$$V_0 = -5.68658111 \text{ MeV}, \quad \hbar\omega = 7 \text{ MeV}$$

- ➡ This problem is a simple tri-diagonal problem easy to solve on a classical computer with a direct matrix diagonalization
- ➡ Still it addresses the problem of a system of fermions and applications on real QC but can be considered as a proof of principle

“Nuclear” Many-Body problem: initial state state preparation



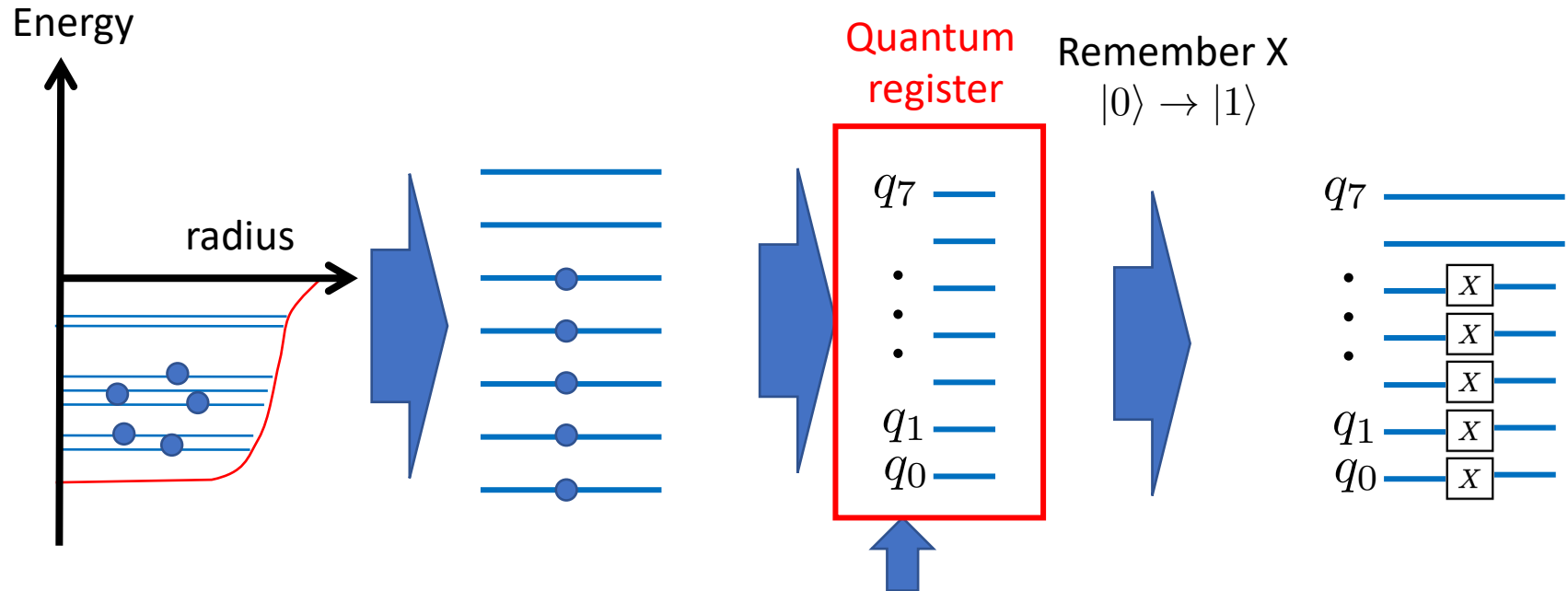
The nuclear “toy” shell in configuration interaction

Direct mapping from Fock to qubit space:

Assign one qubit to each single-particle state

$|0\rangle$ State is occupied

$|1\rangle$ State is unoccupied



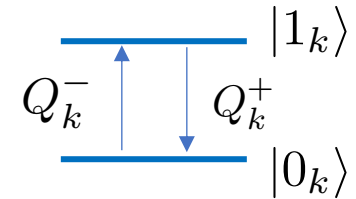
Initially all set to zero

One-difficulty – the antisymmetric nature of fermions

Mapping the Fock space into Qubits

$$|-\rangle = |0 \cdots 0\rangle$$

For qubits



For fermions

$$a_k^\dagger |-\rangle = |0 \cdots 0 \ 1_k \ 0 \cdots 0\rangle \iff Q_k^+ = \frac{1}{2} (X_k - iY_k) = \begin{bmatrix} 0 & 0 \\ 1 & 0 \end{bmatrix}$$

$$\{a_k, a_k^\dagger\} = 1 \qquad \{Q_k^-, Q_k^+\} = 1$$

Problem $\{a_k, a_l^\dagger\} = 0$ while $[Q_k^-, Q_l^+] = 0$

One possible solution (Jordan-Wigner transformation -1928)

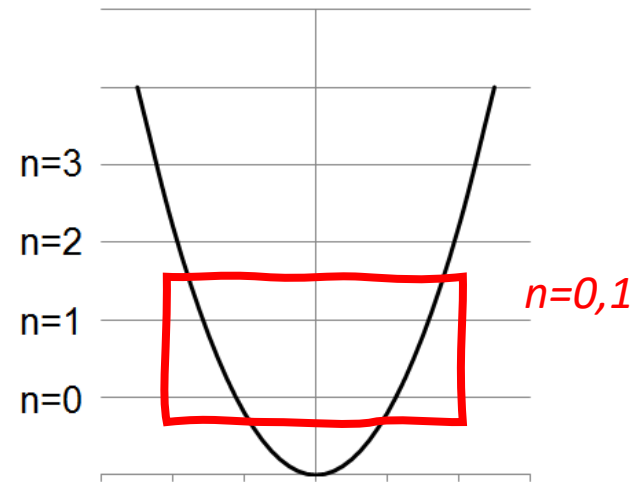
- 1 Order the index like in a lattice
- 2 Define new mapping



$$a_k^\dagger \rightarrow \prod_{k < j} (-\sigma_z^j) \sigma_j^+$$



The JWT method is systematic, how this works:



$$H_N = \sum_{n,n'=0}^{N-1} \langle n' | (T + V) | n \rangle a_{n'}^\dagger a_n$$

➔ $H_1 = \varepsilon_0 a_0^\dagger a_0 + \varepsilon_1 a_1^\dagger a_1 + v_{01} (a_1^\dagger a_0 + a_0^\dagger a_1)$

JWT mapping

$$a_0^\dagger \rightarrow Q_0^+$$

$$a_1^\dagger \rightarrow -Z_0 Q_1^+$$

⬇ $(Q_k^+ Z_k = -Q_k^+)$

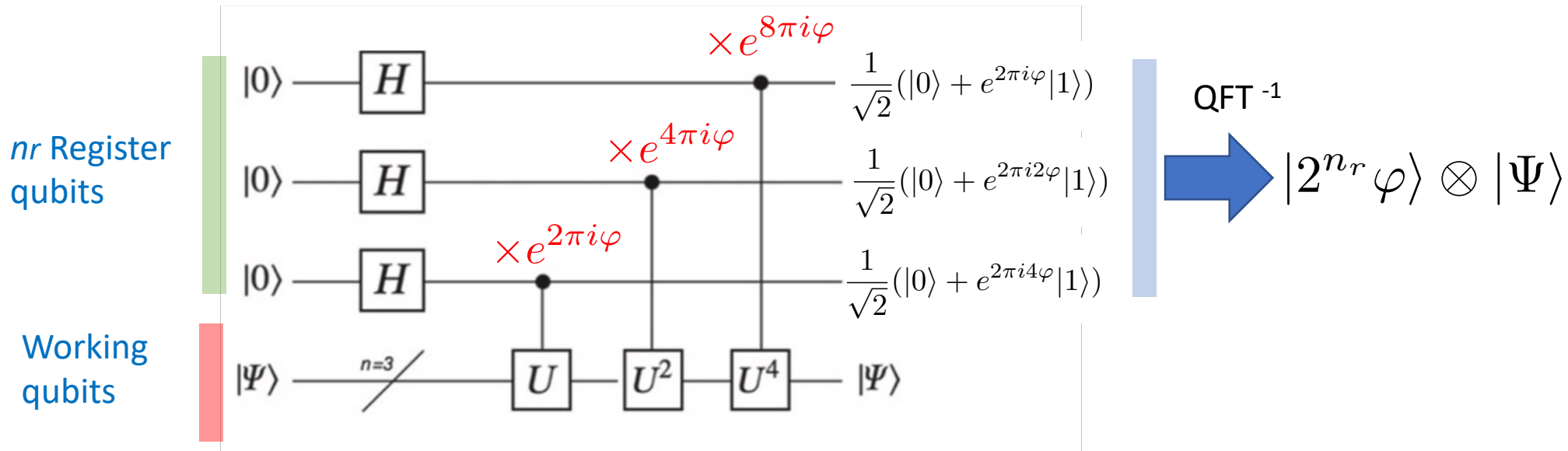
$$H_1 = \frac{\varepsilon_0}{2} (I - Z_0) + \varepsilon_1 (I - Z_1) + v_{01} (X_0 X_1 + Y_0 Y_1)$$

The Quantum-killer application for eigenvalue problem

the quantum-Phase estimation (QPE) algorithm

Assume a unitary operator U

Assume an eigenstate $|\Psi\rangle$ Such that $U|\Psi\rangle = e^{2\pi i\varphi}|\Psi\rangle$ $0 \leq \varphi < 1$

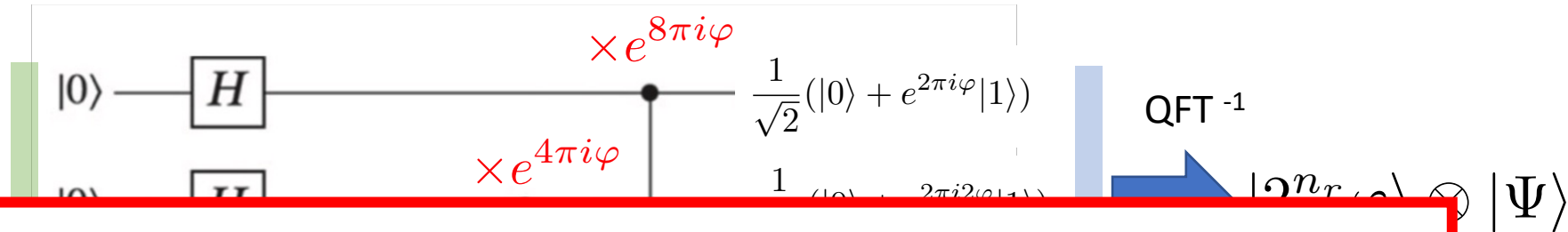


The quantum-Phase estimation (QPE) algorithm

For known eigenvalue problems

Assume a unitary operator U

Assume an eigenstate $|\Psi\rangle$ Such that $U|\Psi\rangle = e^{2\pi i\varphi}|\Psi\rangle$ $0 \leq \varphi < 1$



What is the meaning of $|2^{n_r}\varphi\rangle$?

$\varphi < 1$ Can be written as a binary fraction

$$\varphi = 0.j_1 \dots j_n = \frac{j_1}{2} + \frac{j_2}{2^2} + \dots + \frac{j_n}{2^n} \dots$$

Example $0.1011 \rightarrow 1 \times 0.5 + 0 \times 0.25 + 1 \times 0.125 + 1 \times 0.0625$
 $\rightarrow 0.6875$

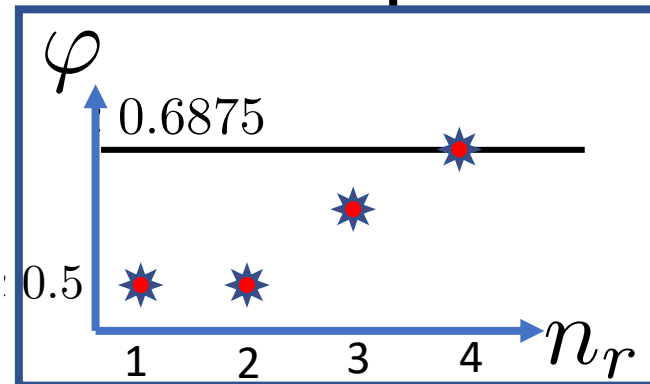
2^{-1}	=	0.5
2^{-2}	=	0.25
2^{-3}	=	0.125
2^{-4}	=	0.0625
2^{-5}	=	0.03125
2^{-6}	=	0.015625
2^{-7}	=	0.0078125

$n_r = 1 \rightarrow |2^{n_r}\varphi\rangle = |1\rangle \rightarrow \varphi \simeq 0.5$

$n_r = 2 \rightarrow |2^{n_r}\varphi\rangle = |10\rangle \rightarrow \varphi \simeq 0.5$

$n_r = 3 \rightarrow |2^{n_r}\varphi\rangle = |101\rangle \rightarrow \varphi \simeq 0.625$

$n_r = 4 \rightarrow |2^{n_r}\varphi\rangle = |1011\rangle \rightarrow \varphi \simeq 0.6875$

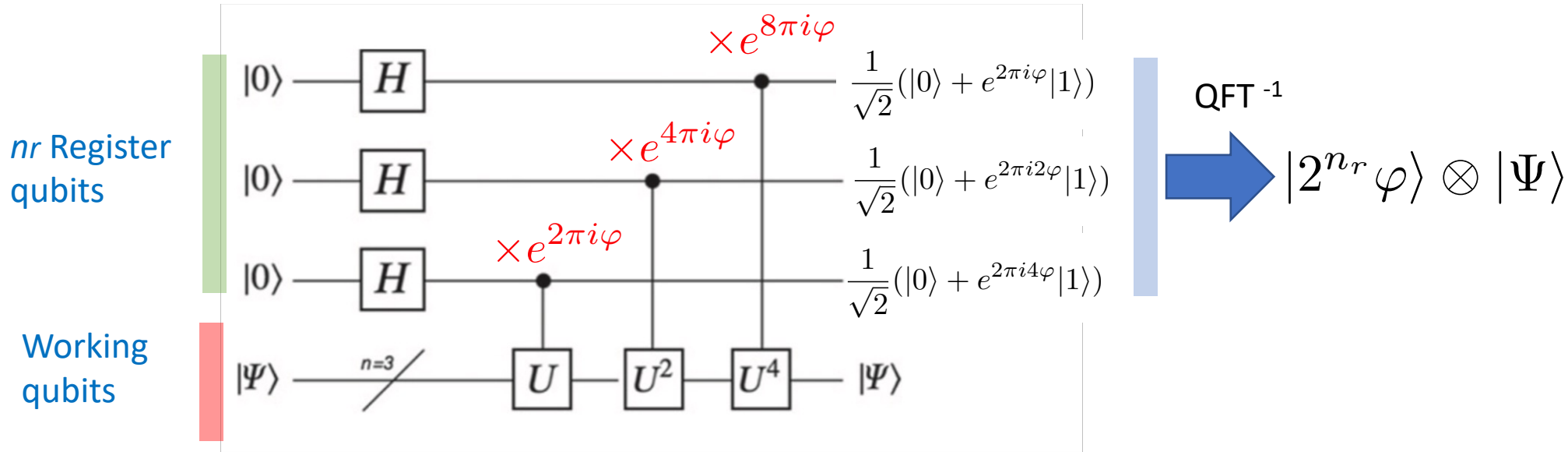


The quantum-Phase estimation (QPE) algorithm

For known eigenvalue problems

Assume a unitary operator U

Assume an eigenstate $|\Psi\rangle$ Such that $U|\Psi\rangle = e^{2\pi i\varphi}|\Psi\rangle$ $0 \leq \varphi < 1$



General Case

For an Hamiltonian problem

$$|\Psi\rangle = \sum_k \alpha_k |\phi_k\rangle \xrightarrow{\text{QPE}} \sum_k \alpha_k \underbrace{|\theta_k 2^{n_r}\rangle}_{\text{register}} \otimes \underbrace{|\phi_k\rangle}_{\text{eigenstate}}$$

One can simply use

$$U = e^{i\alpha(H - E_0)}$$

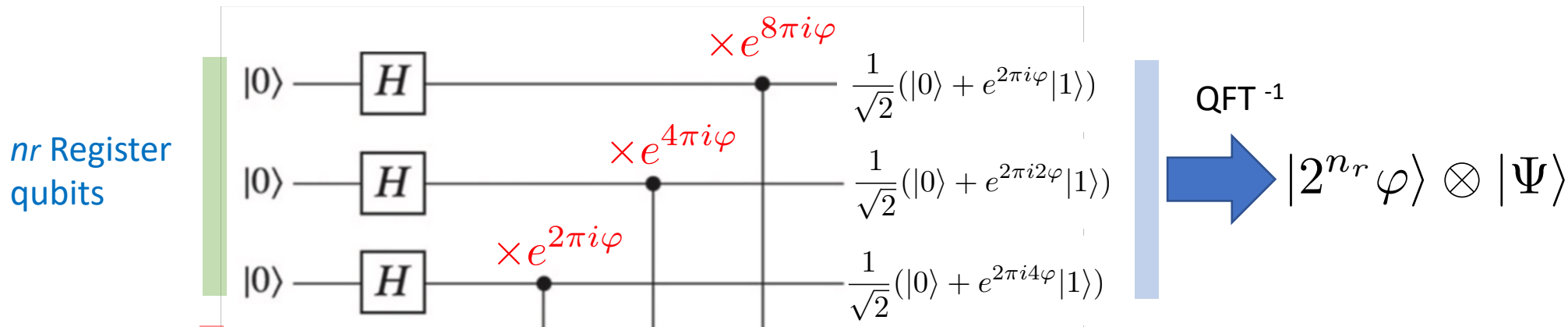
and get any eigenvalues with arbitrary precisions

The quantum-Phase estimation (QPE) algorithm

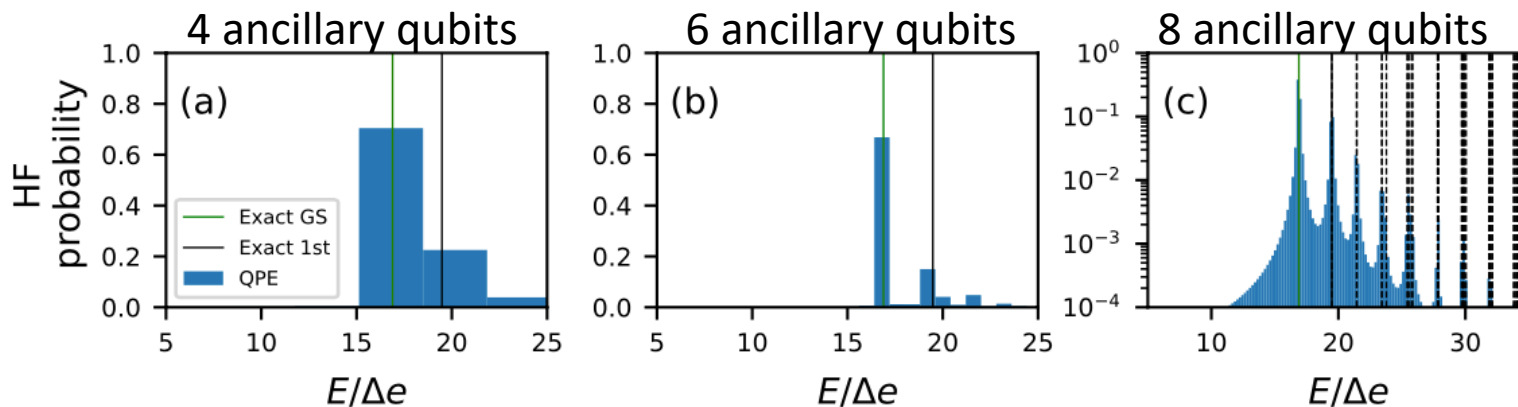
For known eigenvalue problems

Assume a unitary operator U

Assume an eigenstate $|\Psi\rangle$ Such that $U|\Psi\rangle = e^{2\pi i\varphi}|\Psi\rangle$ $0 \leq \varphi < 1$



And this works



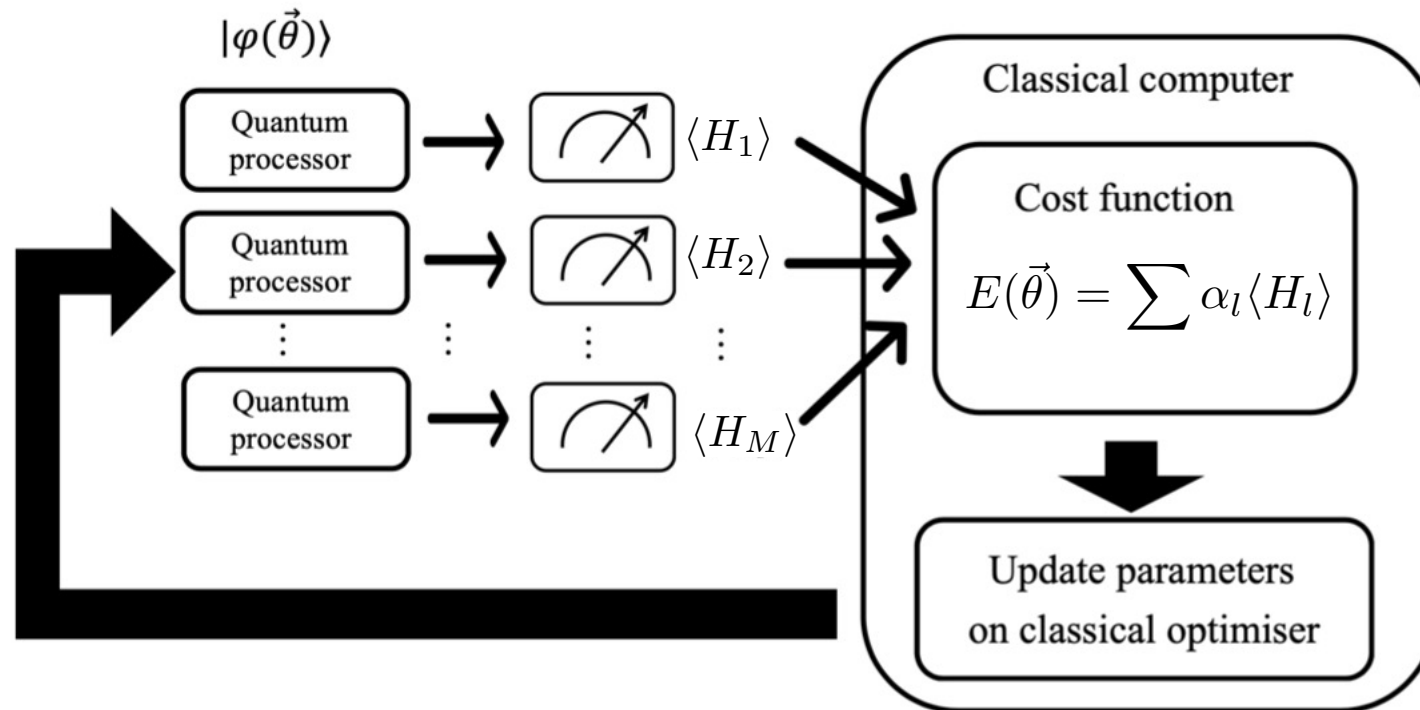
Unfortunately this is beyond the NISQ period

rary

The Variational Quantum Eigensolver (VQE) method

Main Idea of VQE

- Wave-functions can be better controlled by parametrizing them
- One can reduce the QC effort by computing expectation values of simpler operators $H = \sum \alpha_l H_l$
- The optimization task is left to a classical computer



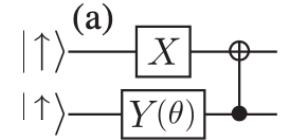
Adapted from S. Endo et al, arxiv:2011.01382
(see also K. Bharti et al., Rev. Mod. Phys. 94, 015004 (2022)).



Wave-function as parametric ansatz $|\Psi(\theta)\rangle = U(\theta)X_0 \cdots X_k|0 \cdots 0\rangle$

Hartree-Fock state with k occupied levels

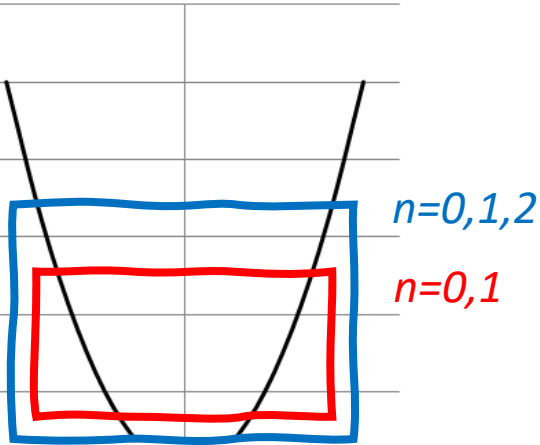
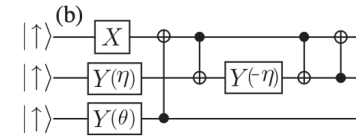
$$n=0,1 \quad U(\theta) \equiv e^{\theta(a_0^\dagger a_1 - a_1^\dagger a_0)} = e^{i(\theta/2)(X_0 Y_1 - X_1 Y_0)}$$



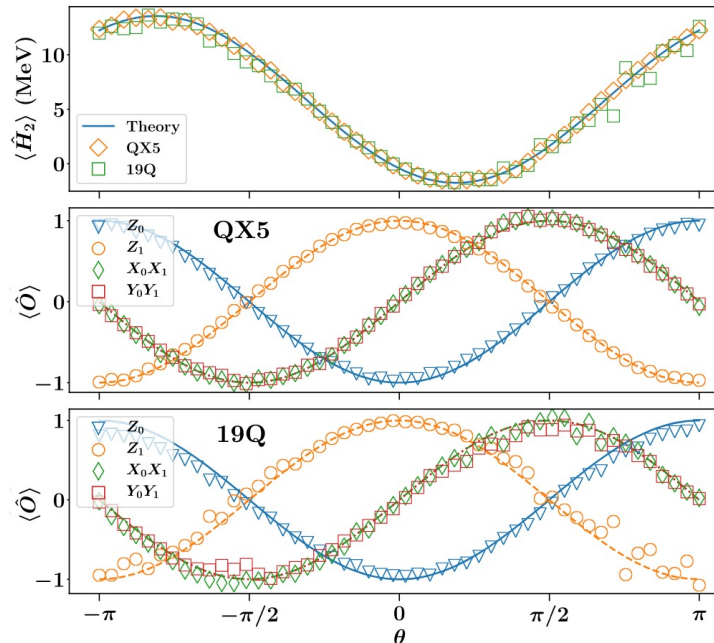
$$n=0,1,2 \quad U(\eta, \theta) \equiv e^{\eta(a_0^\dagger a_1 - a_1^\dagger a_0) + \theta(a_0^\dagger a_2 - a_2^\dagger a_0)}$$

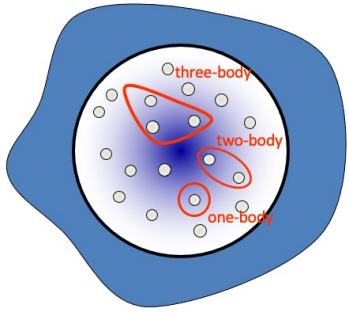
$$\approx e^{i(\eta/2)(X_0 Y_1 - X_1 Y_0)} e^{i(\theta/2)(X_0 Z_1 Y_2 - X_2 Z_1 Y_0)}$$

Note that this is a Unitary Coupled Cluster (UCC)

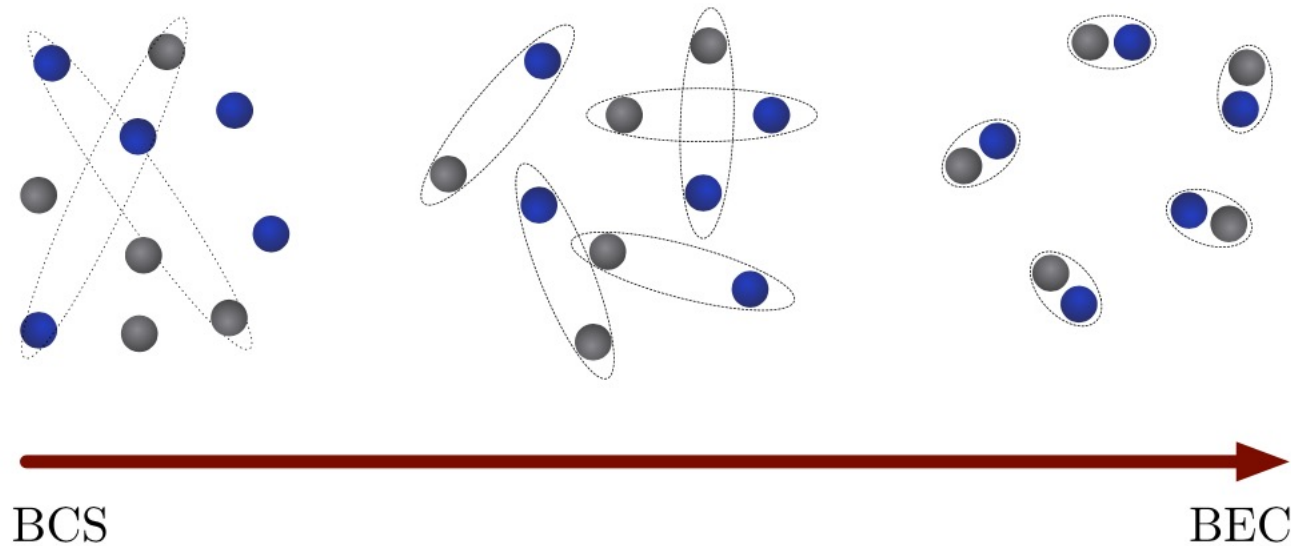


$$H_1 = \frac{\varepsilon_0}{2}(I - Z_0) + \varepsilon_1(I - Z_1) + v_{01}(X_0 X_1 + Y_0 Y_1)$$

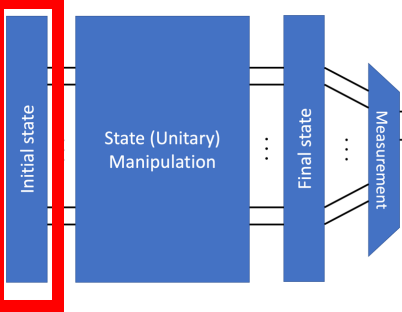




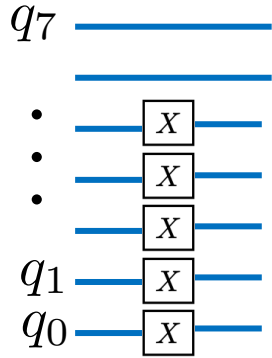
Third Example: Treatment of superconductivity



“Nuclear” Many-Body problem: initial state preparation



Starting point : Slater determinants



Demystifying QC

Illustration with qiskit-IBM

```
import numpy as np
from qiskit import *
%matplotlib inline
import math
```

Creation of the circuit

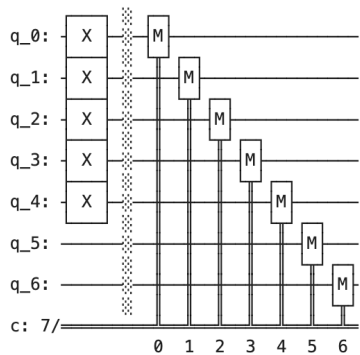
```
nq=7
nc=7
qr = QuantumRegister(nq, 'q') # qubit of interest + register qubits
cr = ClassicalRegister(nc, 'c') # classical register
# name of the circuit
mycircuit = QuantumCircuit(qr, cr)

for j in range(5):
    mycircuit.x(j)

mycircuit.barrier()

for j in range(7):
    mycircuit.measure(j,j)

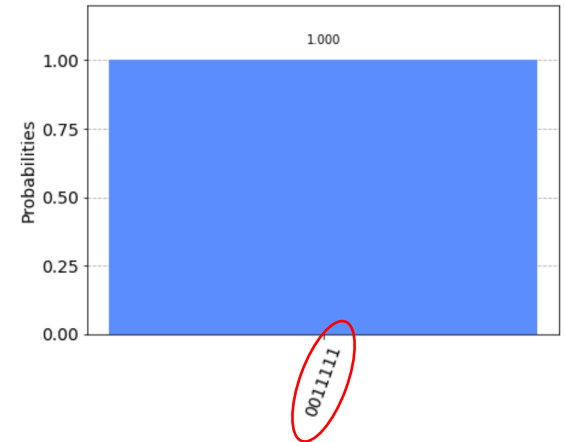
#mycircuit.draw()
print(mycircuit)
```

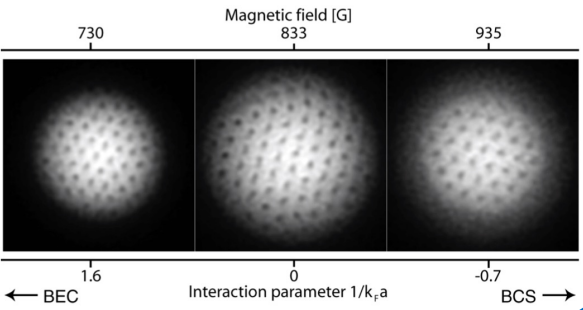


Running the circuit

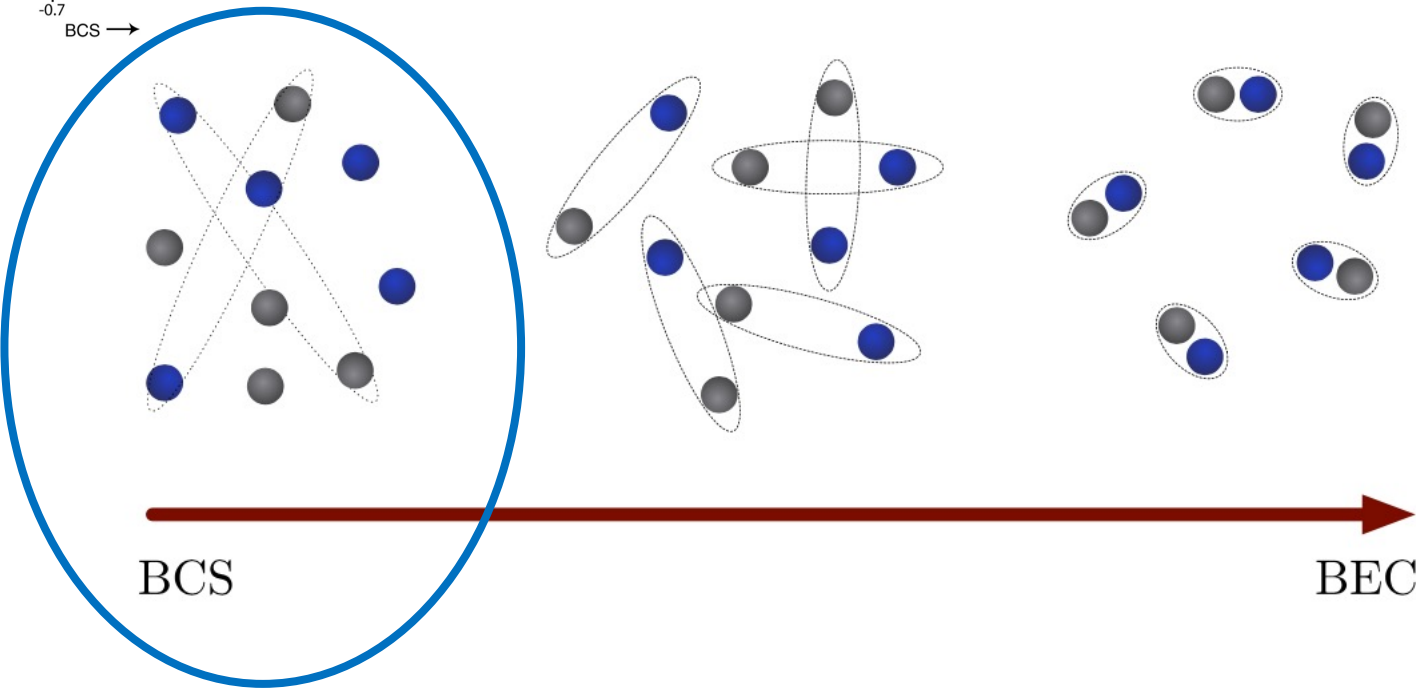
```
[3]: # building our own normalized histo
# Running the code !
backend = Aer.get_backend('qasm_simulator')
shots = 10000
results = execute(mycircuit, backend=backend, shots=shots).result()
answer = results.get_counts()

print(answer)
plot_histogram(answer)
```





Cooper pairs and superfluidity are rather universal phenomena: (condensed matter, Atomic physics, Nuclear physics, ...)



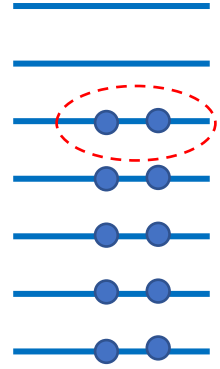
This problem is an archetype of spontaneous symmetry breaking. A “easy” way to describe it is to break the particle number symmetry, i.e. consider wave-function that mixes different particle number

Example

$$|\Phi_0\rangle = \prod_i (u_i + v_i a_i^\dagger a_i^\dagger) |-\rangle$$

➡ Mixes states with 0, 2, 4, ... particles

We say that a symmetry (particle number) is broken

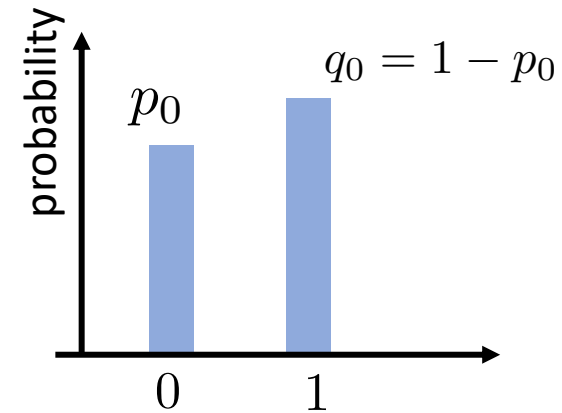
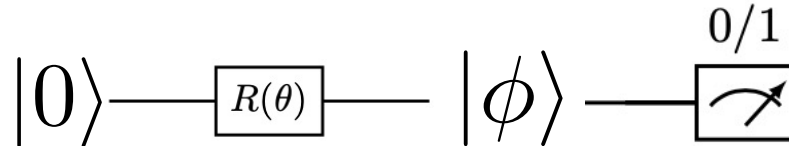


Preparing state that mixes particle number

I assign a qubit to each qubit

$$|\phi\rangle = \sqrt{p_0}|0\rangle + \sqrt{1-p_0}|1\rangle$$

Measuring the qubit gives the probability



On the Qiskit toolkit

```
[1]: import numpy as np
      from qiskit import *
      %matplotlib inline
      import math

      from qiskit.visualization import plot_histogram
```

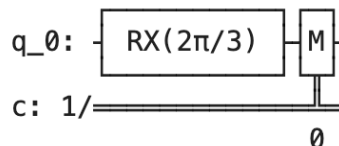
Creation of the circuit

```
[2]: nq=1
      nc=1
      qr = QuantumRegister(nq, 'q') # qubit of interest + register qubits
      cr = ClassicalRegister(nc, 'c') # classical register
      # name of the circuit
      mycircuit = QuantumCircuit(qr, cr)

      #make the rotation
      angle = 4*2*math.pi/12

      mycircuit.rx(angle,0)
      mycircuit.measure(0,0)

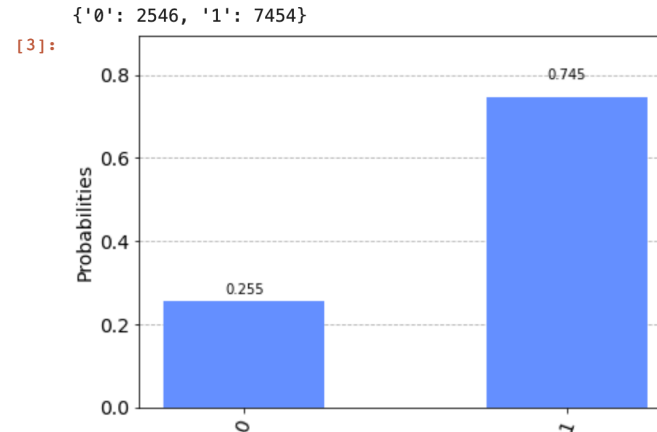
      #mycircuit.draw()
      print(mycircuit)
```



Running the circuit

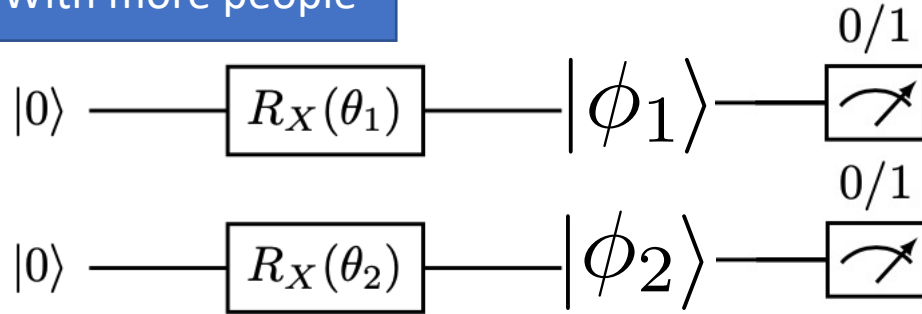
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      print(answer)
      plot_histogram(answer)
```



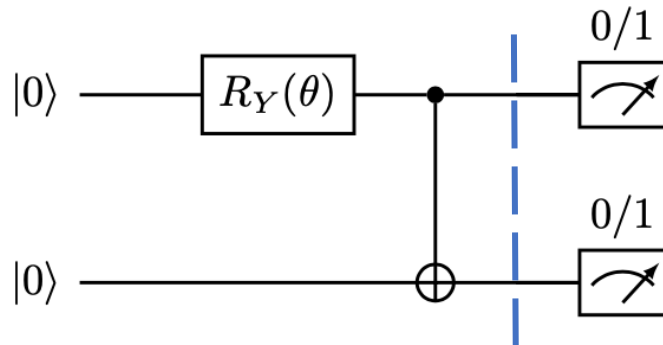
Preparing state that mixes particle number

With more people



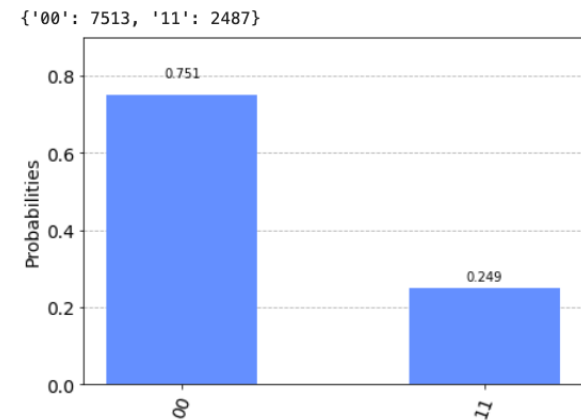
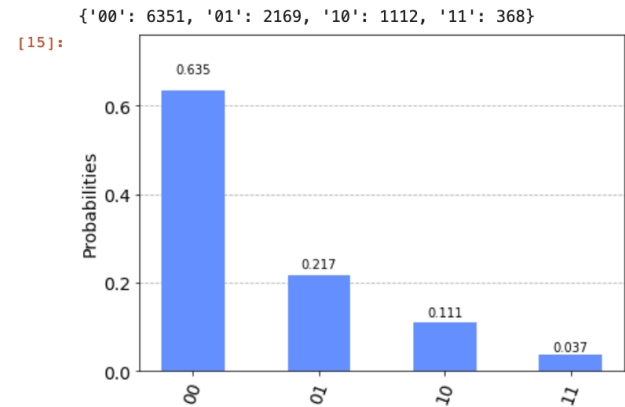
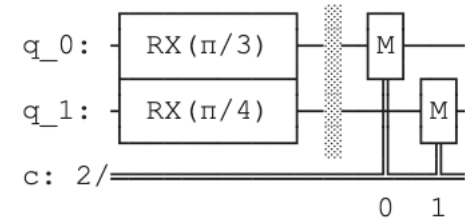
$$|\phi_i\rangle = \alpha_i|0\rangle + \beta_i|1\rangle$$

Qubit can be entangled



$$|\phi_{12}\rangle = \alpha|00\rangle + \beta_i|11\rangle$$

Here I created a Bell state



Application to the N-body pairing problem

Hamiltonian and initial state

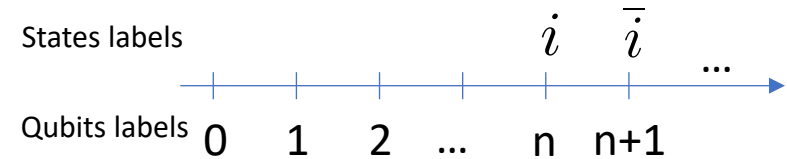
Pairing Hamiltonian

$$H_P = \sum_{i>0} \varepsilon_i (a_i^\dagger a_i + a_{\bar{i}}^\dagger a_{\bar{i}}) - g \sum_{i,j>0} a_i^\dagger a_{\bar{i}}^\dagger a_{\bar{j}} a_j$$

Jordan-Wigner trans:

$$\frac{1}{2}(I_i - Z_i)$$

State ordering is important !



$$a_i^\dagger a_{\bar{i}}^\dagger \longrightarrow Q_n^+ Q_{n+1}^+$$

Initial (symmetry breaking) state preparation

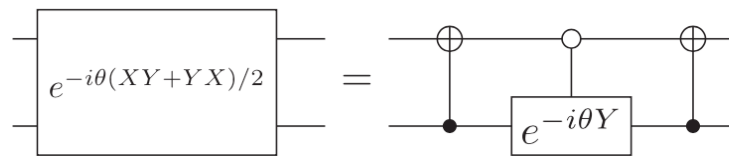
$$|\Psi\rangle = \exp \left\{ - \sum_{i>0} \varphi_i (a_i^\dagger a_{\bar{i}}^\dagger - a_{\bar{i}} a_i) \right\} |0\rangle$$

$$\varphi_i = \varphi \longrightarrow$$

$$|\Psi\rangle = \prod_{n>0} e^{i\varphi(X_n Y_{n+1} + Y_n X_{n+1})/2} |-\rangle$$

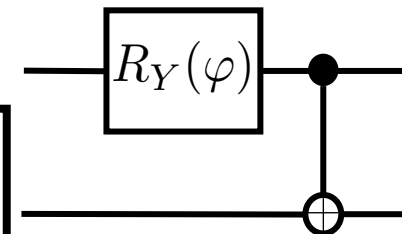
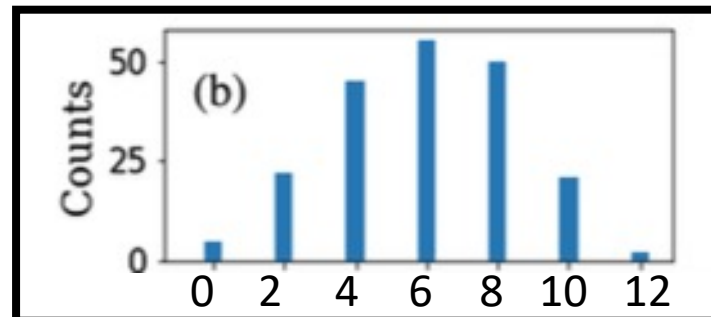
Equivalent universal gate on pairs

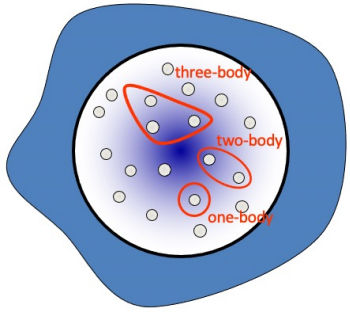
Simplified circuit (generalized Bell state)



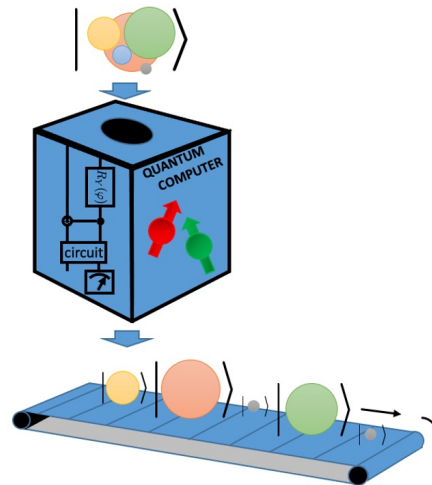
$$|\Psi\rangle = \prod_n \left[\cos\left(\frac{\varphi}{2}\right) I_n \otimes I_{n+1} + \sin\left(\frac{\varphi}{2}\right) Q_n^+ Q_{n+1}^+ \right] |-\rangle$$

Zhang Jiang et al,
Phys. Rev. Applied 9, 044036 (2018).

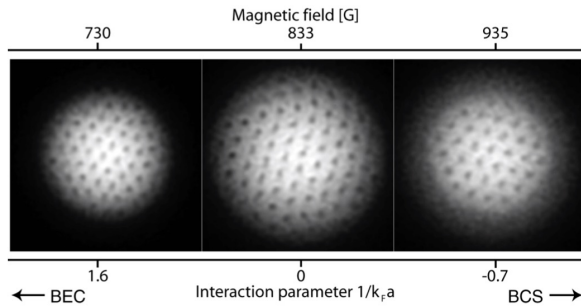




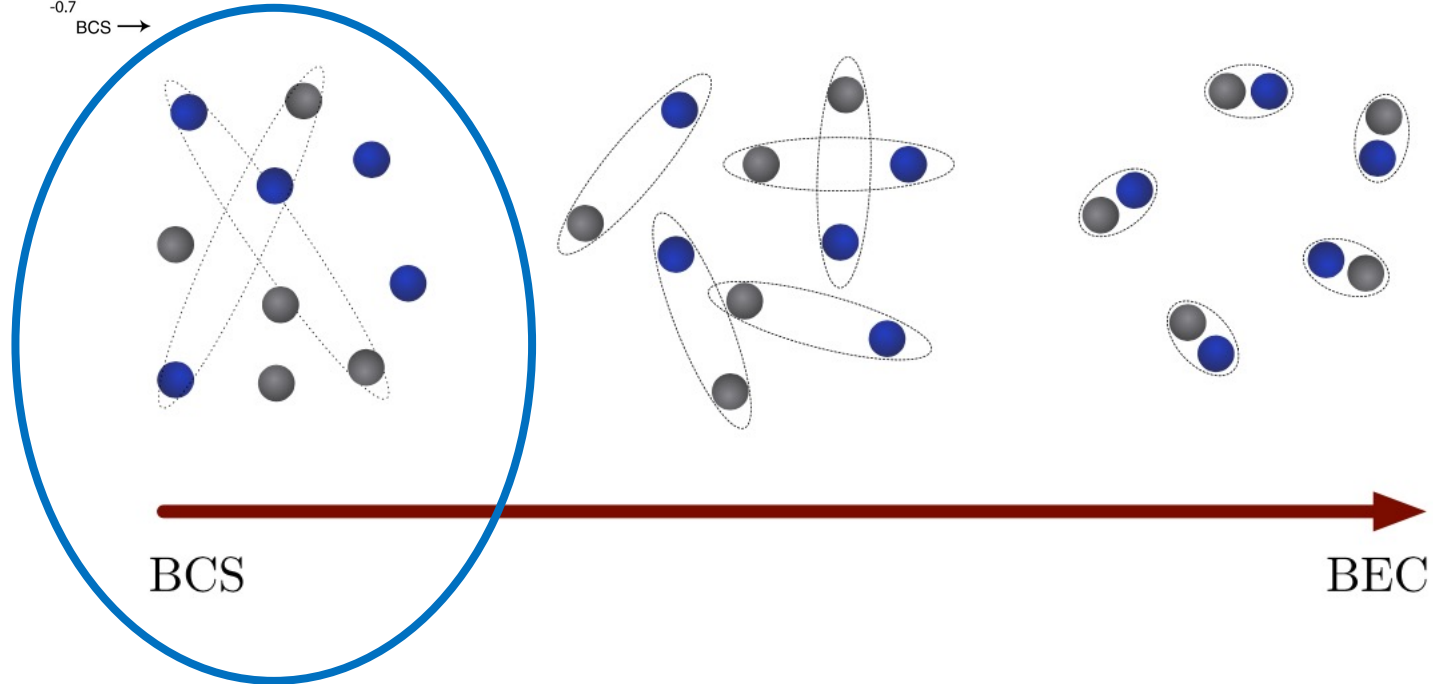
Fourth Example: restoration of broken symmetries



But what is the connection with interacting systems ???



Cooper pairs and superfluidity are rather universal phenomena: (condensed matter, Atomic physics, Nuclear physics, ...)



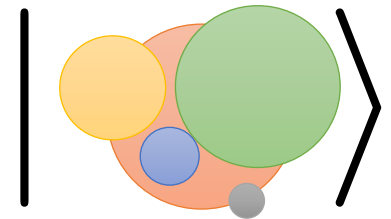
This problem is an archetype of spontaneous symmetry breaking. A “easy” way to describe it is to break the particle number symmetry, i.e. consider wave-function that mixes different particle number

Example

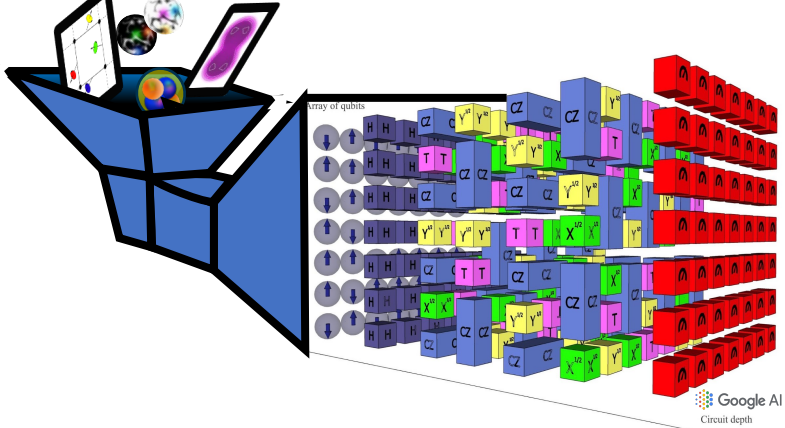
$$|\Phi_0\rangle = \prod_i (u_i + v_i a_i^\dagger a_i^\dagger) |-\rangle$$

➔ Mixes states with 0, 2, 4, ... particles

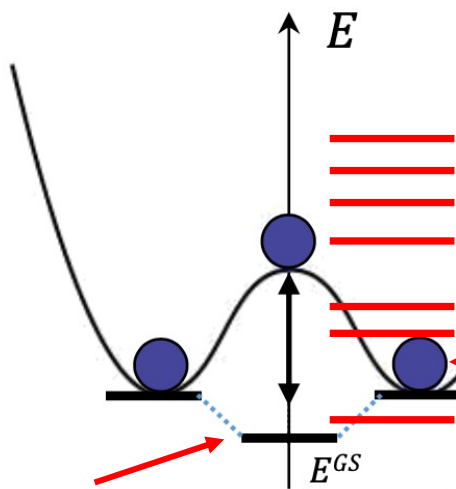
We say that a symmetry (particle number) is broken



But ultimately number of Particle should be restored !



Further
Quantum
or hybrid
Quantum-Classical
Post-processing



2 Symmetry Restored (SR) state (multi-reference)

1 Symmetry Breaking (SB) state

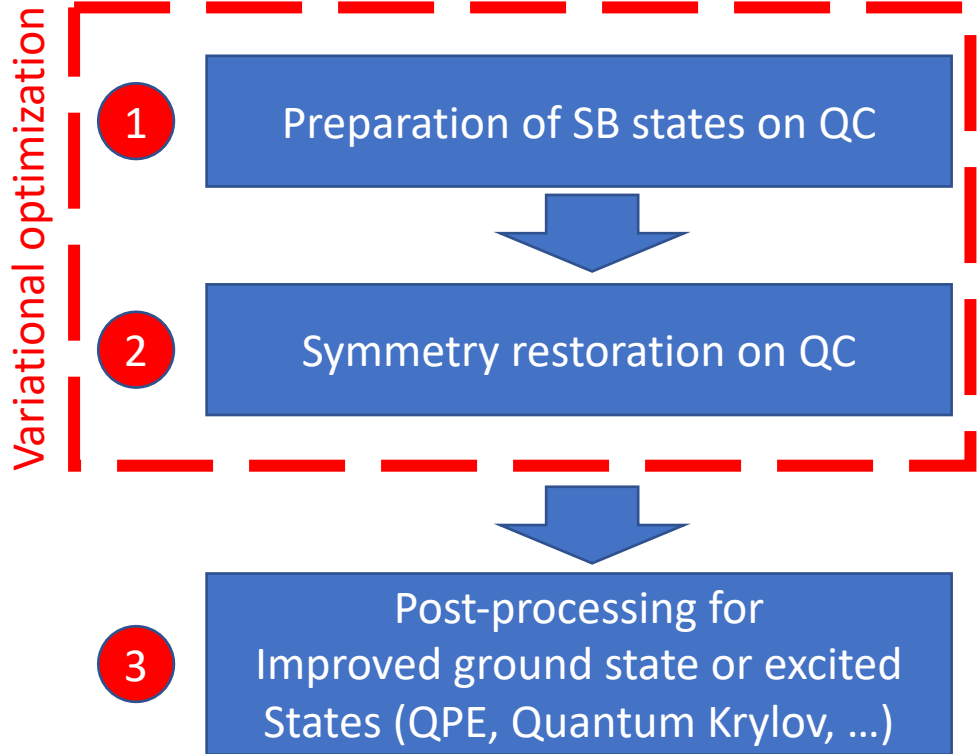


Illustration with small superconductors

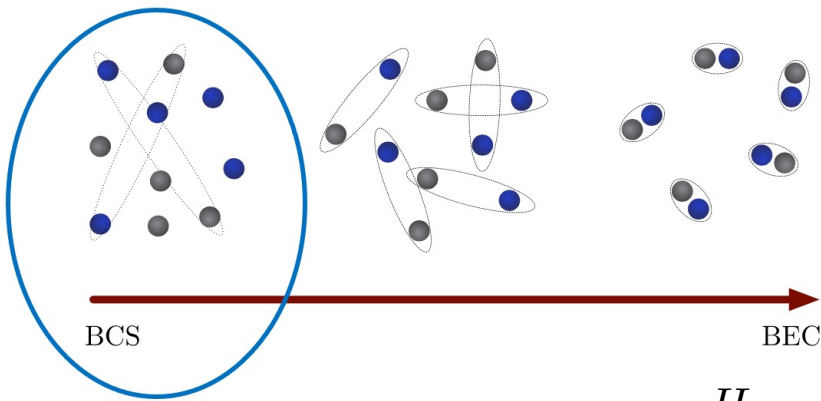
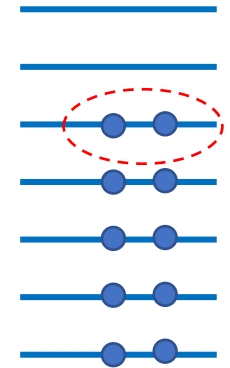


Illustration with the Richardson Hamiltonian

$$H_P = \sum_{i>0} \varepsilon_i (a_i^\dagger a_i + a_{\bar{i}}^\dagger a_{\bar{i}}) - g \sum_{i,j>0} a_i^\dagger a_{\bar{i}}^\dagger a_{\bar{j}} a_j$$



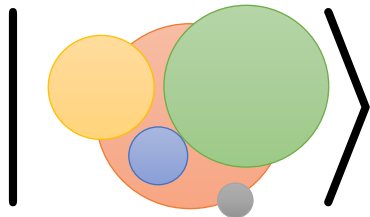
This problem is an archetype of spontaneous symmetry breaking.
An “easy” way to describe it is to break the particle number symmetry, i.e. consider wave-function that mixes different particle number

Example

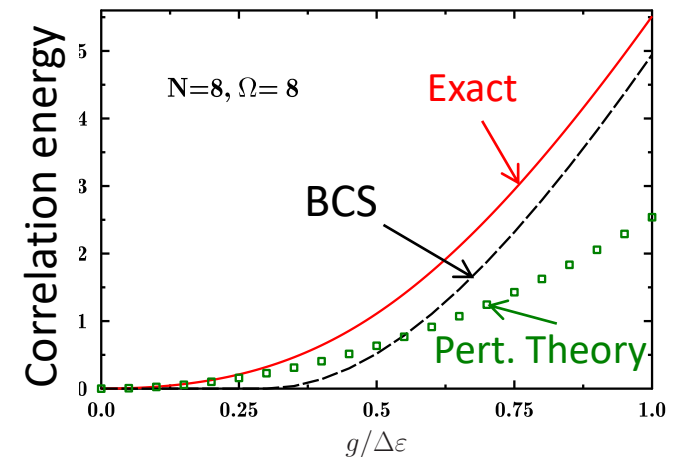
$$|\Phi_0\rangle = \prod_i (u_i + v_i a_i^\dagger a_{\bar{i}}^\dagger) |-\rangle$$

➔ Mixes states with 0, 2, 4, ... particles

The particle number - U(1) symmetry) is broken



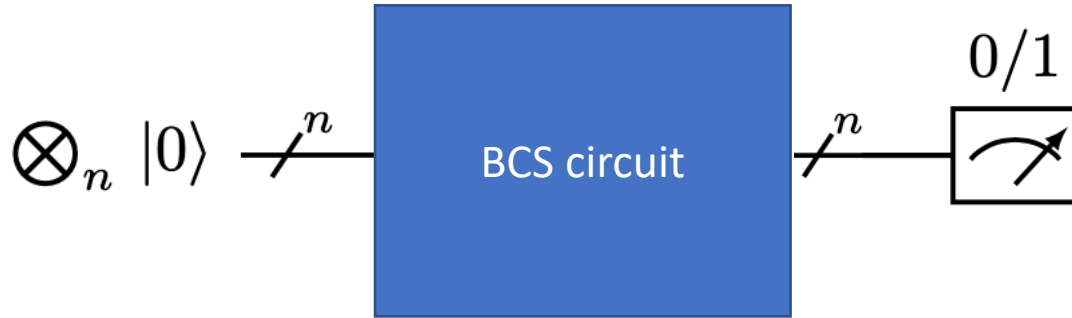
But ultimately number of Particle should be restored !



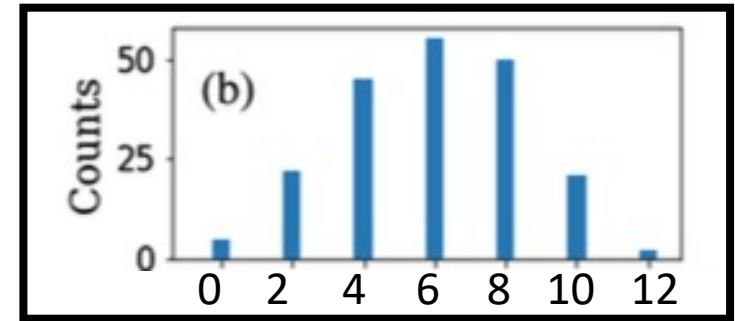
Restoration of particle number symmetry

The counting statistics problem

Direct estimate of Counting statistics



Example of mixing for 12 qubits (with qiskit)



Projection on particle number

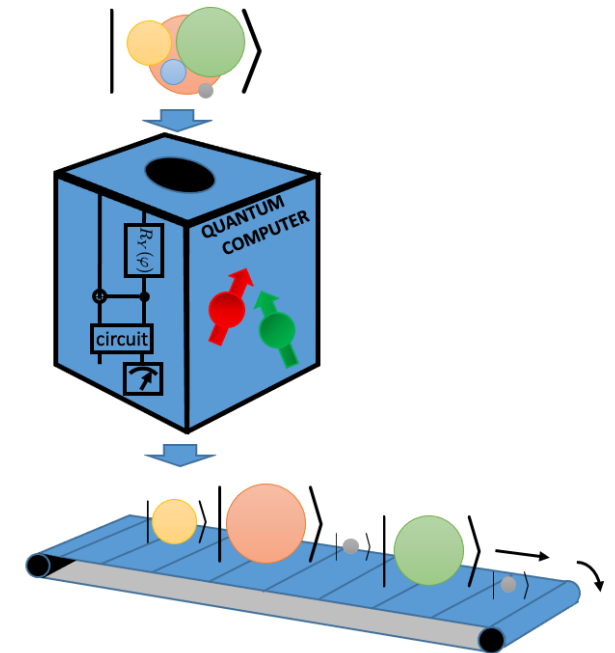
$$|\Psi\rangle = \sum_N c_N |N\rangle \rightarrow |N\rangle$$

For 2 qubits

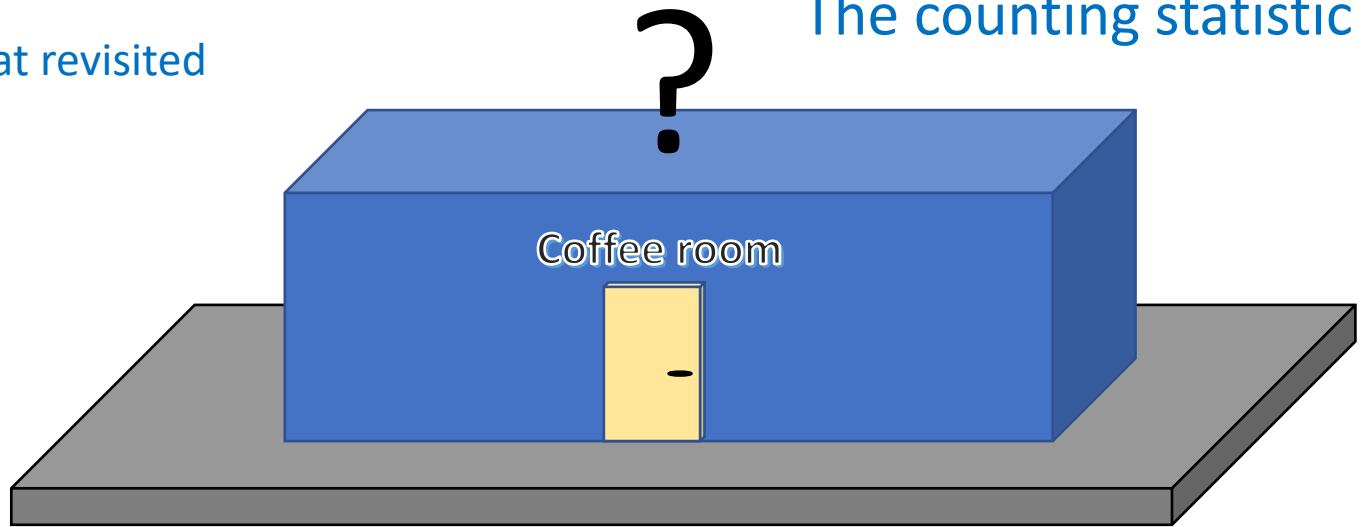
$$|\Psi\rangle = \alpha|00\rangle + \beta|01\rangle + \gamma|10\rangle + \delta|11\rangle$$

$|N=0\rangle$
 $\propto |N=1\rangle$
 $|N=2\rangle$

➔ A possible way to perform the projection is to use The Quantum-Phase-Estimation method with N itself



The schroedinger cat revisited



A more specific example



4 persons



3 persons



2 persons

With many events we can do Probabilities, statistical analysis

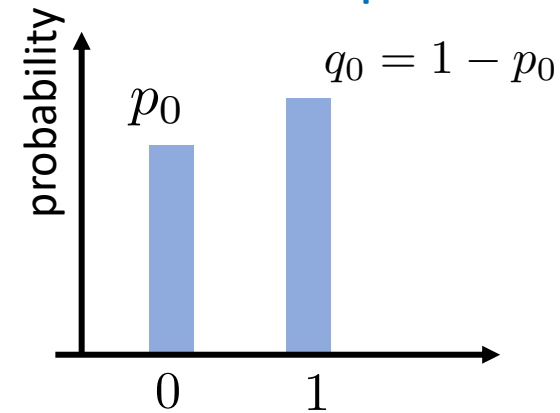
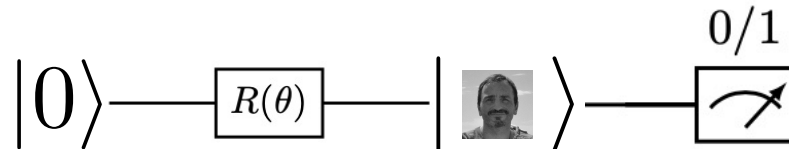
The counting statistic problem

In quantum systems

I assign a qubit to each person

$$|\text{person}\rangle = \sqrt{p_0}|0\rangle + \sqrt{1-p_0}|1\rangle$$

Measuring the qubit gives the probability



Demystifying QC

Illustration with qiskit

```
[1]: import numpy as np
      from qiskit import *
      %matplotlib inline
      import math

      from qiskit.visualization import plot_histogram
```

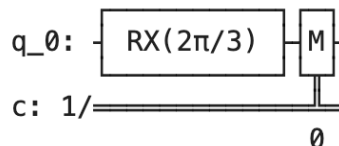
Creation of the circuit

```
[2]: nq=1
      nc=1
      qr = QuantumRegister(nq, 'q') # qubit of interest + register qubits
      cr = ClassicalRegister(nc, 'c') # classical register
      # name of the circuit
      mycircuit = QuantumCircuit(qr, cr)

      #make the rotation
      angle = 4*2*math.pi/12

      mycircuit.rx(angle,0)
      mycircuit.measure(0,0)

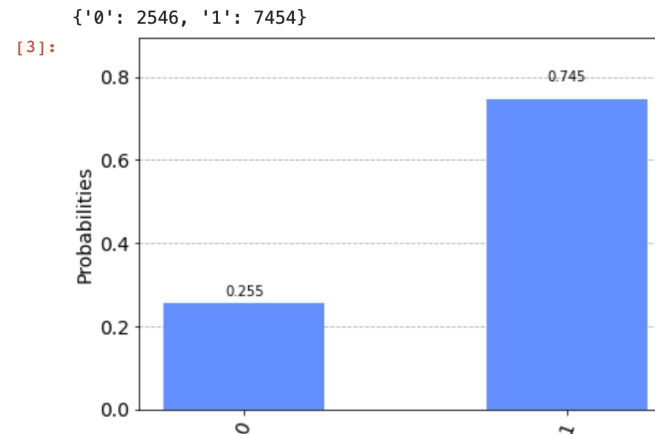
      #mycircuit.draw()
      print(mycircuit)
```



Running the circuit

```
[3]: # building our own normalized histo
      # Running the code !
      backend = Aer.get_backend('qasm_simulator')
      shots = 10000
      results = execute(mycircuit, backend=backend, shots=shots).result()
      answer = results.get_counts()

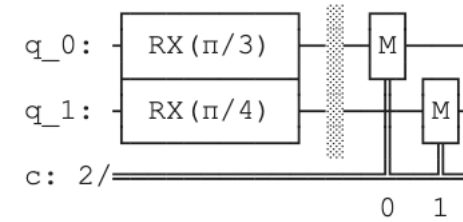
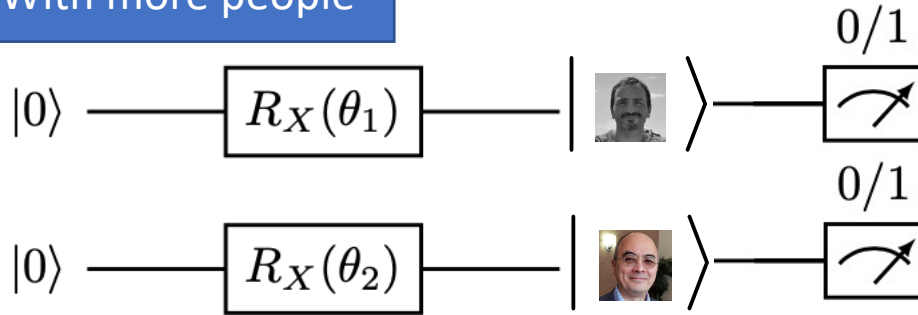
      print(answer)
      plot_histogram(answer)
```



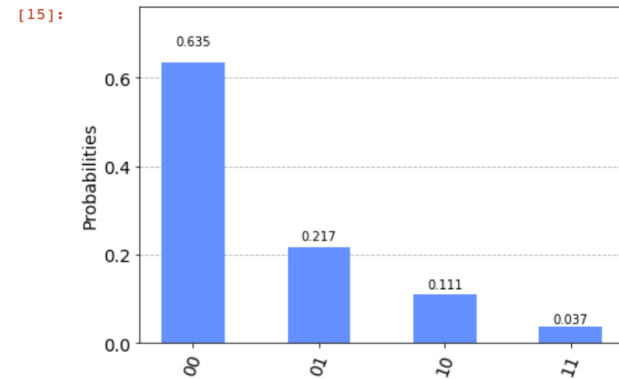
The counting statistic problem

In quantum systems

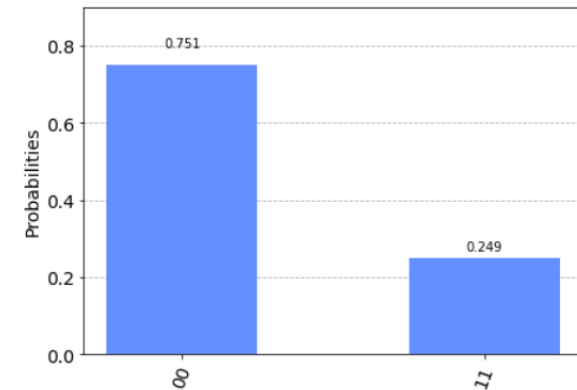
With more people



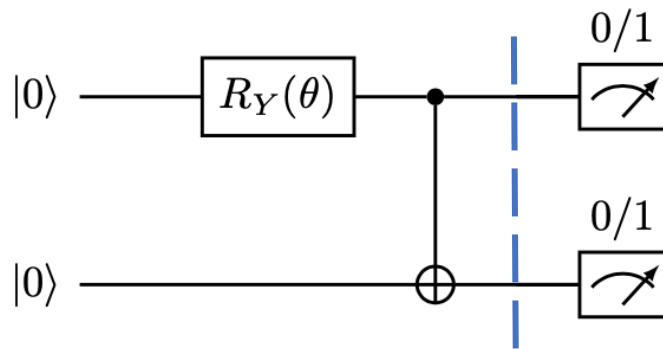
{'00': 6351, '01': 2169, '10': 1112, '11': 368}



{'00': 7513, '11': 2487}



People can be entangled



$$|\Phi\rangle = \alpha |\text{[Person 1] [Person 2]}\rangle + \beta |\text{[Person 1] [Person 2]}\rangle$$

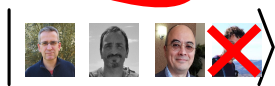
Here I created a Bell state

The counting statistic problem without destroying the wave-function

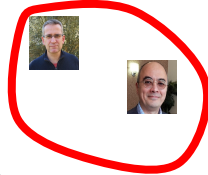
Initial wave-function

$$|\Phi\rangle = \alpha|\text{img1 img2 img3 img4}\rangle + \beta|\text{X img2 img3 img4}\rangle + \gamma|\text{X X img3 img4}\rangle + \delta|\text{img1 X img3 img4}\rangle + \dots$$

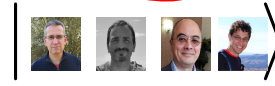
Event 1



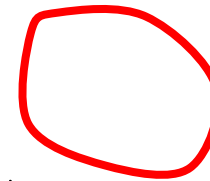
Event 2



Event 3



Event 4



...

After the measurement the wave-function collapse to one of the state



Schrodinger's Cat

$$\frac{1}{\sqrt{2}}|\text{cat alive}\rangle + \frac{1}{\sqrt{2}}|\text{cat dead}\rangle$$

If I open the box



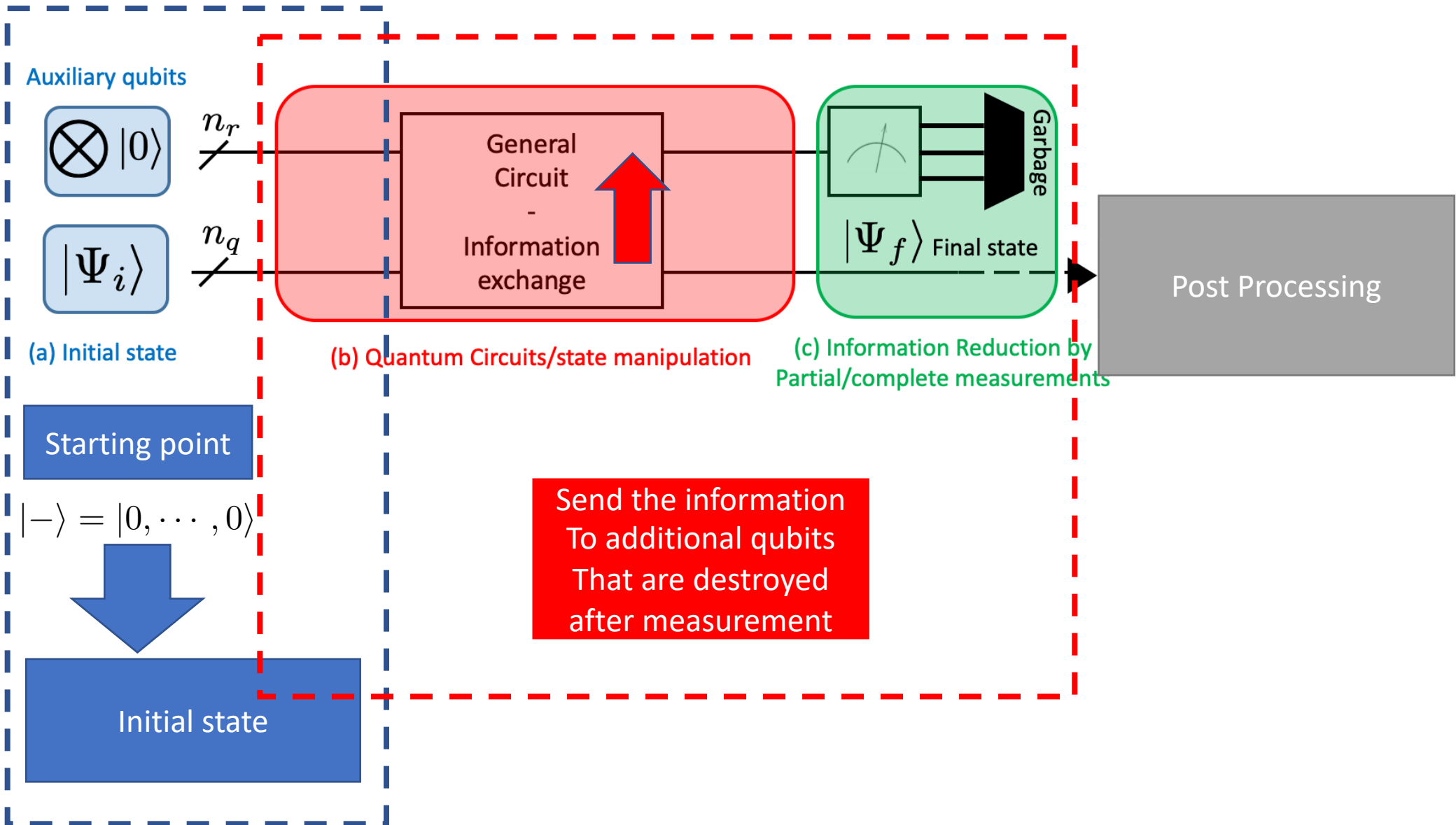
or

A more difficult problem

I want to select the component with 3 persons without completely destroying it

$$|\Phi\rangle = +\beta'|\text{X img2 img3 img4}\rangle + \delta'|\text{img1 X img3 img4}\rangle + \dots$$

Non-destructive counting on a quantum computer

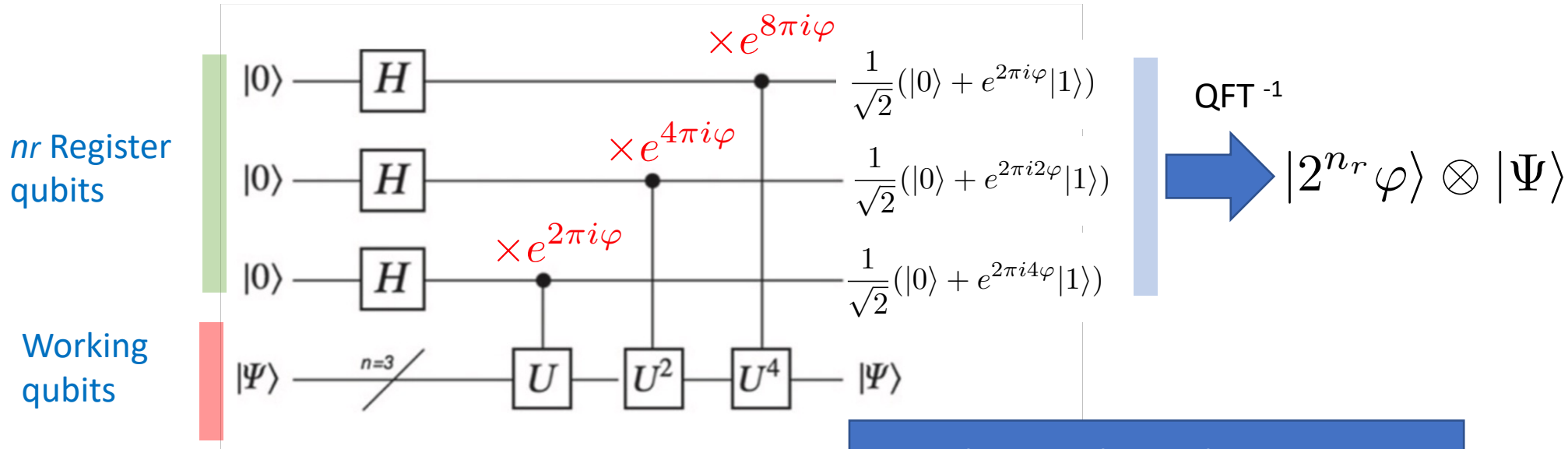


The quantum-Phase estimation (QPE) algorithm

For known eigenvalue problems

Assume a unitary operator U

Assume an eigenstate $|\Psi\rangle$ Such that $U|\Psi\rangle = e^{2\pi i\varphi}|\Psi\rangle$ $0 \leq \varphi < 1$



General Case

For the particle number projection

$$|\Psi\rangle = \sum_k \alpha_k |\phi_k\rangle \xrightarrow{\text{QPE}} \sum_k \alpha_k \underbrace{|2^{n_r}\varphi\rangle}_{\text{register}} \otimes \underbrace{|\phi_k\rangle}_{\text{eigenstate}}$$

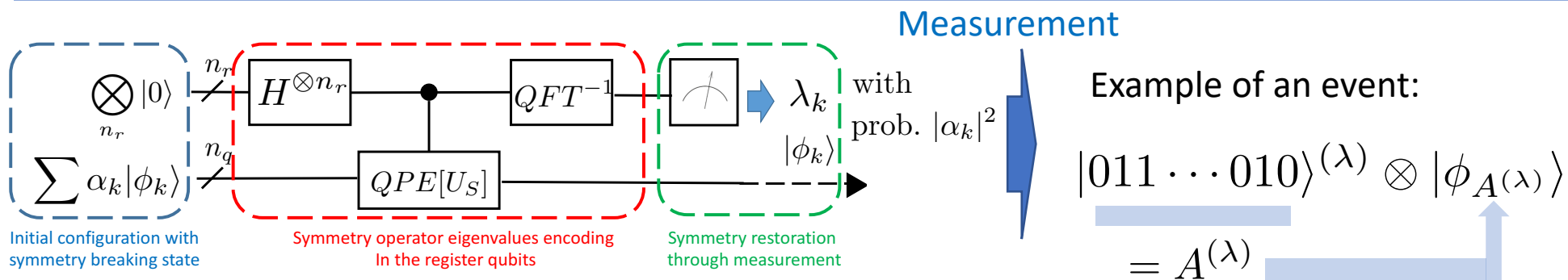
$$U = U_N = e^{2\pi i \frac{N}{2^{n_r}}} \quad \text{with} \quad N = \frac{1}{2} \sum_i (I_i - Z_i)$$

Assume eigenvalues $\{0, 1, \dots, A\}$

Constraint: $0 \leq \frac{A}{2^{n_r}} < 1$ then $\frac{A}{2^{n_r}} = 0.a_1 \dots a_{n_r-1}$

If I measure given binary number in the ancillary qubit. After measurement, I have the projection on the associated particle number component

Eigenvalues-Ground state and excited states



BCS/HFB state

$$|\Psi\rangle = \prod_n \left[\cos\left(\frac{\varphi}{2}\right) I_n \otimes I_{n+1} + \sin\left(\frac{\varphi}{2}\right) Q_n^+ Q_{n+1}^+ \right] |-\rangle$$

Measurement

Projected BCS or with varying number of particles

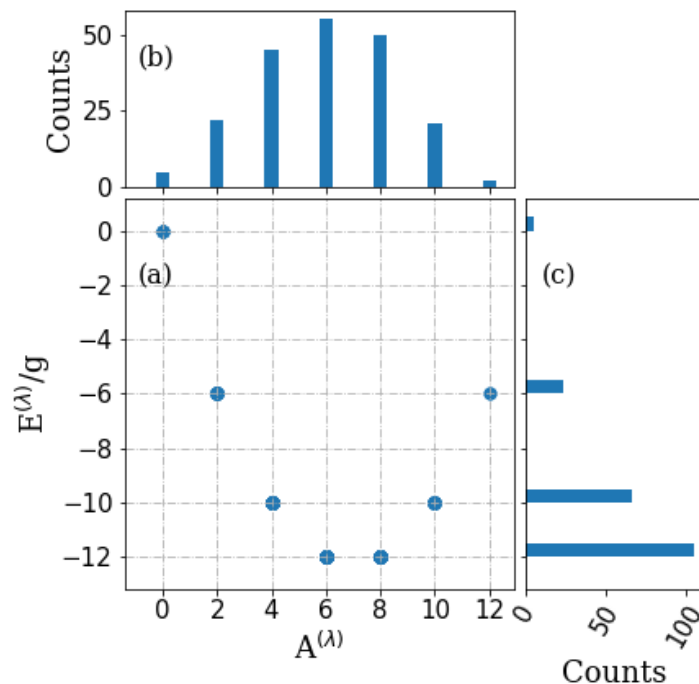
Degenerate case

$$H_P = -g \sum_{i,j>0} a_i^\dagger a_j^\dagger a_{\bar{j}} a_{\bar{i}}$$

$$\langle \phi_{A^{(\lambda)}} | H | \phi_{A^{(\lambda)}} \rangle$$

H was encoded on the full Fock space with $A < n_q$
 For the degenerate case, this should give the exact solution

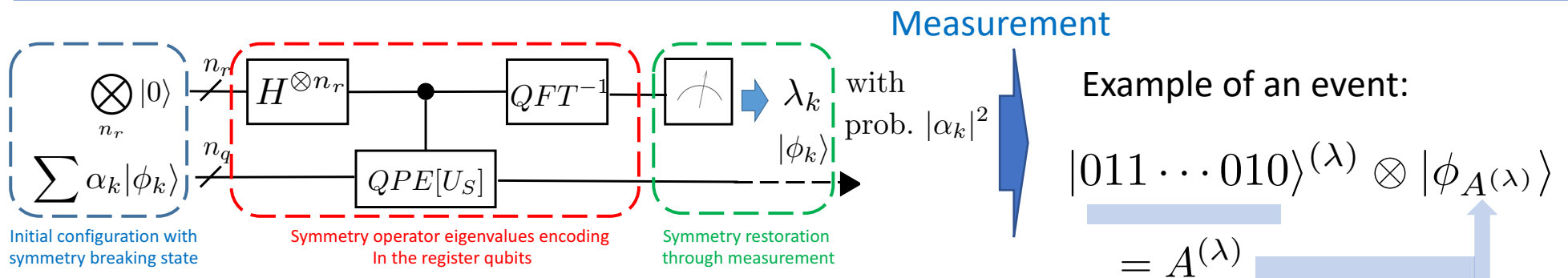
6 pairs



Exact solution

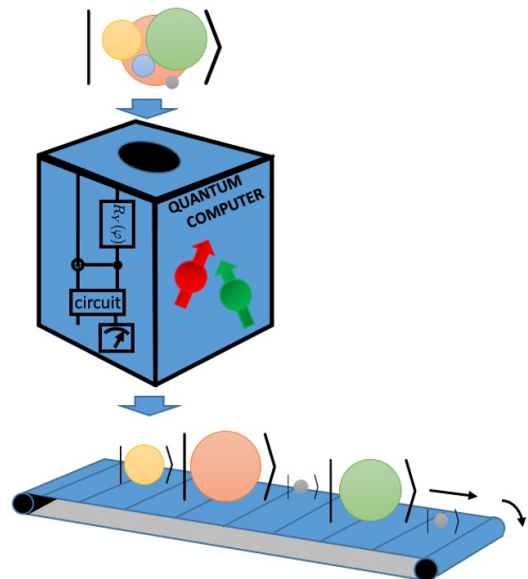
$$E/g = -\frac{1}{4}(A - \nu)(2n_q - A - \nu - 2).$$

Eigenvalues-Ground state and excited states

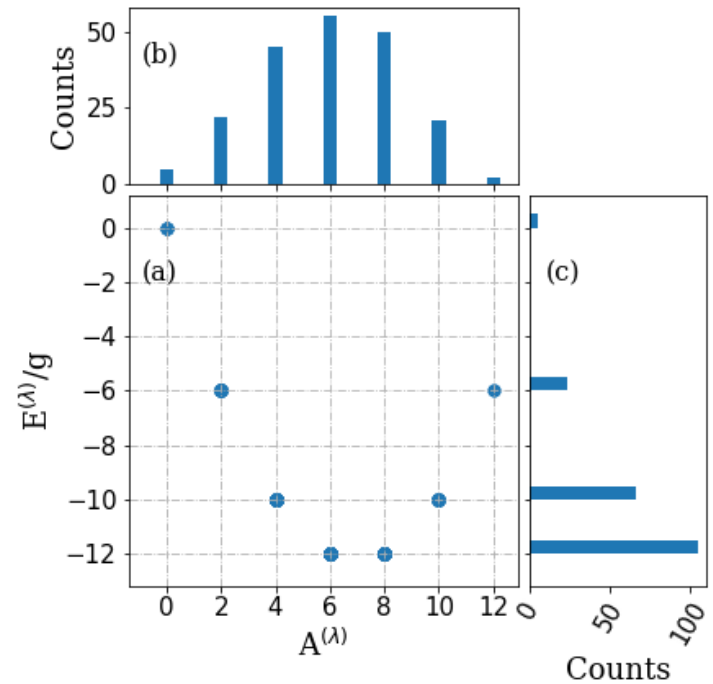


BCS/HFB state

$$|\Psi\rangle = \prod_n \left[\cos\left(\frac{\varphi}{2}\right) I_n \otimes I_{n+1} + \sin\left(\frac{\varphi}{2}\right) Q_n^+ Q_{n+1}^+ \right] |-\rangle$$



6 pairs



Exact solution

$$E/g = -\frac{1}{4}(A - \nu)(2n_q - A - \nu - 2).$$

A schematic view

Making projection on particle number

$$\bigotimes_{n_r} |0\rangle$$

Information
Transfer on the mixing
of particle number



$$\sum_k \alpha_k |01001\dots 1\rangle \otimes |\varphi_k\rangle$$

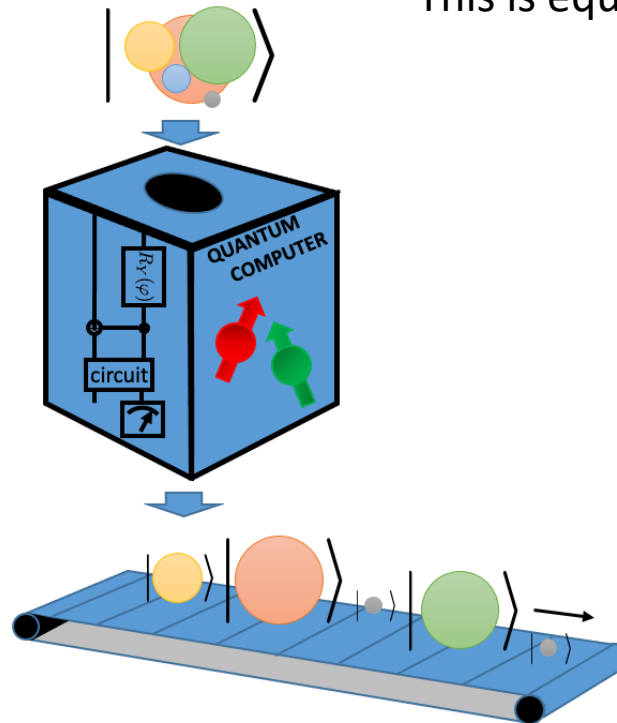
= Particle number
written as a binary number



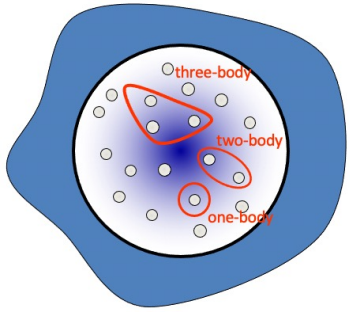
We can measure the register qubit
This is equivalent to project on $|\varphi_k\rangle$

$$|\Psi\rangle = \sum_k \alpha_k |\varphi_k\rangle$$

An even more schematic view



Then I can use this
Wave-function for
post-processing



Combining restoration of broken symmetries and variational method: The Quantum Variation After rojection (Q-VAP)

Coming back to our superconducting problem

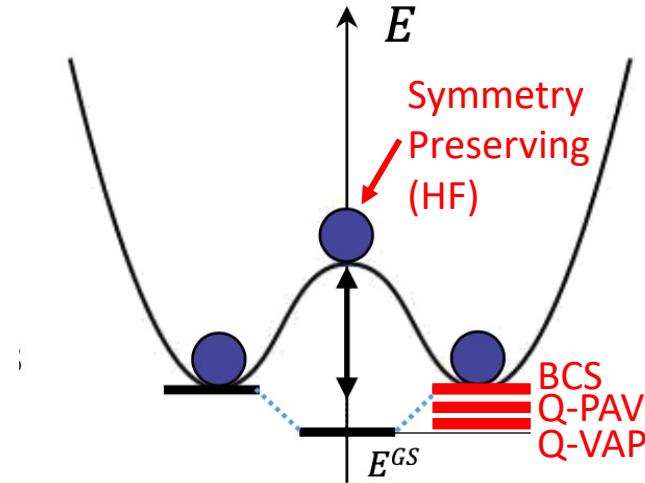
Combining projection with variational method

Possible optimization schemes

Variational

Symmetry-Breaking ansatz $|\Psi(\{\theta_p\})\rangle = \bigotimes_{p=1}^{N-1} [\sin(\theta_p)|0_p\rangle + \cos(\theta_p)|1_p\rangle]$

Pair occupation are now encoded

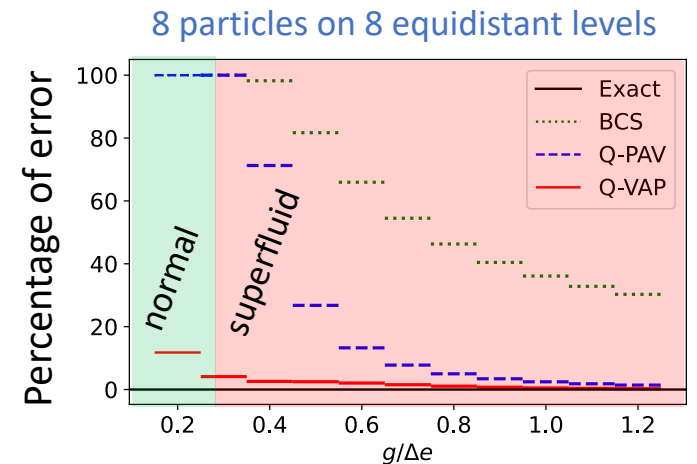
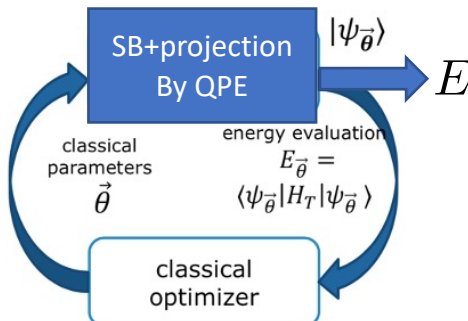
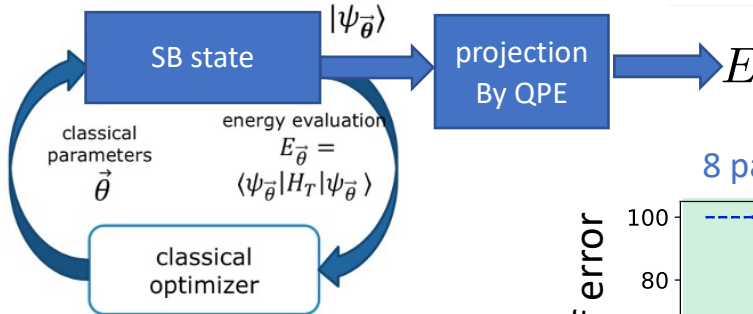
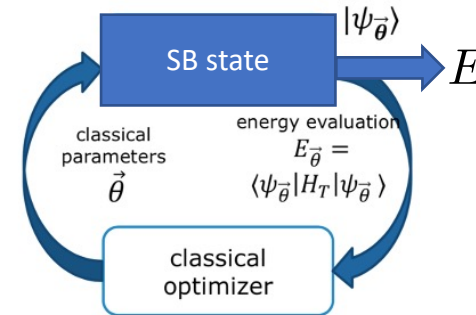


Quantum-Classical optimizers

➔ Standard BCS theory

➔ Project after optimization
Q-PAV: Quantum Projection After Variation

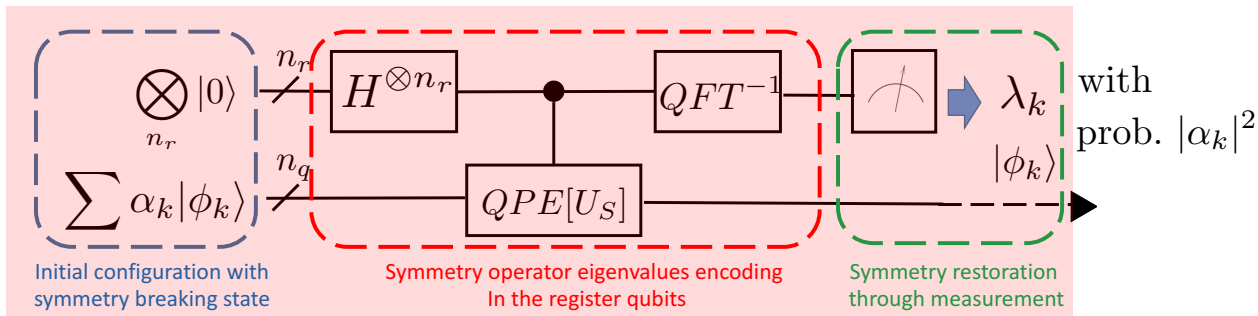
➔ The optimization is made on the Symmetry restored state.
Q-VAP: Quantum Variation After Projection



Complete strategy



E. A. Ruiz Guzman



Initial state preparation
(HF, BCS, P-VAP, Q-VAP)

➔ Post-processing to get excited states:
Quantum Phase Estimation

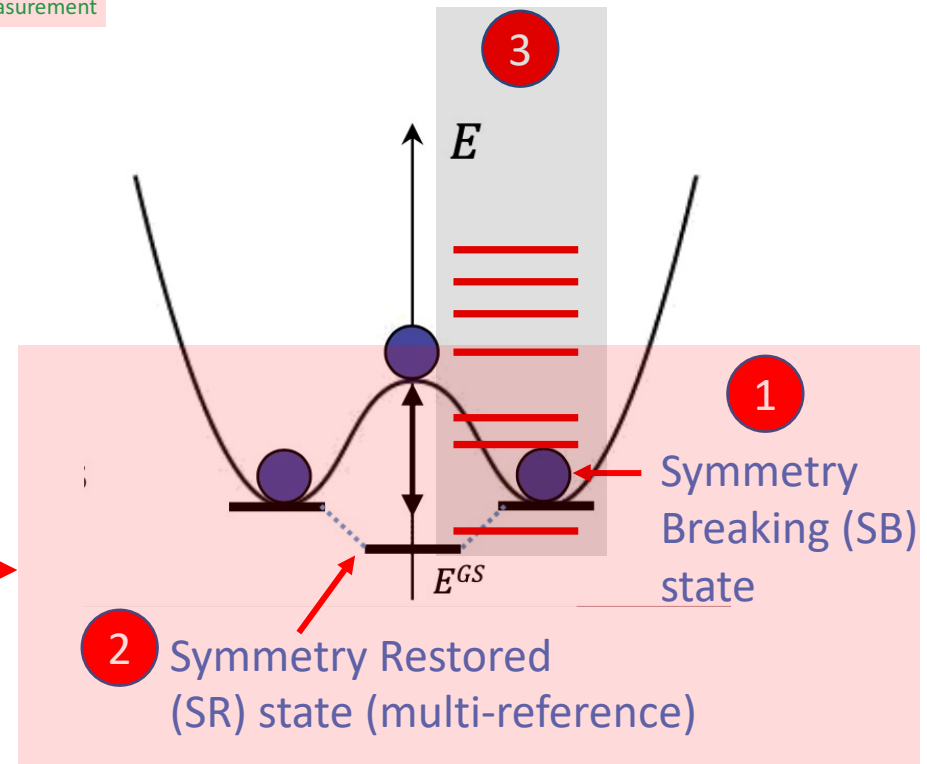
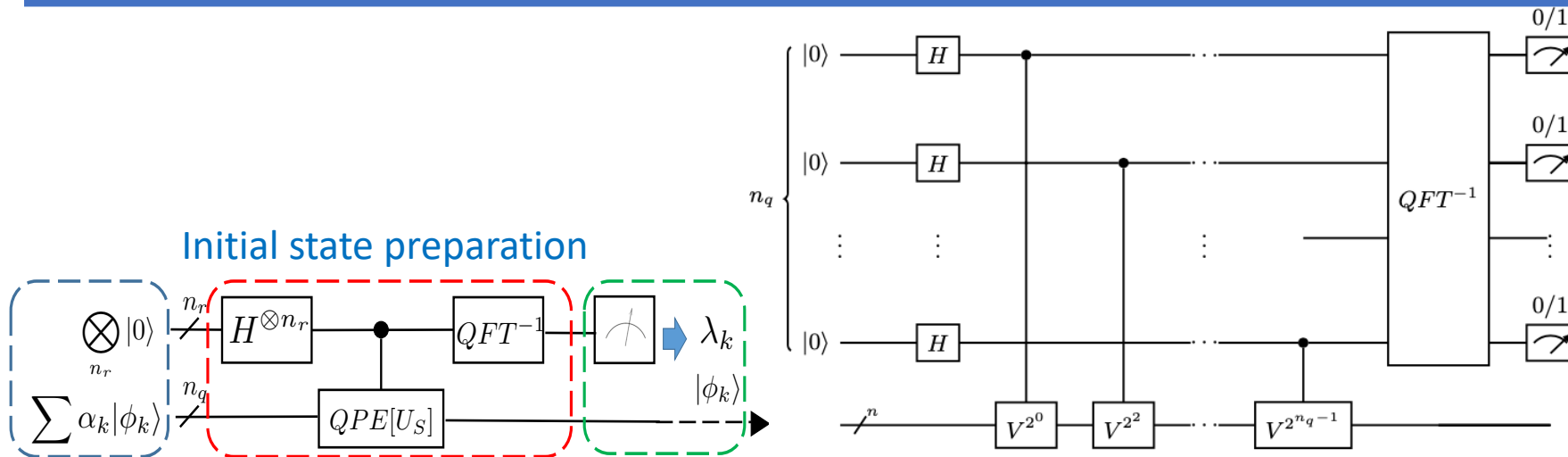


Illustration of the QPE method with projected state



Some technical details

$$V = \exp \left\{ -2\pi i \left(\frac{H - E_{\min}}{E_{\max} - E_{\min}} \right) \right\}$$

➔ For the propagator, we used the Trotter-Suzuki method

$$U(\tau) = e^{-i\tau H}$$

$$U(\tau) = \prod U(\Delta\tau) \longrightarrow \prod U_\varepsilon(\Delta\tau) U_g(\Delta\tau)$$

$$H_P = \sum_{i>0} \varepsilon_i (a_i^\dagger a_i + a_i^\dagger a_{\bar{i}}) - g \sum_{i,j>0} a_i^\dagger a_i^\dagger a_j a_j$$

$$\prod_p \begin{pmatrix} 1 & 0 \\ 0 & \exp(-2i\tilde{\varepsilon}_p \Delta t) \end{pmatrix} \prod_{p>q} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & \cos(\lambda_{pq}) & i \sin(\lambda_{pq}) & 0 \\ 0 & i \sin(\lambda_{pq}) & \cos(\lambda_{pq}) & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

with
 $\lambda_{pq} = g\Delta t$

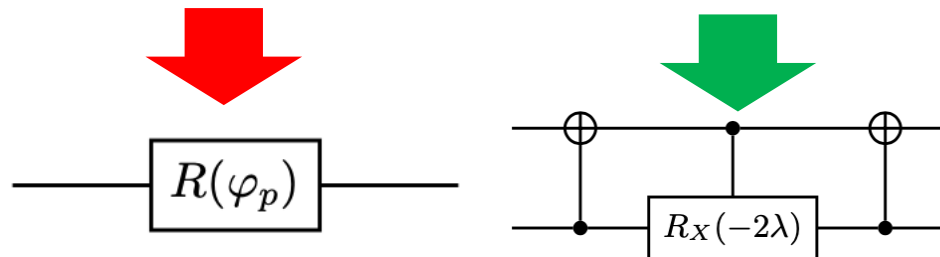
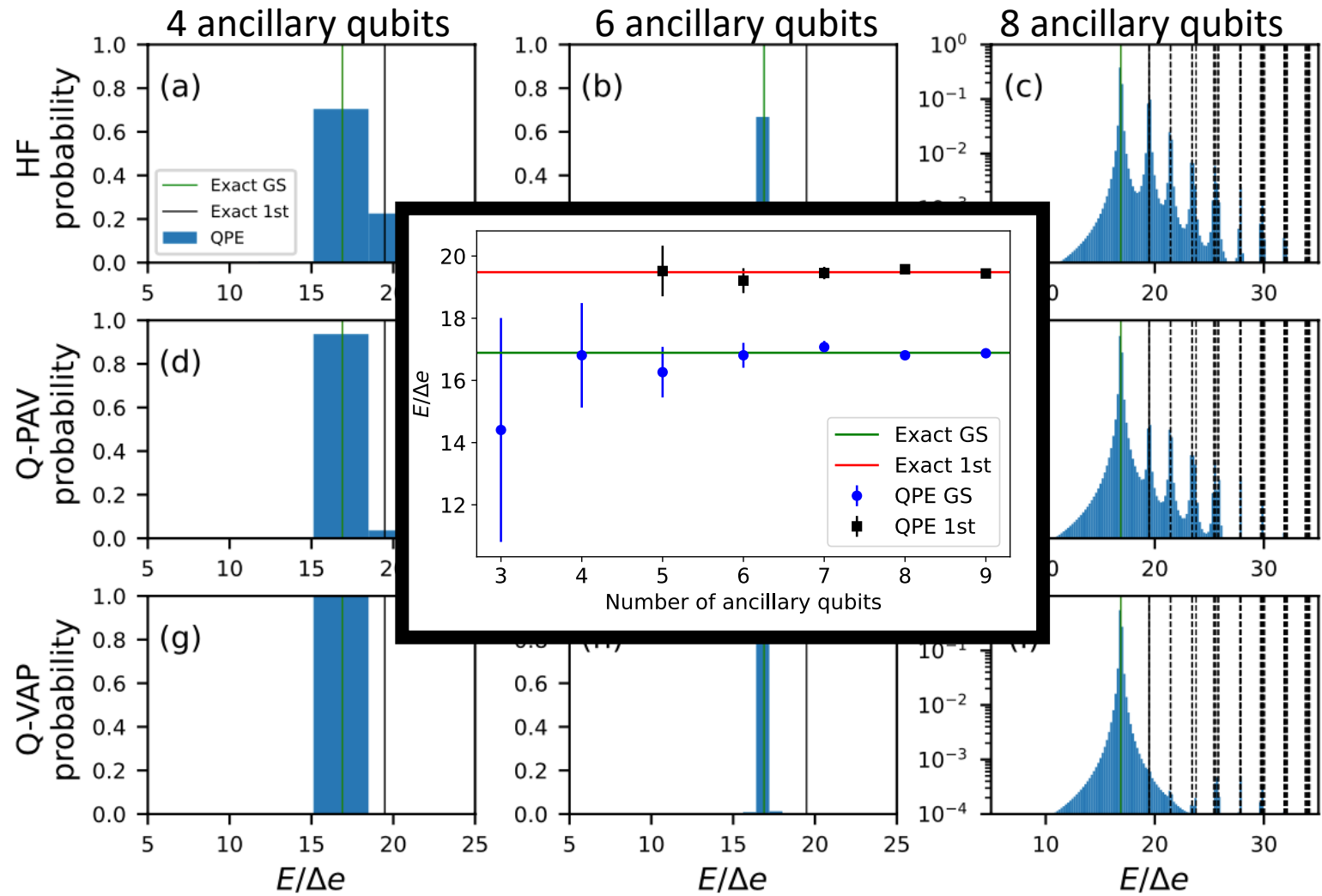
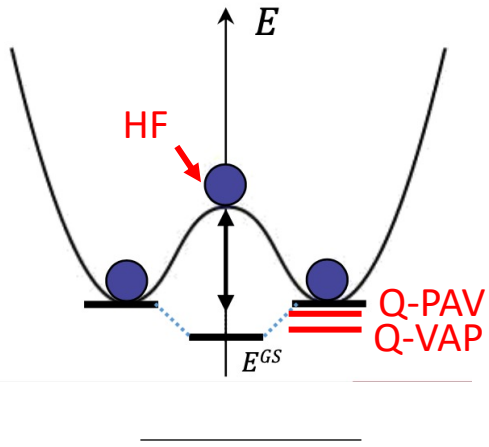
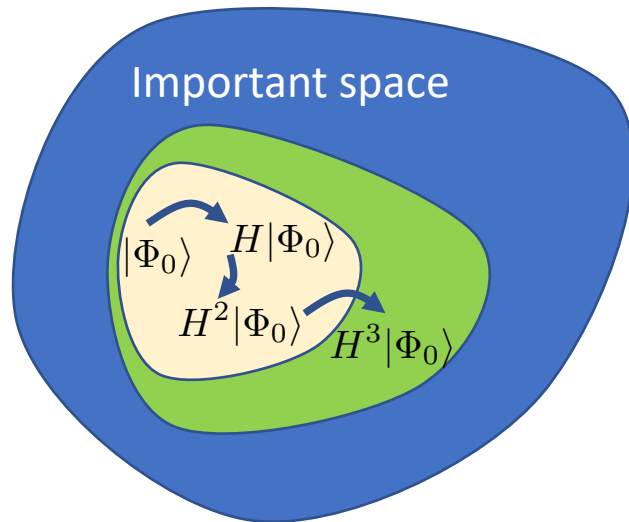


Illustration of the QPE method with projected state



Approximate method : Krylov Based methods

Hilbert space



Our strategy

Compute overlap and Hamiltonian matrix elements on the quantum computer



Solve the eigenvalue problem on the classical computer

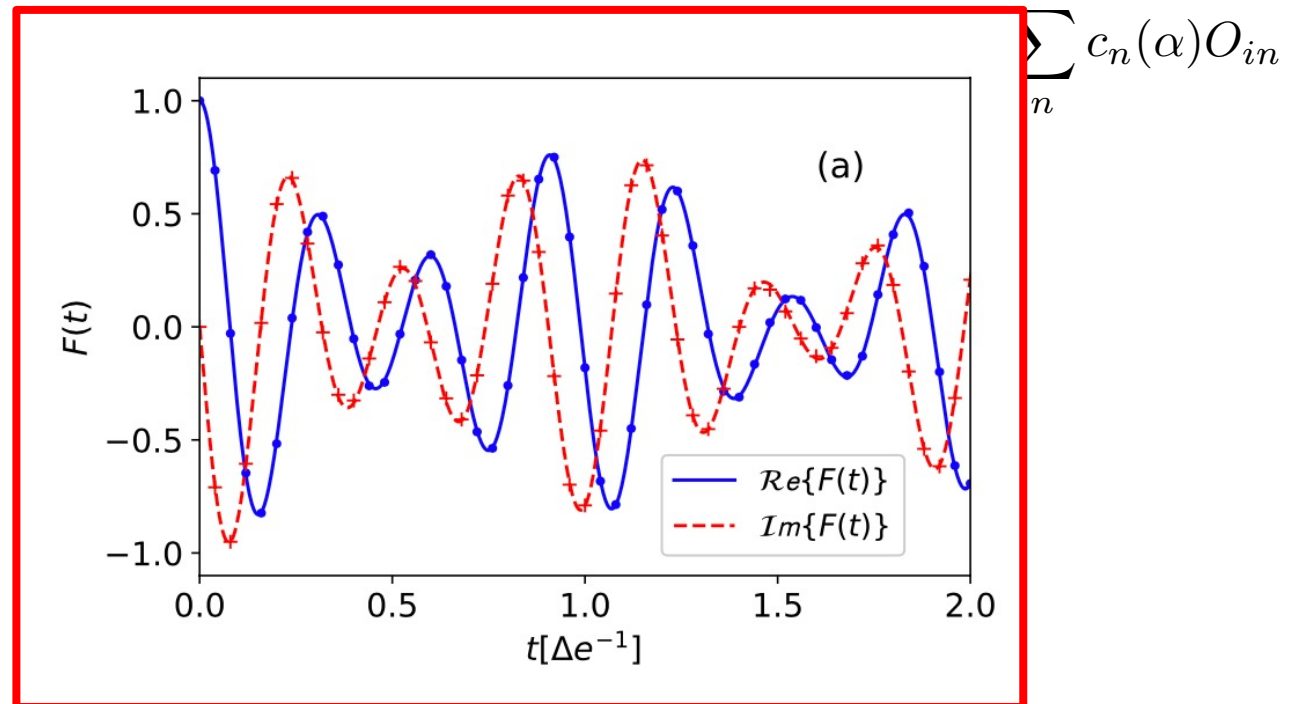
$$\{|\Psi\rangle, H|\Psi\rangle, \dots, H^{M-1}|\Psi\rangle\} \equiv \{|\Phi_i\rangle\}_{i=0, M-1}$$



Diagonalize in the non-orthogonal subspace

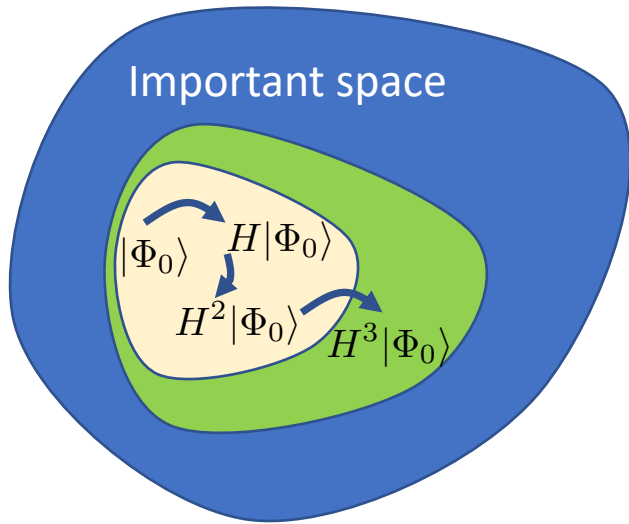
$$O_{ij} = \langle \Phi_i | \Phi_j \rangle \quad H_{ij} = \langle \Phi_i | H | \Phi_j \rangle$$

Generalized eigenvalue problem



Approximate method : Krylov Based methods

Highly Truncated Hilbert space



$$\{|\Psi\rangle, H|\Psi\rangle, \dots, H^{M-1}|\Psi\rangle\} \equiv \{|\Phi_i\rangle\}_{i=0, M-1}$$

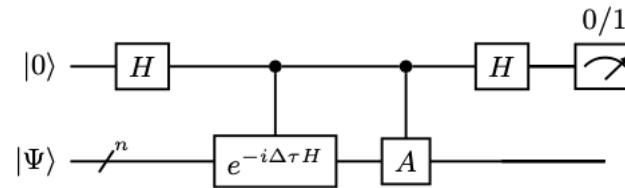


$$\{|\Psi\rangle, e^{-i\tau_1 H}|\Psi\rangle, \dots, e^{-i\tau_{M-1} H}|\Psi\rangle\}$$

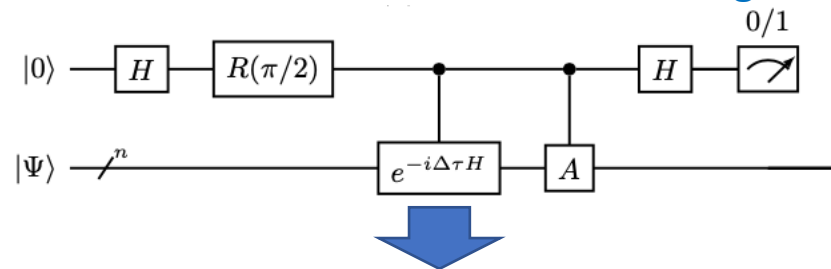


$$O_{ij} = \langle \Phi_i | \Phi_j \rangle = \langle \Psi | e^{-i(\tau_j - \tau_i) H} | \Psi \rangle \quad H_{ij} = \langle \Psi | H e^{-i(\tau_j - \tau_i) H} | \Psi \rangle$$

Hadamard test for the real part of O and H

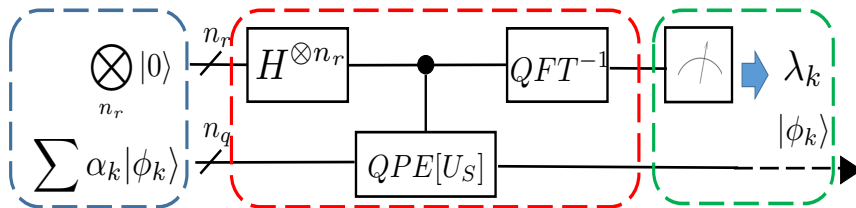


Modified Hadamard test for the imaginary part

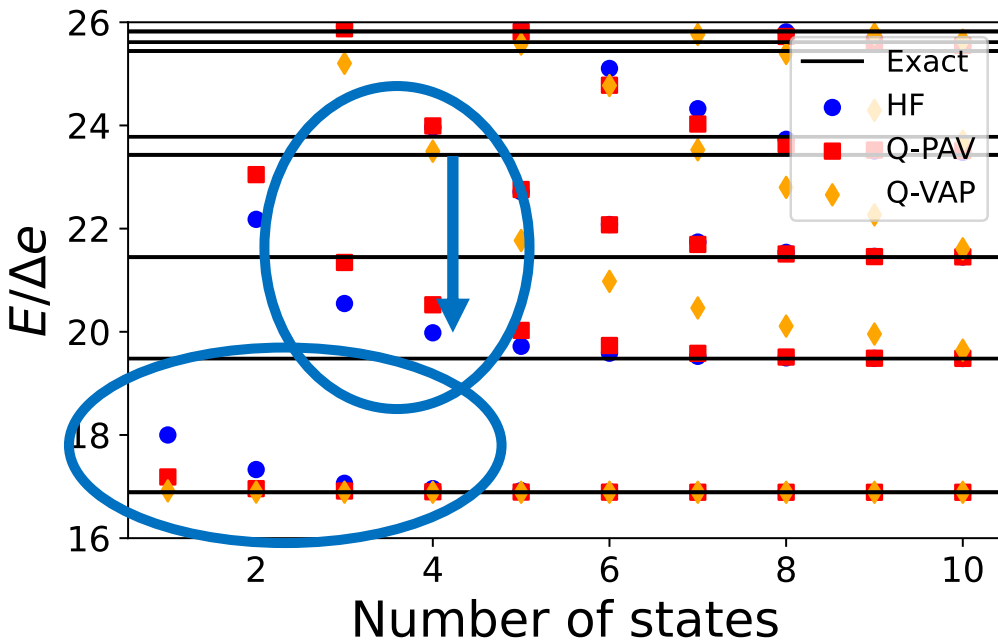


Diagonalization on a classical computer

Initial state preparation



Comparison QPE vs Quantum Kr



➔ The combination of Q-VAP + Quantum Krylov Is very good for the Ground state

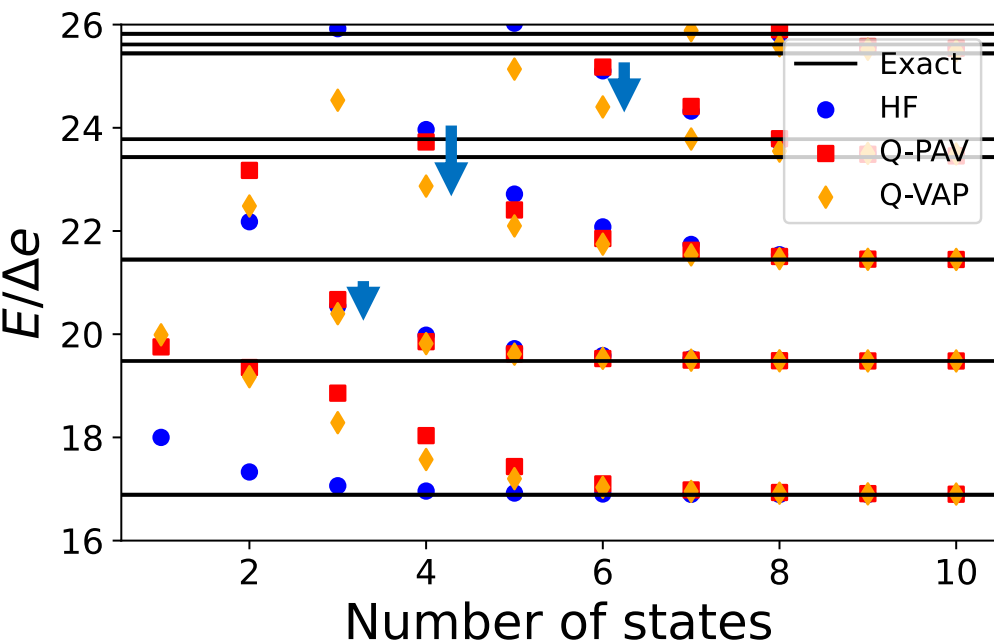
➔ But Q-VAP + Quantum Krylov is worth than others for excited states

A possible solution

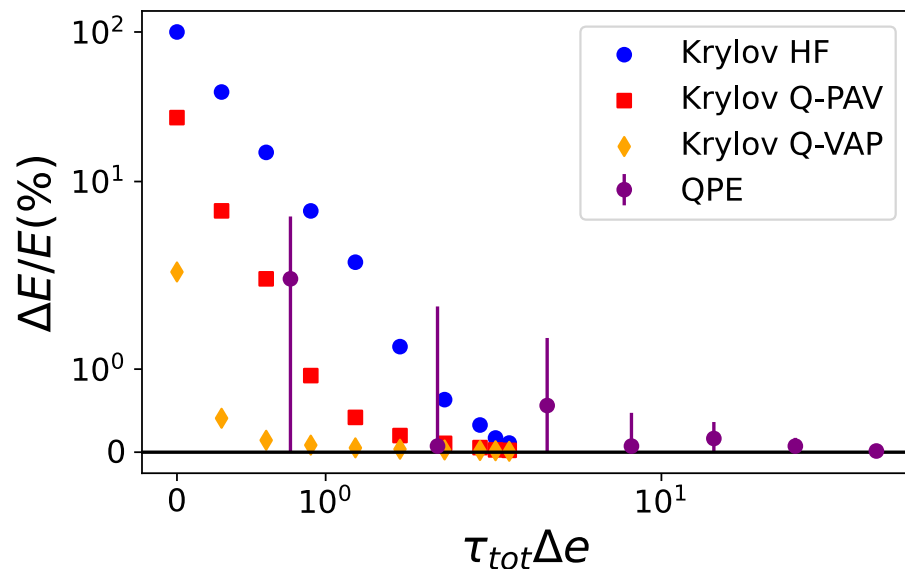
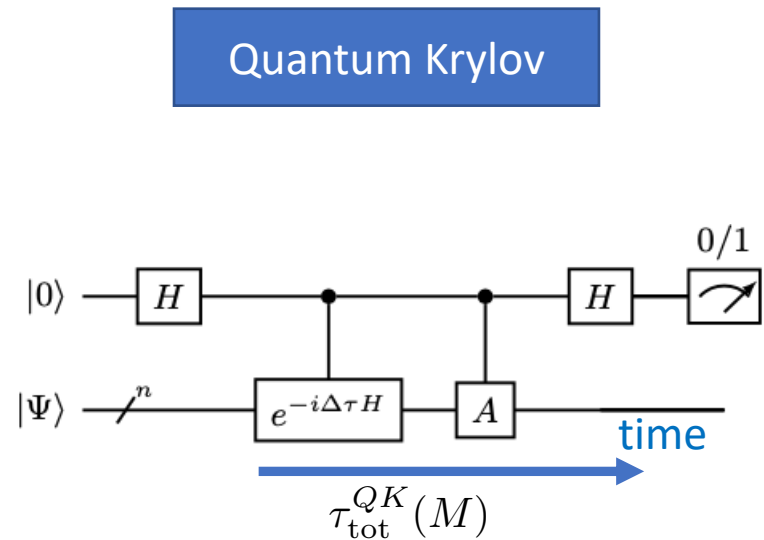
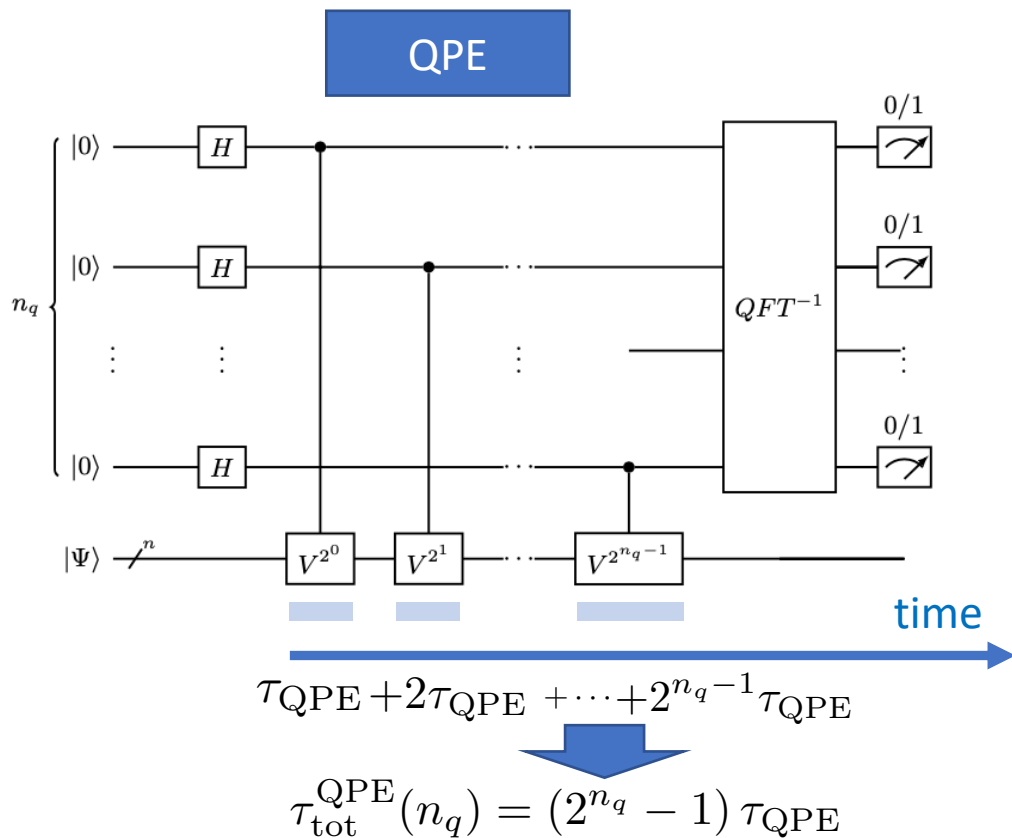
$$|\Psi\rangle = \bigotimes_p [\sin(\theta_p)|0_p\rangle + \cos(\theta_p)|1_p\rangle]$$

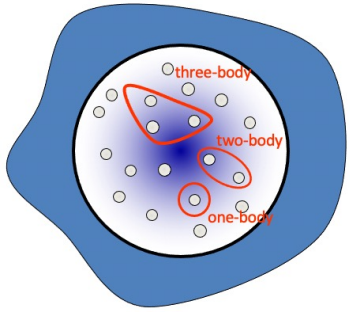
$$|\Psi'\rangle = [-\cos(\theta_i)|0_i\rangle + \sin(\theta_i)|1_i\rangle] \bigotimes_{p \neq i} [\sin(\theta_p)|0_p\rangle + \cos(\theta_p)|1_p\rangle]$$

$$\langle \Psi' | \Psi \rangle = 0$$



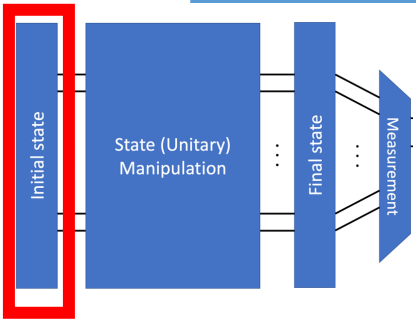
Comparison QPE vs Quantum Krylov after Q-VAP





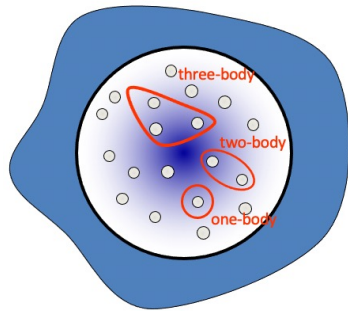
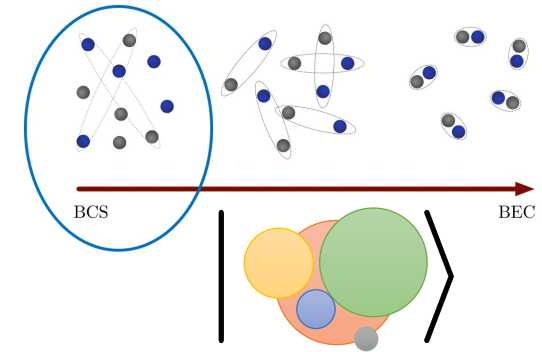
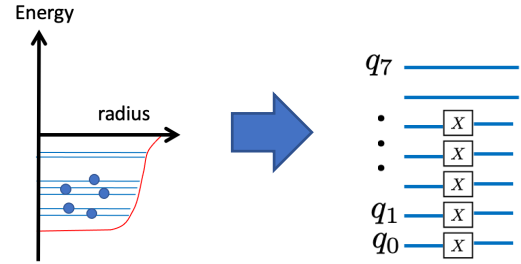
More on many-body ansatz in quantum computing

Some flash of what is happening now: nuclear and other many-body systems

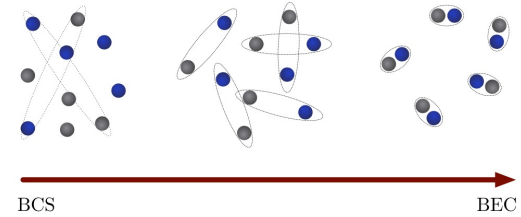


The many ways to prepare a system

Independent particles



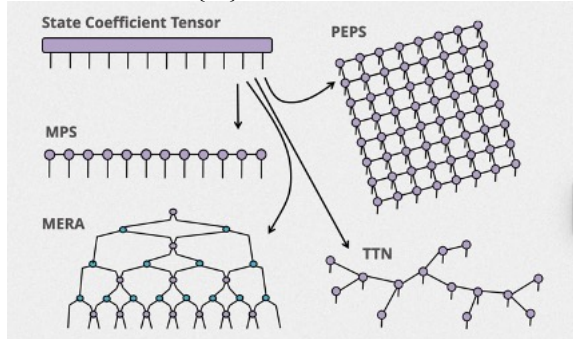
Superfluid systems



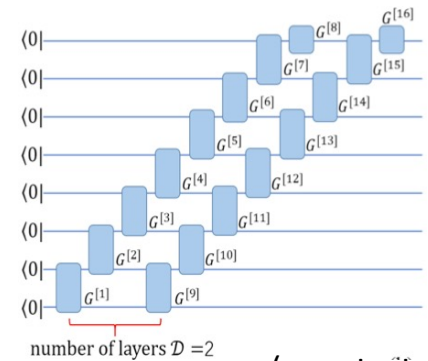
Requires symmetry breaking and restoration tools

Parametrizing general states: Tensor network

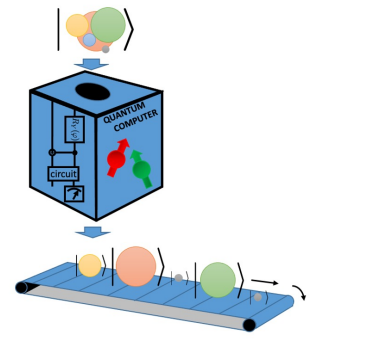
$$\Psi^{s_1, s_2, \dots, s_{n-1}, s_n} = \sum_{\{\alpha_i\}} G_{\alpha_1}^{s_1} G_{\alpha_1 \alpha_2}^{s_2} \dots G_{\alpha_{n-2} \alpha_{n-1}}^{s_{n-1}} G_{\alpha_{n-1}}^{s_n}$$



Examples:

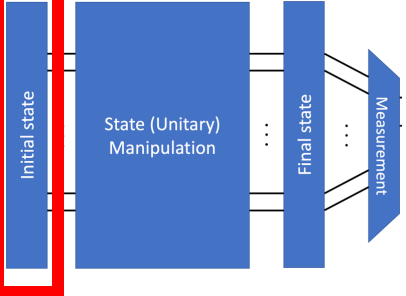


(now tested with A. Ruiz Guzman [PhD-IJCLab])



D. Lacroix, PRL 125, 230502 (2020).

Other ways to prepare correlated systems

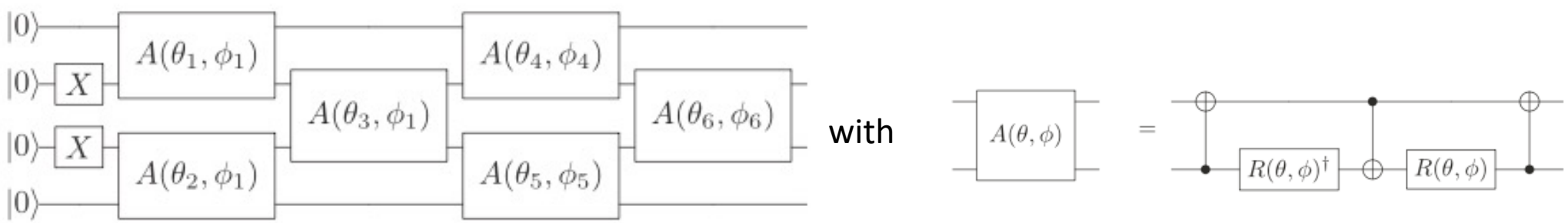


What is the most general way for 2 qubits ?

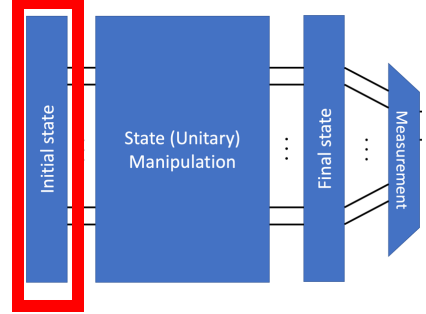
➔ In total, this gives a priori $4 + 6 \times 2 - 1 = 15$ parameters

$$\begin{array}{l}
 \alpha|00\rangle \\
 +\beta|01\rangle \\
 +\gamma|10\rangle \\
 +\delta|11\rangle
 \end{array}
 \begin{array}{c}
 \text{---} \\
 \text{---} \\
 \text{---} \\
 \text{---}
 \end{array}
 \begin{array}{c}
 \boxed{G(\{\lambda_i\})} \\
 \\
 \\
 \\
 \end{array}
 \begin{array}{l}
 \alpha'|00\rangle \\
 +\beta'|01\rangle \\
 +\gamma'|10\rangle \\
 +\delta'|11\rangle
 \end{array}$$

An example of structure that do not mixes particle number



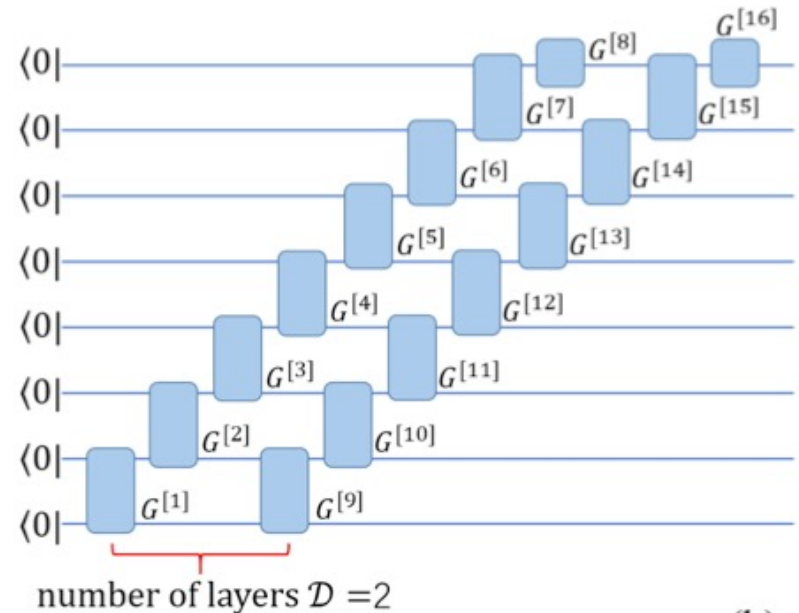
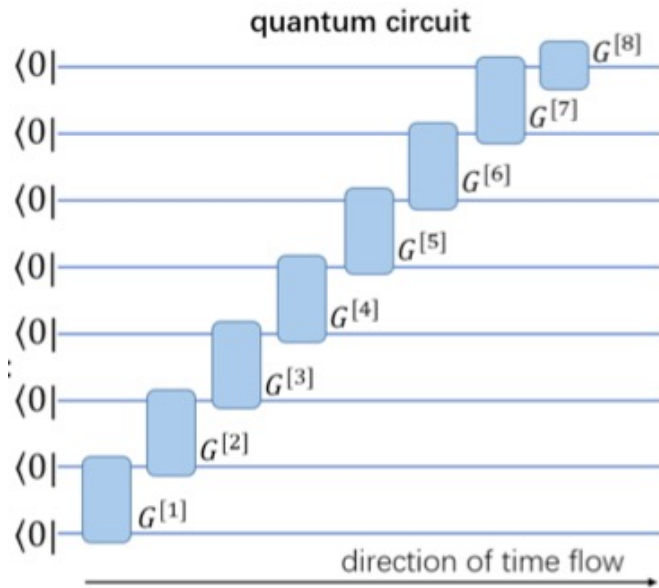
Other ways to prepare correlated systems



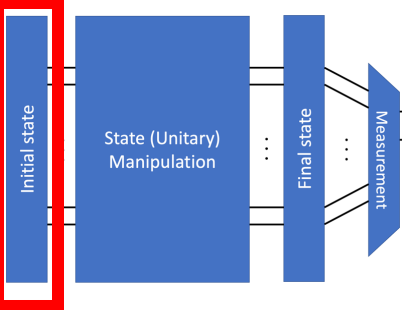
From Two to many qubits

We might use a “product like” form

$$|\Psi\rangle = \sum_{s_i=0,1} \Psi_{s_1, \dots, s_N} |s_1, \dots, s_N\rangle \longrightarrow \Psi^{s_1, s_2, \dots, s_{n-1}, s_n} = \sum_{\{\alpha_i\}} G_{\alpha_1}^{s_1} G_{\alpha_1 \alpha_2}^{s_2} \dots G_{\alpha_{n-2} \alpha_{n-1}}^{s_{n-1}} G_{\alpha_{n-1}}^{s_n}$$



Other ways to prepare correlated systems



$$\Psi^{s_1, s_2, \dots, s_{n-1}, s_n} = \sum_{\{\alpha_i\}} G_{\alpha_1}^{s_1} G_{\alpha_1 \alpha_2}^{s_2} \cdots G_{\alpha_{n-2} \alpha_{n-1}}^{s_{n-1}} G_{\alpha_{n-1}}^{s_n}$$

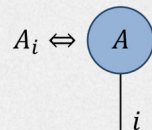
A few words on tensor network

<https://www.tensors.net/>
<https://tensornetwork.org/>

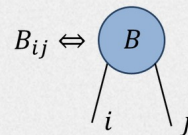
For our purposes, a tensor can simply be understood as a multi-dimensional array of numbers. Typically we use a diagrammatic notation for tensors, where each tensor is drawn as a solid shape with a number of 'legs' corresponding to its order:

$$A = \begin{bmatrix} A_1 \\ A_2 \\ \vdots \\ A_m \end{bmatrix} \quad B = \begin{bmatrix} B_{11} & \cdots & B_{1n} \\ \vdots & \ddots & \vdots \\ B_{m1} & \cdots & B_{mn} \end{bmatrix} \quad C = \begin{bmatrix} \begin{bmatrix} C_{111} & \cdots & C_{1n1} \\ \vdots & \ddots & \vdots \\ C_{m11} & \cdots & C_{mn1} \end{bmatrix}^1 \\ \vdots \\ \begin{bmatrix} C_{111} & \cdots & C_{1n1} \\ \vdots & \ddots & \vdots \\ C_{m11} & \cdots & C_{mn1} \end{bmatrix}^3 \end{bmatrix}^l$$

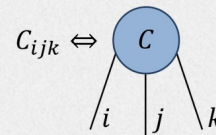
• vector



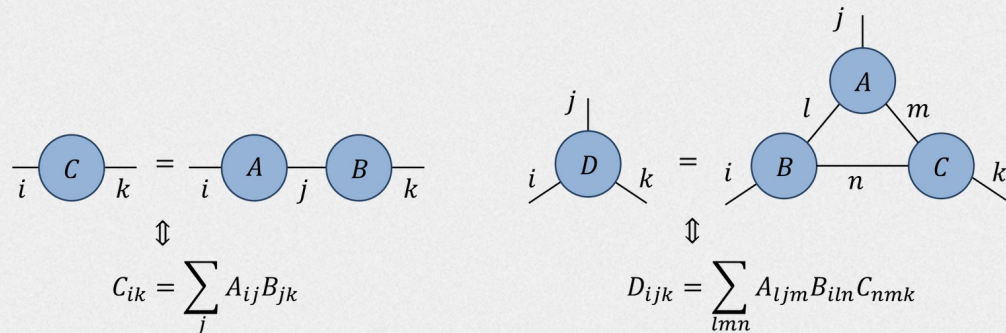
• matrix



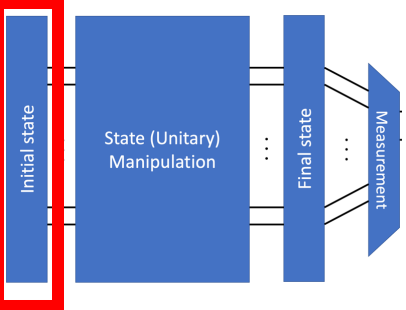
• order-3 tensor



We can form networks comprised of multiple tensors, where an index shared by two tensors denotes a contraction (or summation) over this index:



Other ways to prepare correlated systems

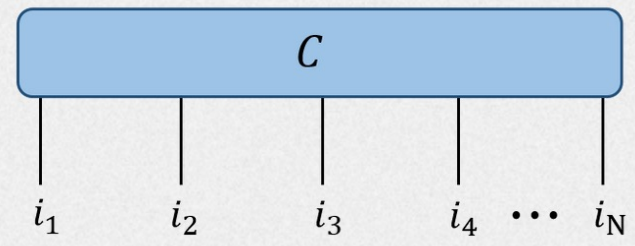


$$\Psi^{s_1, s_2, \dots, s_{n-1}, s_n} = \sum_{\{\alpha_i\}} G_{\alpha_1}^{s_1} G_{\alpha_1 \alpha_2}^{s_2} \cdots G_{\alpha_{n-2} \alpha_{n-1}}^{s_{n-1}} G_{\alpha_{n-1}}^{s_n}$$

A few words on tensor network

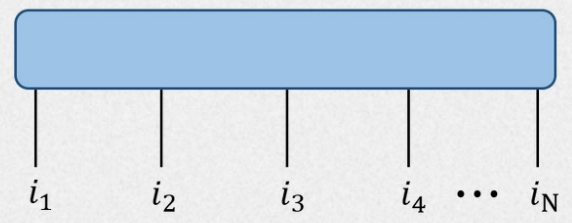
<https://www.tensors.net/>
<https://tensornetwork.org/>

$$|\psi\rangle = \sum_{i_1, i_2, \dots, i_N} C_{i_1, i_2, \dots, i_N} |i_1\rangle |i_2\rangle \cdots |i_N\rangle$$



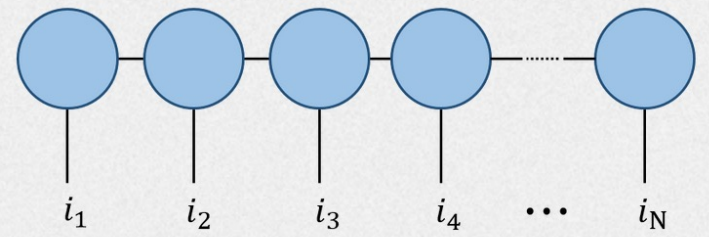
Matrix Product State

Order-N tensor:
contains $\sim \exp(N)$ parameters

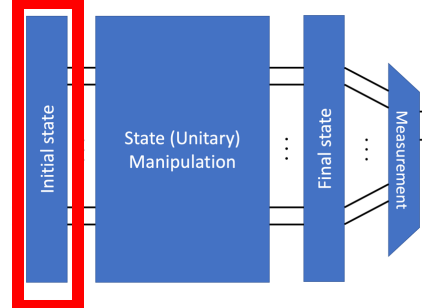


\Rightarrow

Network of low-order tensors:
contains $\sim \text{poly}(N)$ parameters



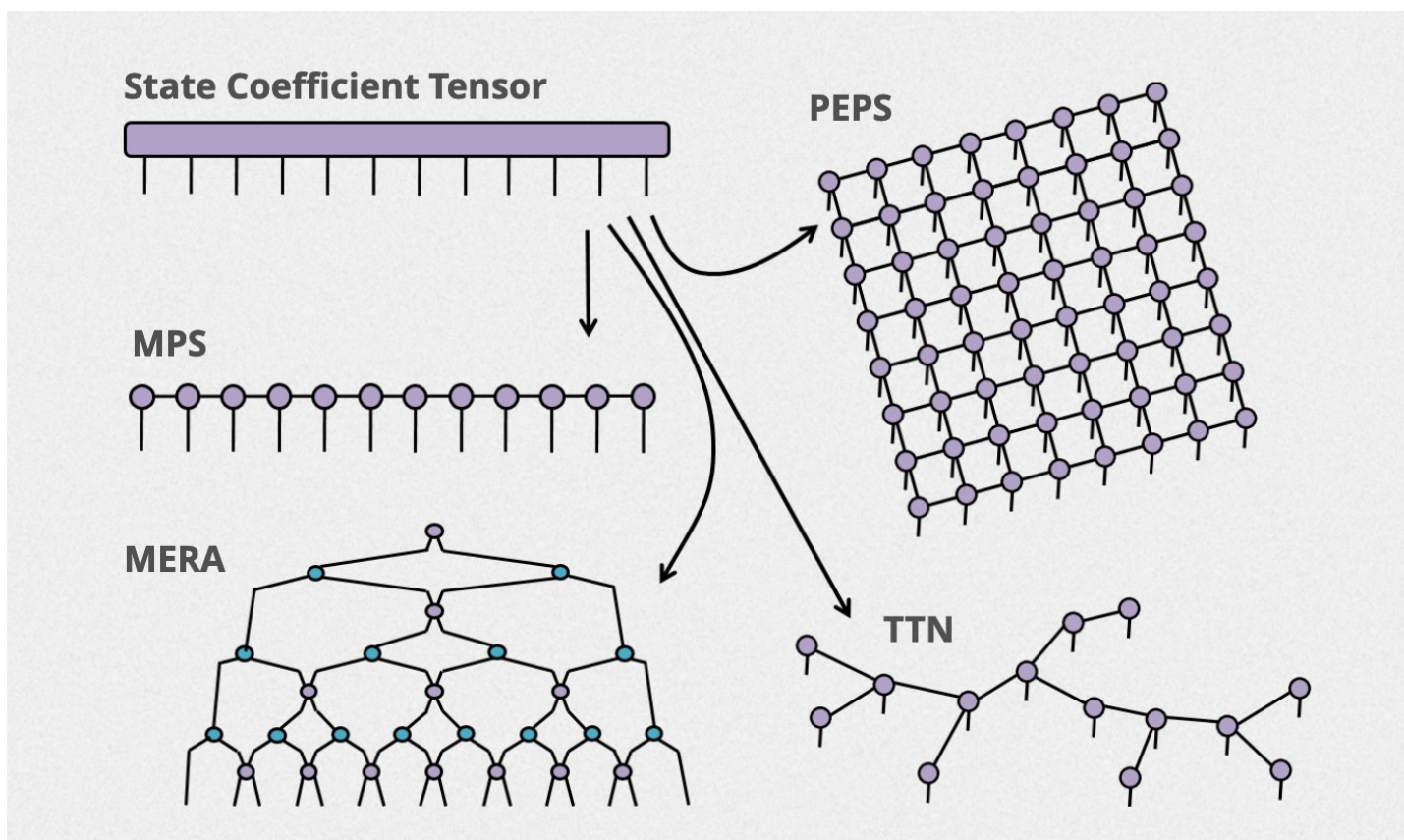
Other ways to prepare correlated systems

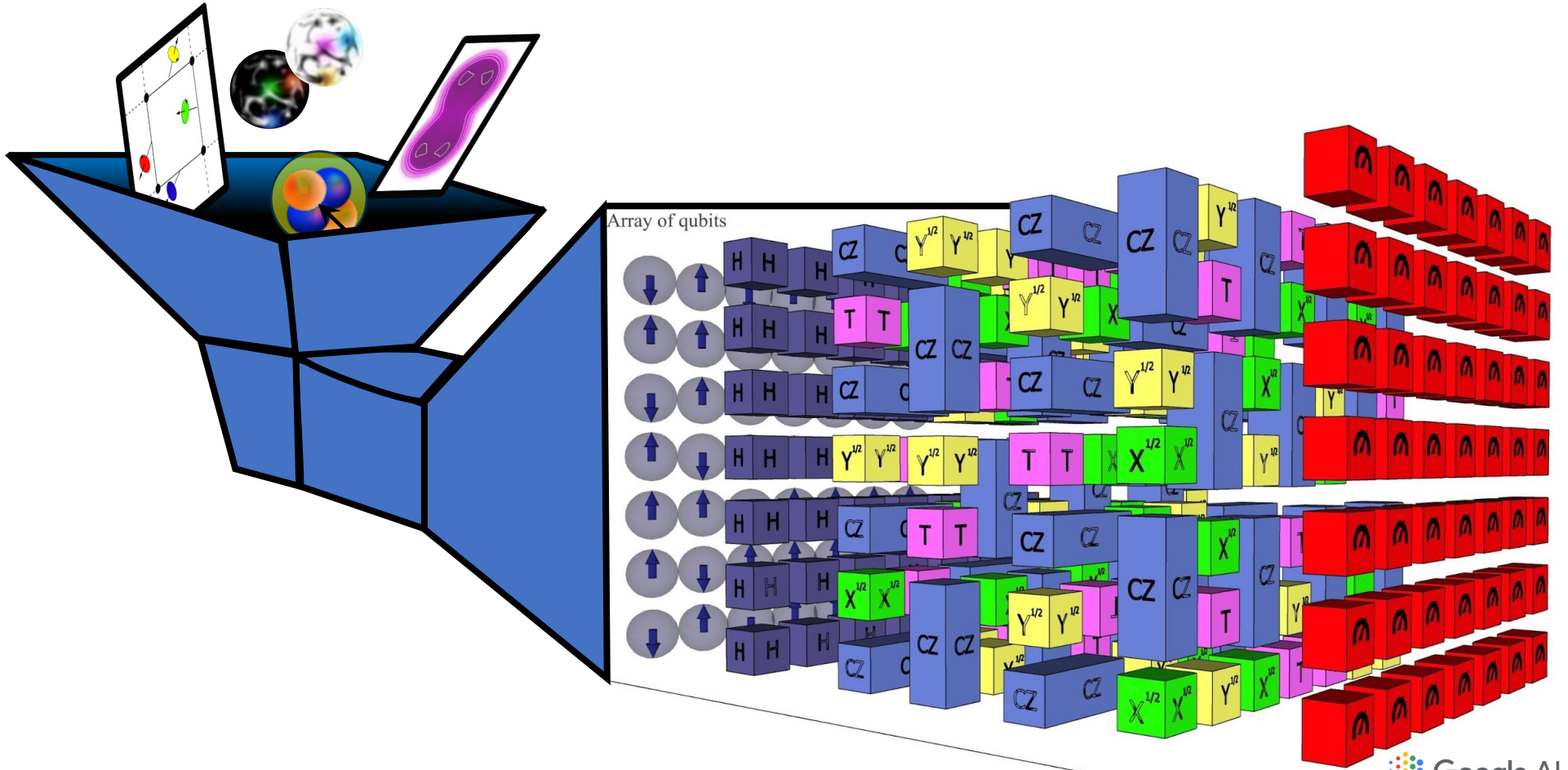


$$\Psi^{s_1, s_2, \dots, s_{n-1}, s_n} = \sum_{\{\alpha_i\}} G_{\alpha_1}^{s_1} G_{\alpha_1 \alpha_2}^{s_2} \dots G_{\alpha_{n-2} \alpha_{n-1}}^{s_{n-1}} G_{\alpha_{n-1}}^{s_n}$$

A few words on tensor network

<https://www.tensors.net/>
<https://tensornetwork.org/>





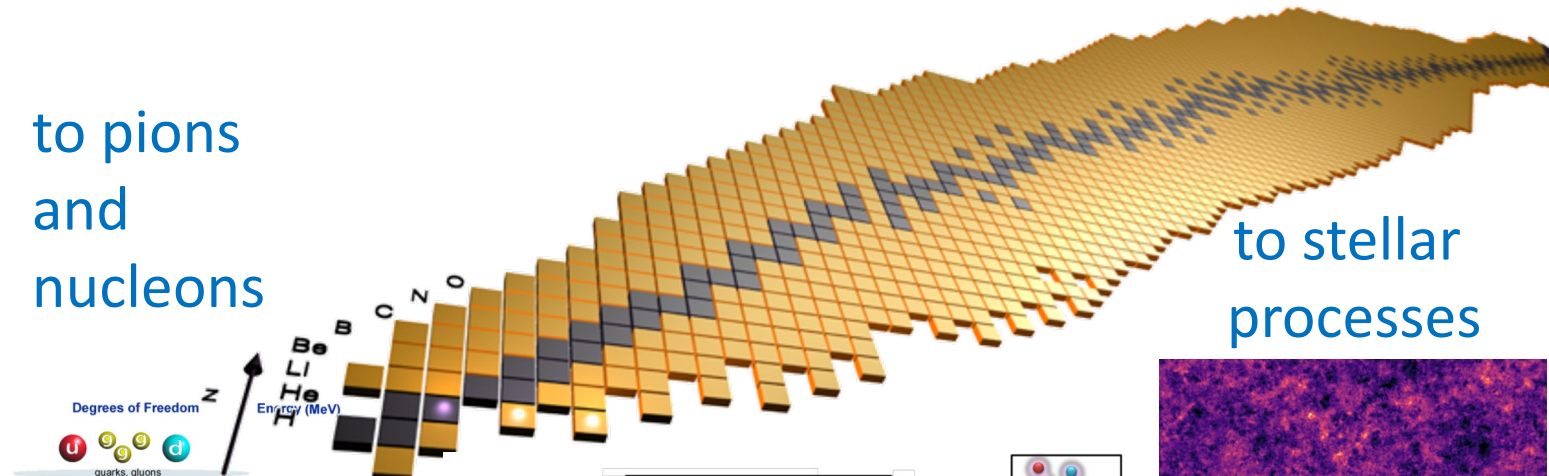
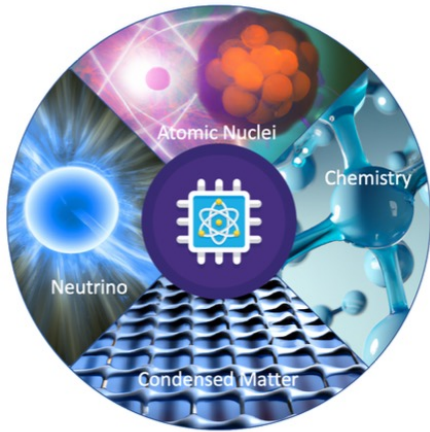
Quantum computers are often themselves many-body interacting systems

Quantum computing for the infinities

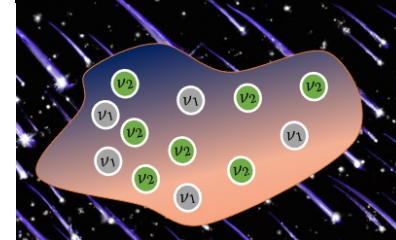
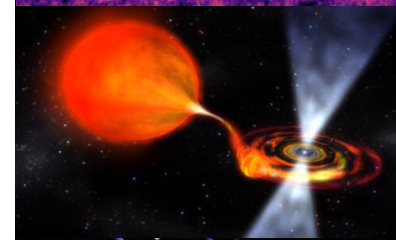
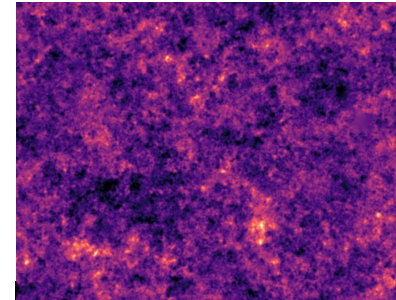
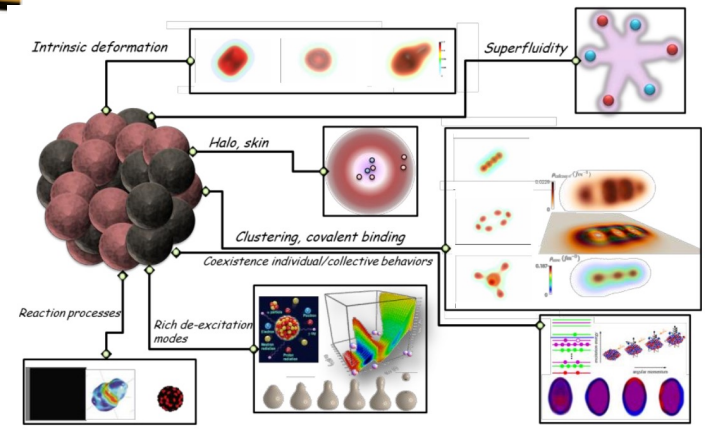
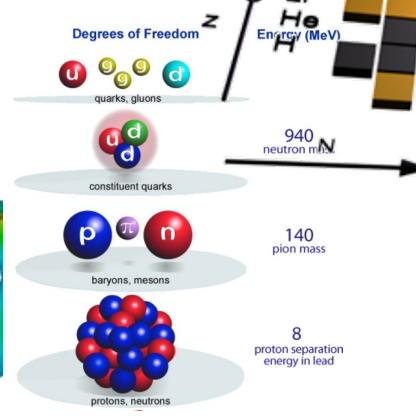
to atomic nuclei

to pions
and
nucleons

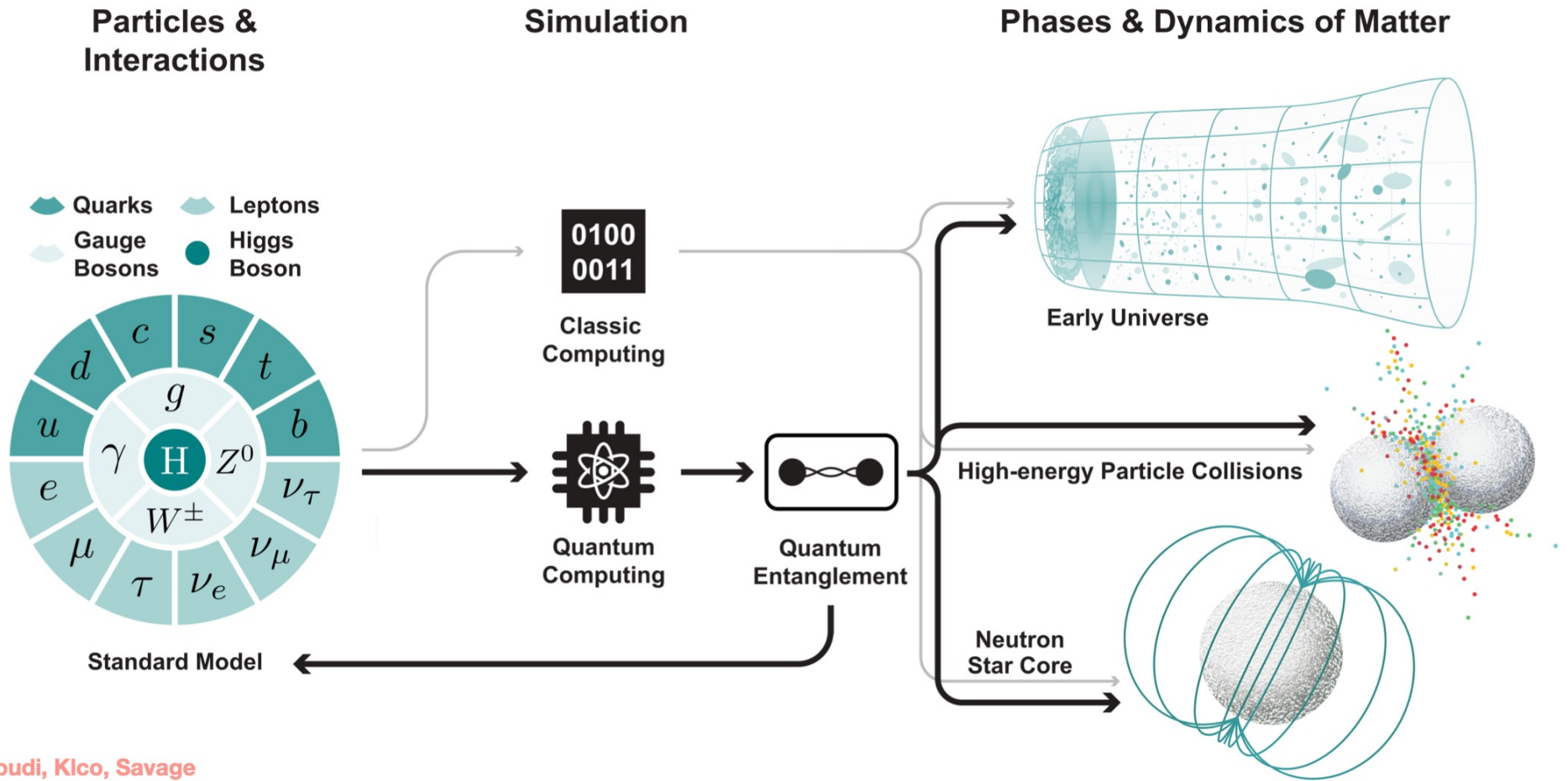
to stellar
processes



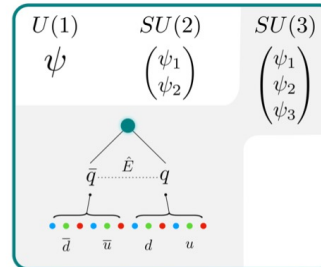
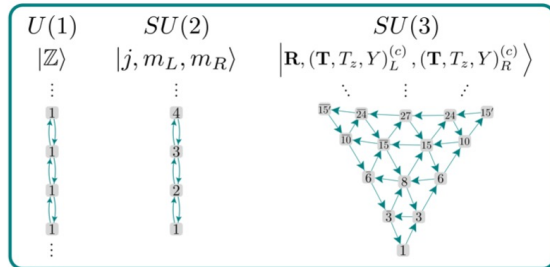
from quark



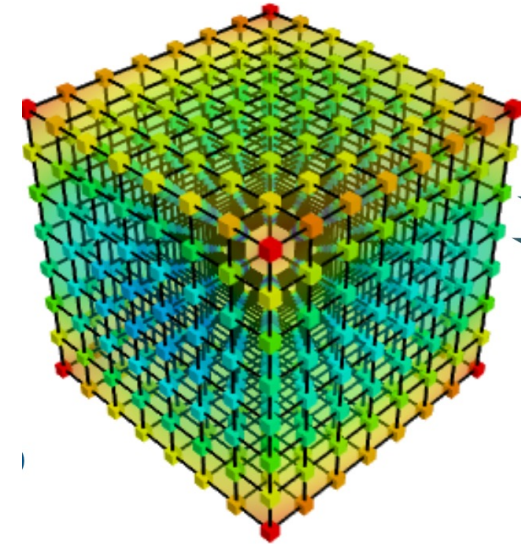
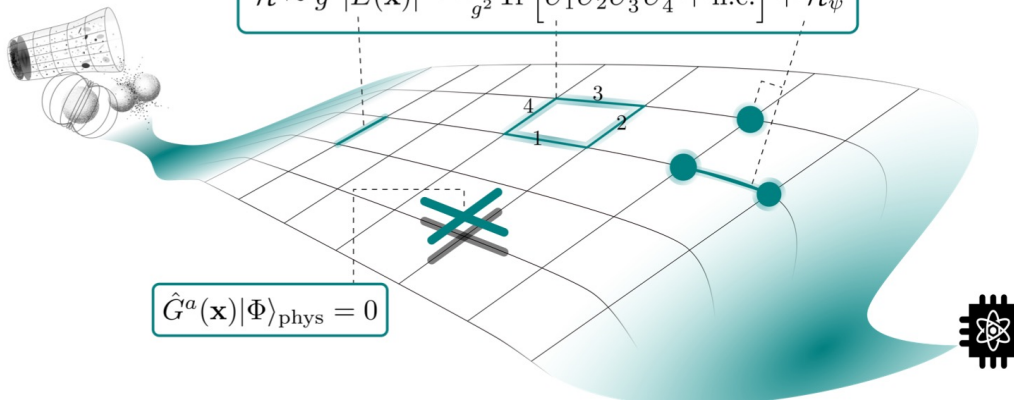
Courtesy J.P. Ebran



Digital Quantum Chromodynamics



$$\hat{\mathcal{H}} \sim g^2 |\hat{E}(\mathbf{x})|^2 - \frac{1}{g^2} \text{Tr} [\hat{U}_1 \hat{U}_2 \hat{U}_3^\dagger \hat{U}_4^\dagger + \text{h.c.}] + \hat{\mathcal{H}}_\psi$$



- ➔ Map quarks and gluons on quantum register
- ➔ Develop unitary operators for their evolution
- ➔ Obtain relevant observables from measurements

Digital Quantum Chromodynamics

$$\begin{matrix} U(1) & SU(2) \\ |\mathbb{Z}\rangle & |j, m\rangle \\ \vdots & \vdots \end{matrix}$$

PRX QUANTUM 4, 027001 (2023)

Quantum Simulation for High-Energy Physics

Christian W. Bauer,^{1,*} Zohreh Davoudi,^{2,†} A. Baha Balantekin,³ Tanmoy P. Ghosh,⁴ Atcharya,⁴ Marcela Carena,^{5,6,7,8} Wibe A. de Jong,¹ Patrick Draper,⁹ Aida El-Khadra,⁹ Masanori Hanada,¹¹ Dmitri Kharzeev,^{12,13} Henry Lamm,⁵ Ying-Ying Li,¹⁴ Mikhail Lukin,¹⁸ Yannick Meurice,¹⁹ Christopher Monroe,^{20,21,22,23} Guido Pagano,²⁴ John Preskill,²⁵ Enrico Rinaldi,^{26,27,28} Alessandro Roggero,²⁹ Martin J. Savage,³³ Irfan Siddiqi,^{31,32,34} George Siopsis,³⁵ David V. Vukobratovic,³⁶ Yukari Yamauchi,² Kübra Yeter-Aydeniz,³⁸ and ...

- ¹ Physics Division, Lawrence Berkeley National Laboratory
- ² Department of Physics, Maryland Center for Fundamental Physics Simulation, University of Maryland, College Park
- ³ Department of Physics, University of Wisconsin-Madison
- ⁴ T-2, Los Alamos National Laboratory
- ⁵ Fermi National Accelerator Laboratory
- ⁶ Enrico Fermi Institute, University of Chicago
- ⁷ Kavli Institute for Cosmological Physics, Northwestern University
- ⁸ Department of Physics, University of Wisconsin-Madison
- ⁹ Department of Physics and Illinois Quantum Center, University of Illinois Urbana-Champaign
- ¹⁰ Department of Physics, University of Illinois Urbana-Champaign
- ¹¹ Department of Mathematics, University of Illinois Urbana-Champaign
- ¹² Center for Nuclear Theory, Brookhaven National Laboratory

[hep-lat] 1 Feb 2023

Review on Quantum Computing for Lattice Field Theory

Lena Funcke,^{a,b,*} Tobias Hartung,^c Karl Jansen^d and Stefan König^e

^aTransdisciplinary Research Area "Building Blocks of Matter and Fundamental Matter" and Helmholtz Institute for Radiation and Nuclear Physics (HIR) at Nußallee 14-16, 53115 Bonn, Germany

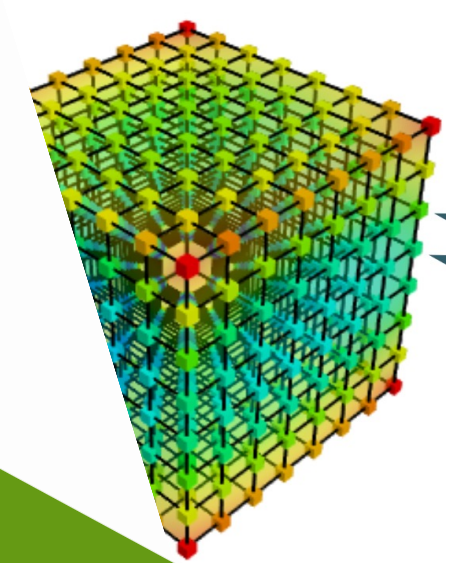
^bCenter for Theoretical Physics, Co-Design Center for Quantum Computing, Massachusetts Avenue, Cambridge, MA 02139

^cNortheastern University, Cambridge, MA 02125

^dDeutsches Elektronen-Synchrotron DESY, Notke 85, 22607 Hamburg, Germany

^eComputation-Based Science Center, 2121 Nicosia, Cyprus

E-mail: lena.funcke@desy.de, tobias.hartung@desy.de, karl.jansen@desy.de, stefan.koenig@desy.de



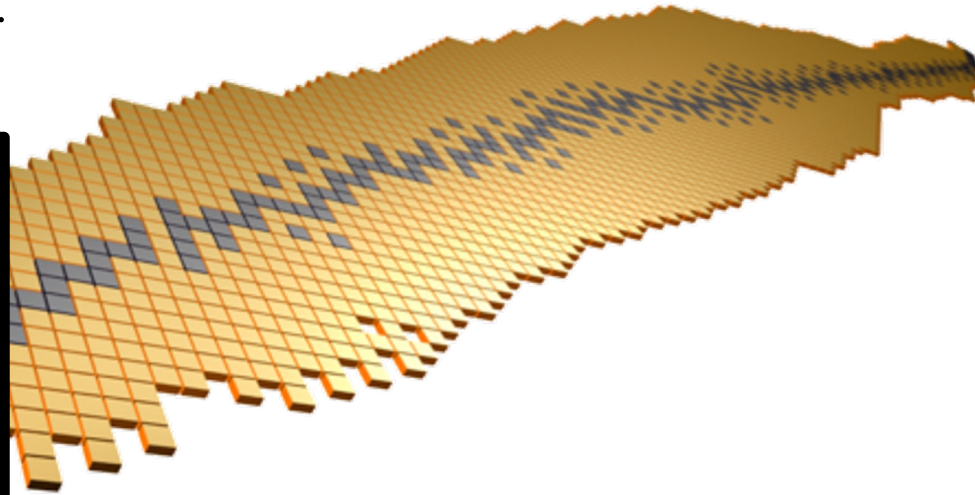
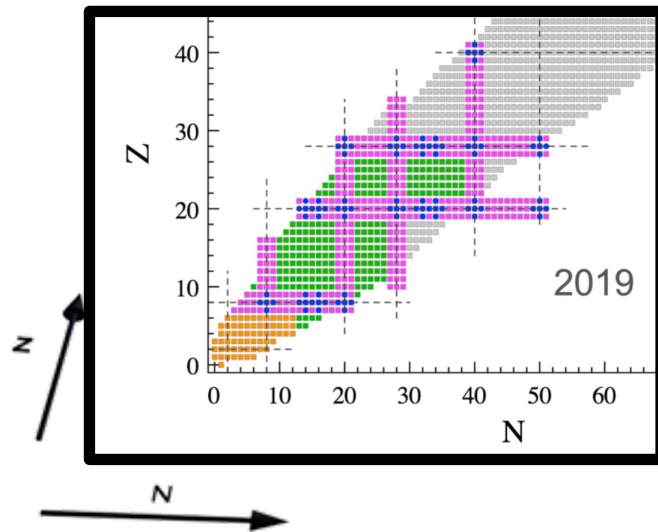
PROCEEDINGS OF SCIENCE

Roadmap

Actual tendency : Towards Full configuration-Interaction description ?

$$H = H_{1\text{-body}} + H_{2\text{-body}} + \dots$$

Current status

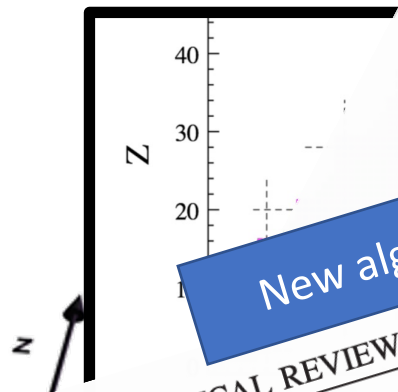


Actual tendency : Towards

Interaction description ?

$$H = H_{1\text{-body}} + H_{2\text{-body}}$$

Current s'



New algorithms
 PHYSICAL REVIEW LETTERS

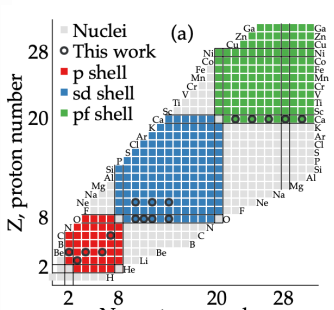
Rodeo Algorithm for Quantum
 Kenneth Choi¹, Dean Lee², Joey Bonitati², Zh...

¹Ridgefield High School, Ridgefield, Connecticut 06877, USA
²Facility for Rare Isotope Beams and Department of Physics and Astronomy, Michigan State University, East Lansing, Michigan 48824, USA
 PHYSICAL REVIEW LETTERS 125, 230502 (2020)

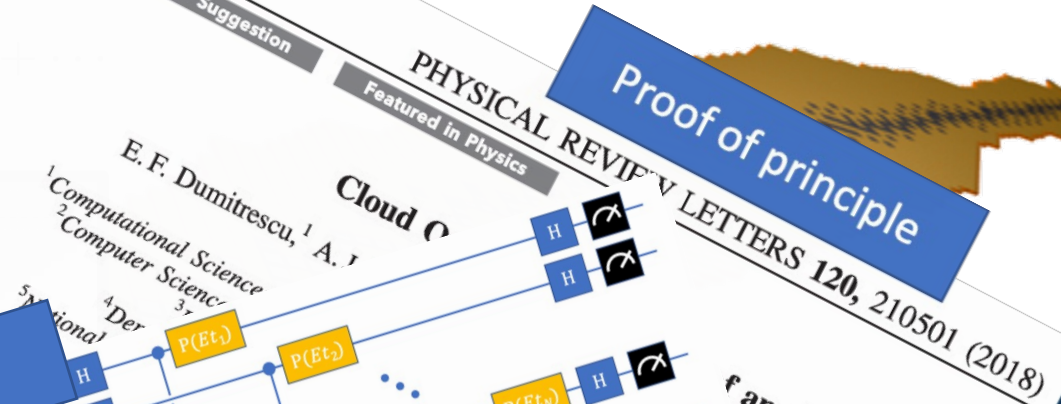
Systematic strategy

Nuclear shell-model simulation in digital quantum computers

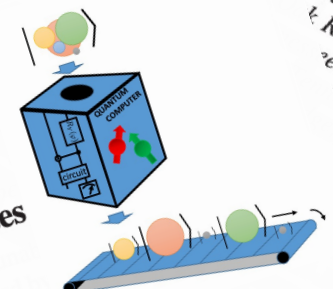
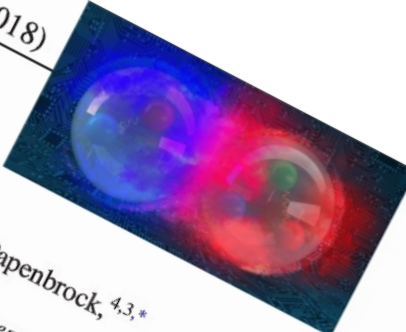
A. Pérez-Obiol*
 Barcelona Supercomputing Center, 08034 Barcelona, Spain

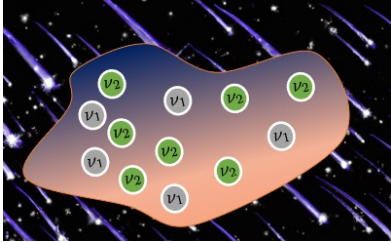


A. M. Romero[†], J. Menéndez[‡], and A. Rios[§]
 Departament de Física Quàntica i Astrofísica (FQA),
 Universitat de Barcelona (UB), c/ Martí i Franquès 1, 08008, Barcelona, Spain

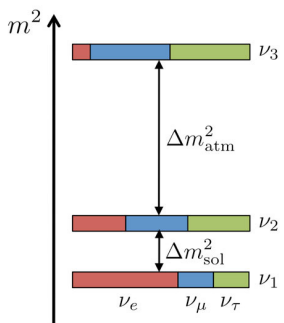


of an Atomic Nucleus
 T.D. Morris^{1,3}, T. Papenbrock^{4,3,*}
¹Wvski^{1,3}
²Oak Ridge, Tennessee 37831, USA
³37831, USA
⁴37996, USA
 Tennessee 37831, USA

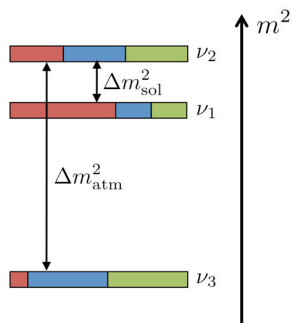




normal hierarchy (NH)



inverted hierarchy (IH)



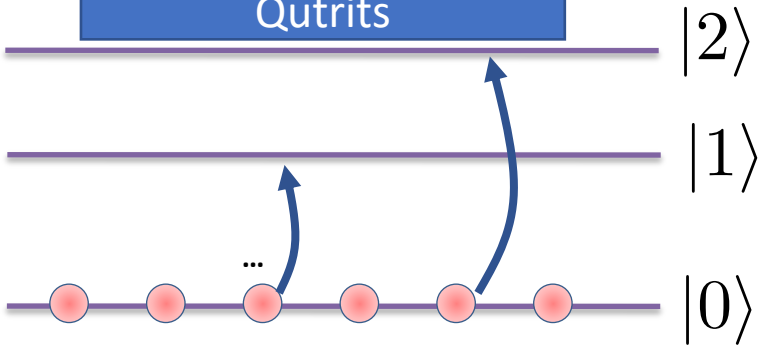
2-flavor approx.
directly treated as
Qubit



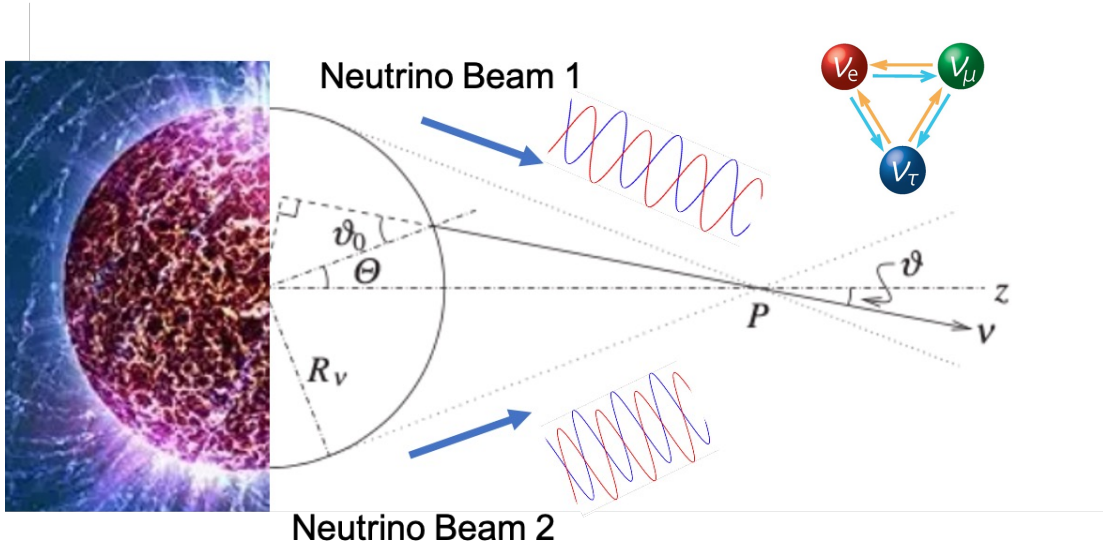
— $|1\rangle$

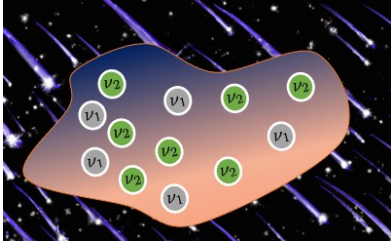
— $|0\rangle$

Natural treatment as
Qutrits



Neutrino oscillations in beams





Nowadays: Increasing number of applications

2-flavor approx. directly treated as

PHYSICAL REVIEW D **104**, 063009 (2021)

Simulation of collective neutrino oscillations on a quantum computer

Benjamin Hall,¹ Alessandro Roggero,^{2,3} Alessandro Baroni,⁴ and Joseph Carlson⁴

¹Facility for Rare Isotopes Beams (FRIB), Michigan State University, East Lansing, Michigan 48824, USA

PHYSICAL REVIEW D **107**, 023007 (2023)

Trapped-ion quantum simulation of collective neutrino oscillations

Valentina Amitrano,^{1,2,*} Alessandro Roggero,^{1,2} Piero Luchi,^{1,2} Francesco Luca Vespucci,^{1,2,3} and Francesco Pederiva^{1,2}

¹Dipartimento di Fisica, University of Trento, via Sommarive 14, I-38122 Borgo Tosseno, TN, Italy

PHYSICAL REVIEW LETTERS **130**, 221003 (2023)

Multi-Neutrino Entanglement and Correlations

Marc Illa^{1,*} and Martin

InQubator for Quantum Simulation (IQUS), Department of Physics, UI

(Received 7 December 2022; revised 27 April 2023; accepted ...)

The time evolution of multi-neutrino entanglement and correlations in neutrino oscillations, relevant for dense neutrino environments, has been simulated using quantum computers. Simulations performed of systems with up to 12 neutrinos using quantum computers are used to compute n -tangles, and two- and three-tangle mean-field descriptions. n -tangle rescalings are found to converge in the presence of genuine multi-neutrino entanglement.

DOI: 10.1103/PhysRevLett.130.221003

▪ $|1\rangle$

▪ $|0\rangle$

Quantum Information And entanglement

Eur. Phys. J. A manuscript No. (will be inserted by the editor)

Quantum information and quantum simulation of neutrino physics

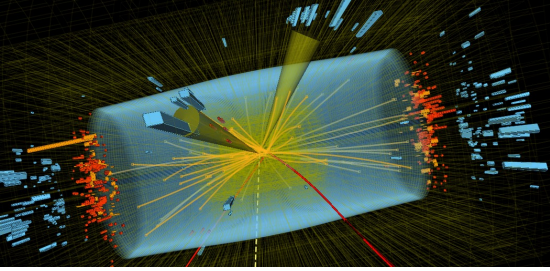
A. B. Balantekin^{a,1}, Pooja Siwach^{a,1}, Michael J. Cervia^{b,2,3}, Amol V. Patwardhan^{c,4}, Eermal Rrapaj^{d,5,6,7},

¹ University of Wisconsin, 1150 University Ave, Madison, WI 53706
² George Washington University, 725 21st St NW, Washington, DC 20052
³ University of Maryland, College Park, MD, USA 20742
⁴ SLAC National Accelerator Laboratory, 2575 Sand Hill Rd, Menlo Park, CA 94025
⁵ ERSC, Lawrence Berkeley National Laboratory, Berkeley, California 94720, USA
⁶ CERN, CH-1211, Geneva, Switzerland
⁷ UC Berkeley, California, Berkeley, CA 94720-7300

1 Neutrinos in extreme astrophysical environments

such as ... neutrinos, owing to their feeble interactions ... efficient at transporting energy ... extreme astrophysical ... binary comp

Quantum Machine Learning and event classification



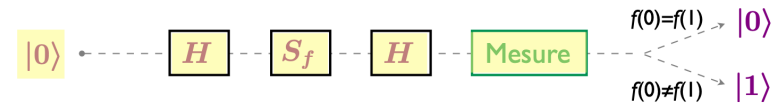
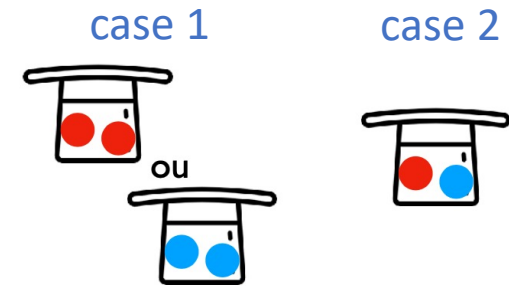
Some quantum historical algorithms are very fast for pattern recognition.

Deutsch (1985), Simon (1994), ...

The “simple” Deutsch problem: $f: \{0,1\} \rightarrow \{0,1\}$ (Oracle)

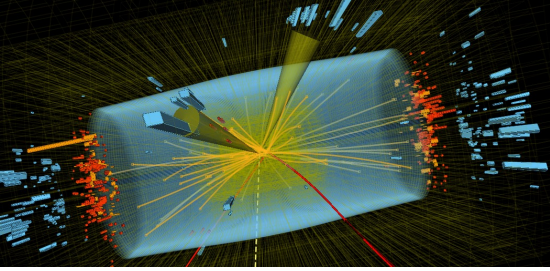
Q: determine if $f(0)=f(1)$

- Classically requires to have 2 answers $f(0) ? f(1) ?$
- Quantum: one can directly ask $f(0)=f(1)$



Quantum Machine Learning and event classification

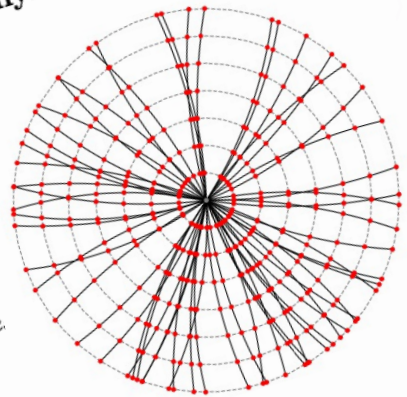
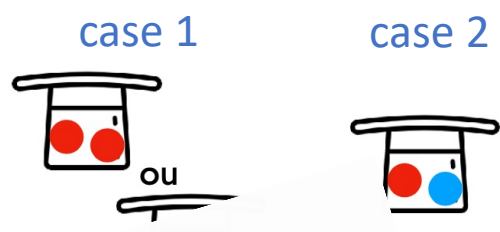
Some quantum historical algorithms are very fast for pattern recognition.



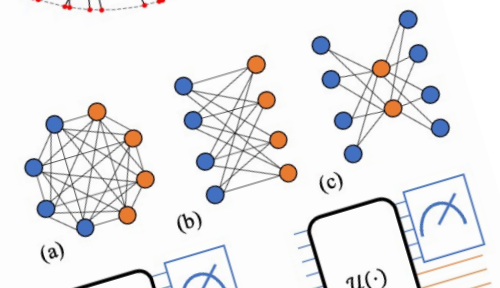
994), ...

1} (Oracle)

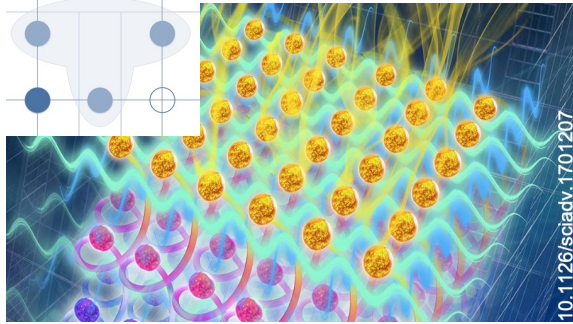
?



Quantum computing for data analysis in high energy physics*
Andrea Delgado,[†] Kathleen E. Hamilton, and Prasanna
Oak Ridge National Laboratory, Oak Ridge, Tennessee U
Jean-Roch Vimant[‡]
California Institute of Technology
Duarte Magano
Instituto Superior Técnico, Universidade de Lisboa, Portuga
Instituto de Telecomunicações, Physics of Information and Quantum Technologie
Yasser Omar
Instituto Superior Técnico, Universidade de Lisboa, Portu
Instituto de Telecomunicações, Physics of Information and Quantum Technologies (I
Portuguese Quantum Institute, Portugal
Pedrame Bargassa
Quantum Institute, Portugal and
Experimental de Partículas, Lis



Lattice gauge theories



Zohar, Kolck, Savage, ...

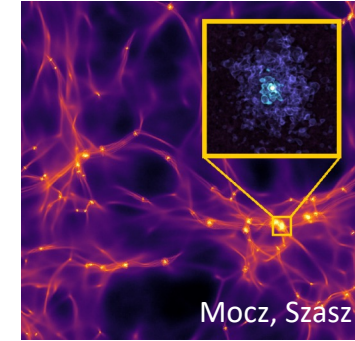
N-body problem

N-body nuclear systems



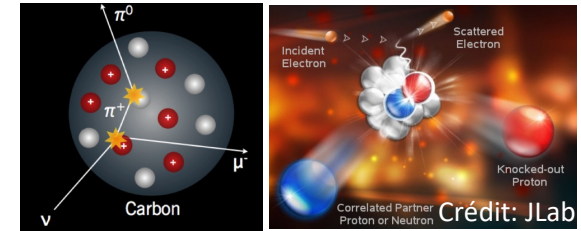
Dumitrescu, Hagen, Carlson, Papenbrock...

Dark matter



Mocz, Szász

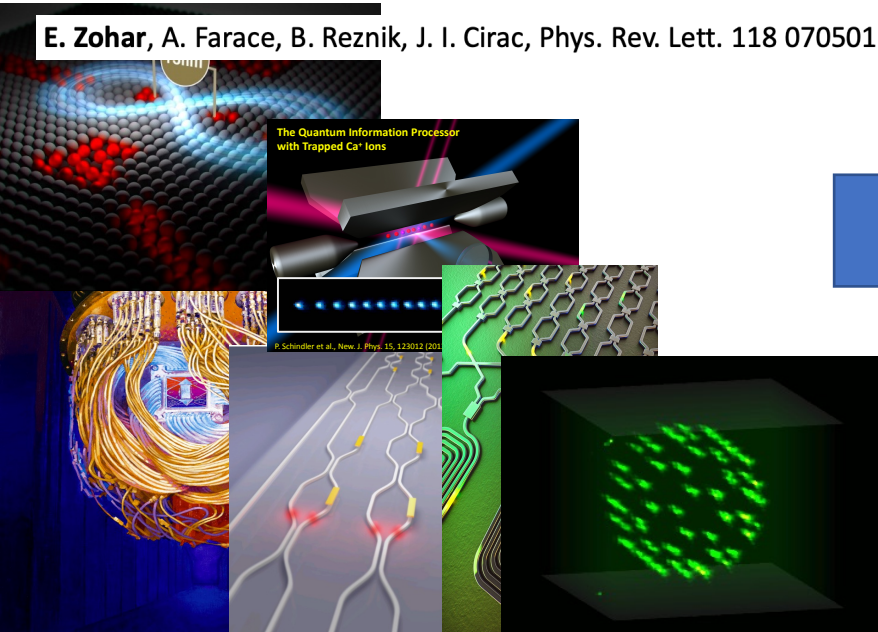
Dynamics: e, ν scattering



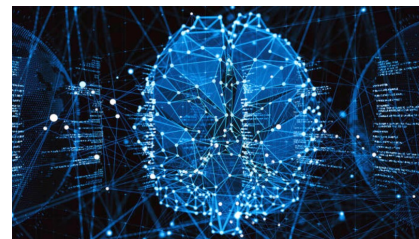
Roggero, Carlson, ...

- E. Zohar, J. I. Cirac, B. Reznik, Phys. Rev. Lett. **110**, 125304 (2013)
- E. Zohar, J. I. Cirac, B. Reznik, Phys. Rev. A **88** 023617 (2013)
- E. Zohar, J. I. Cirac, B. Reznik, Rep. Prog. Phys. **79**, 014401 (2016)
- D. González Cuadra, E. Zohar, J. I. Cirac, New J. Phys. **19** 063038 (2017)

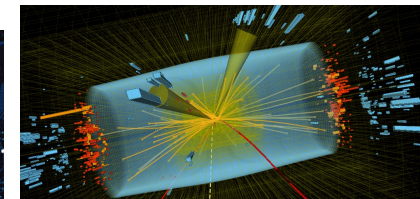
E. Zohar, A. Farace, B. Reznik, J. I. Cirac, Phys. Rev. Lett. 118 070501 (2017)

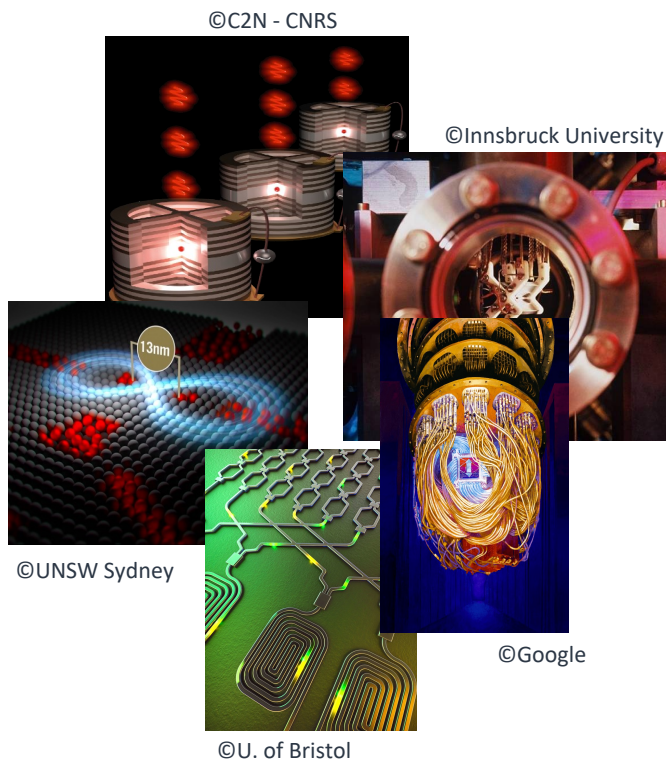


Applications to data mining (event classification)



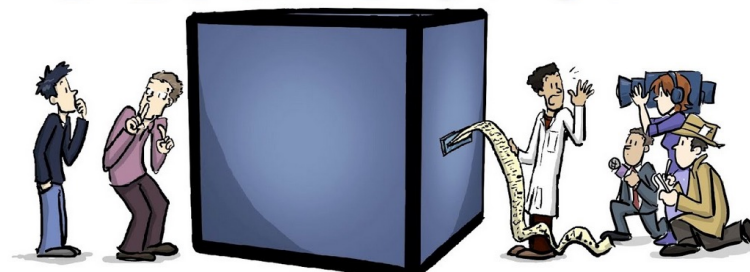
CMS-detector



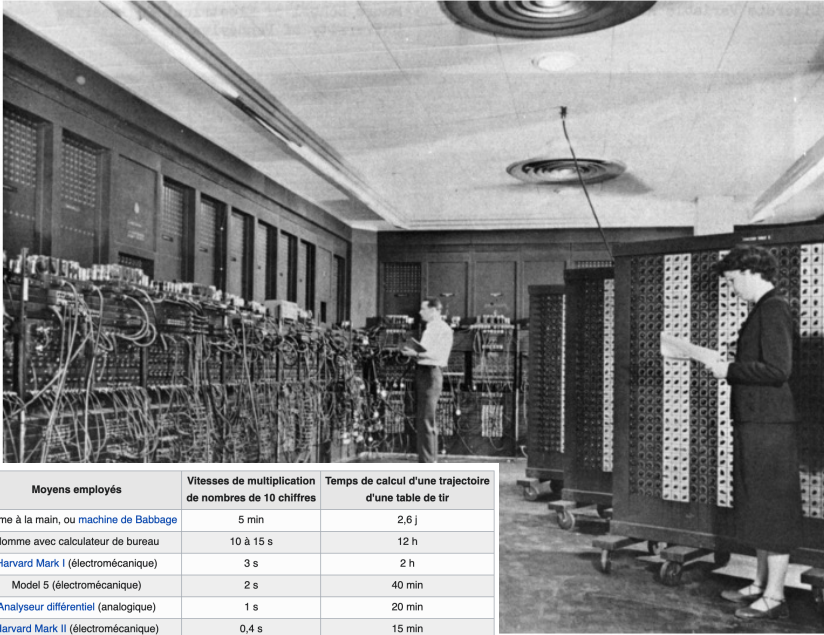


- ➔ Quantum computing is a high risk/high benefit interdisciplinary field
- ➔ It might lead to unprecedented boost in theory (or more generally in complex problems)
- ➔ It leads to natural link between public research and private companies (IBM, Google, ...)
- ➔ Emerging QC programs in France

A Quantum COMPUTER



Eniac ~1950



Moyens employés	Vitesses de multiplication de nombres de 10 chiffres	Temps de calcul d'une trajectoire d'une table de tir
Homme à la main, ou <i>machine de Babbage</i>	5 min	2,6 j
Homme avec calculateur de bureau	10 à 15 s	12 h
<i>Harvard Mark I</i> (électromécanique)	3 s	2 h
Model 5 (électromécanique)	2 s	40 min
<i>Analyseur différentiel</i> (analogique)	1 s	20 min
<i>Harvard Mark II</i> (électromécanique)	0,4 s	15 min
ENIAC (électronique)	0,001 s	3 s



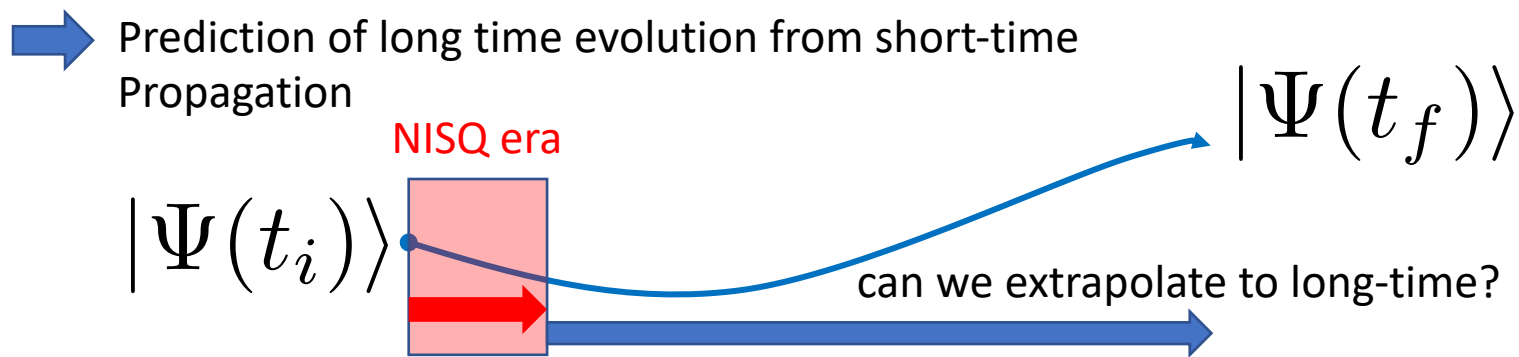
IBM ~2020



Thank you

Additional material

Predicting long time dynamics from short-time evolution



What is the physical content of short-time evolution ?

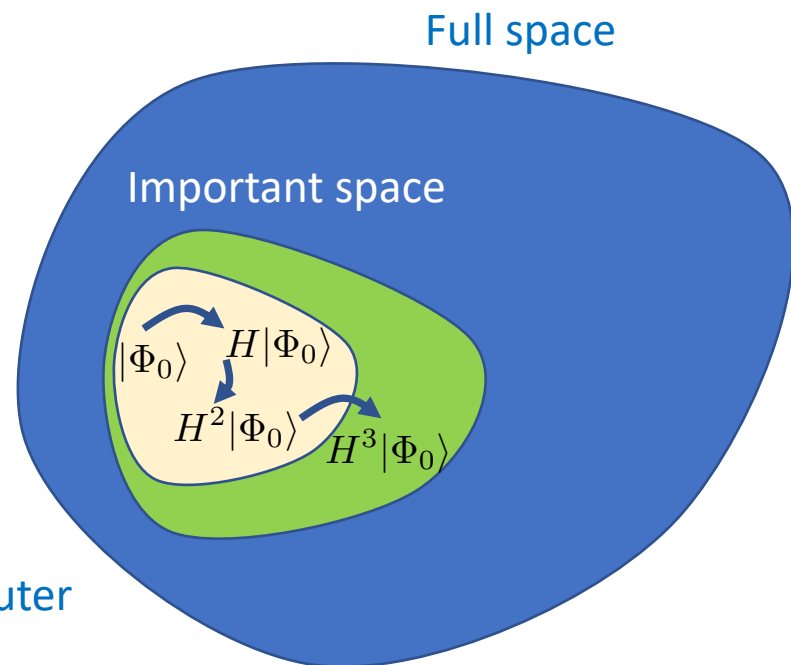
$$|\Phi(t)\rangle = \left(1 - itH + \frac{(-it)^2}{2!} H^2 + \dots \right) |\Phi(0)\rangle$$

➔ $H^K |\Phi(0)\rangle$

Are the so-called Krylov states

But they cannot be computed easily on a quantum computer

➔ We propose instead to compute $\langle H^K \rangle_0$



Hamiltonian moments calculation on a quantum computer

With minimal qubits number

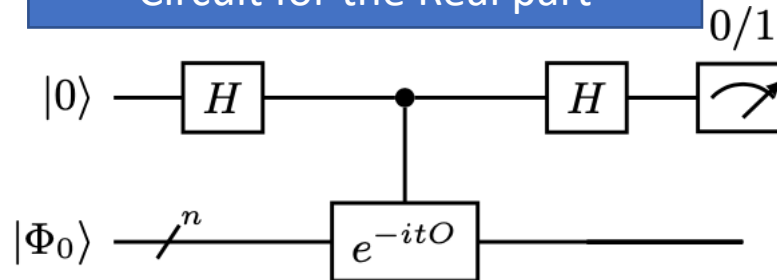
Generating function concept

$$F(t) = \langle \Phi_0 | e^{-itH} | \Phi_0 \rangle$$

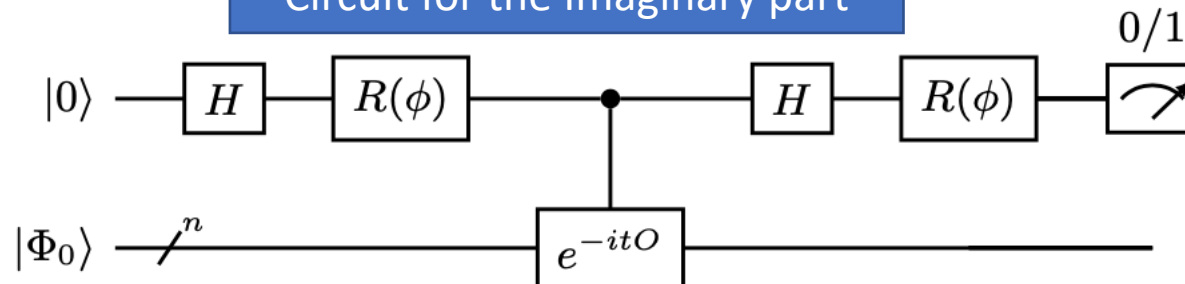
$$F(t) = 1 - it\langle H \rangle_0 + \frac{(-it)^2}{2} \langle H^2 \rangle_0 + \dots \Rightarrow \langle H^K \rangle_0 = i^K \left. \frac{d^K F(t)}{dt^K} \right|_{t=0}$$

Practical method to get $F(t)$

Circuit for the Real part



Circuit for the Imaginary part



Hamiltonian moments calculation on a quantum computer

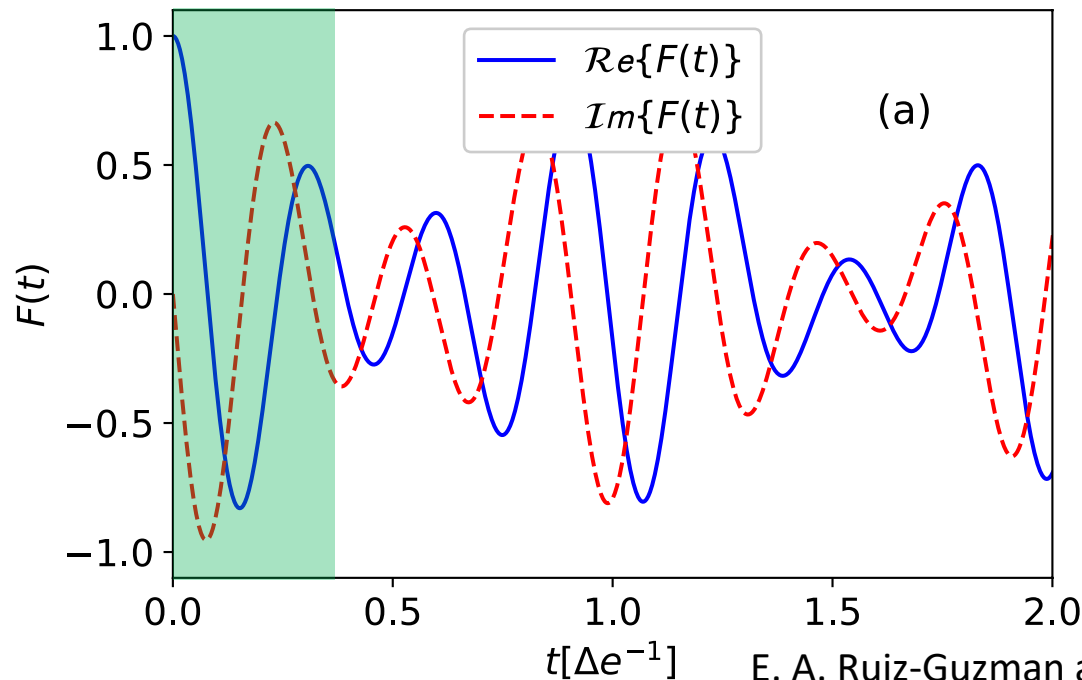
With minimal qubits number

Generating function concept

$$F(t) = \langle \Phi_0 | e^{-itH} | \Phi_0 \rangle$$

$$F(t) = 1 - it\langle H \rangle_0 + \frac{(-it)^2}{2} \langle H^2 \rangle_0 + \dots \Rightarrow \langle H^K \rangle_0 = i^K \left. \frac{d^K F(t)}{dt^K} \right|_{t=0}$$

Illustration for the cooper pair problem



finite difference
made on a classical
computer

➔ $\langle H^K \rangle_0$

Next use the moments for post-processing

$$\langle H^K \rangle_0$$

Ground state property
(imaginary time evolution)

Evolution: Krylov without
Krylov states

$$E(\tau) = \frac{\langle H e^{-\tau H} \rangle}{\langle e^{-\tau H} \rangle}$$

$$\frac{d}{d\tau} E(\tau) \simeq - \sum_{K=0}^{L-2} \frac{(-\tau)^K}{K!} \kappa_{K+2}$$

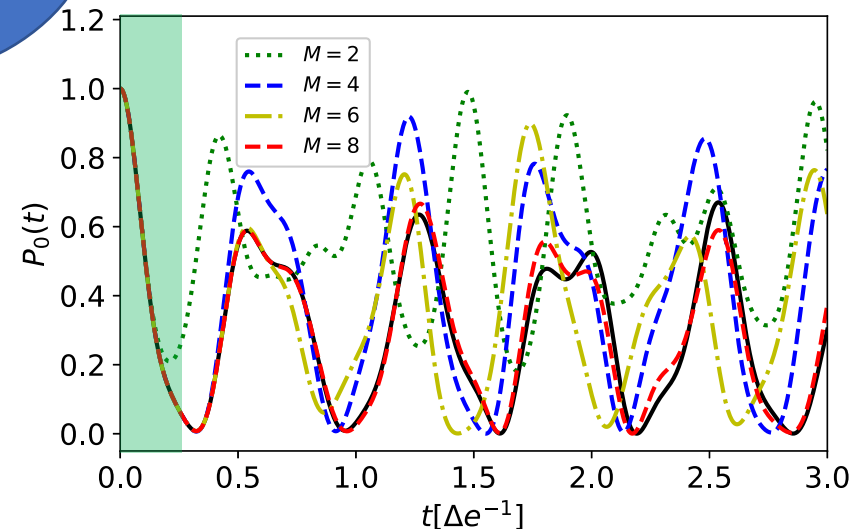
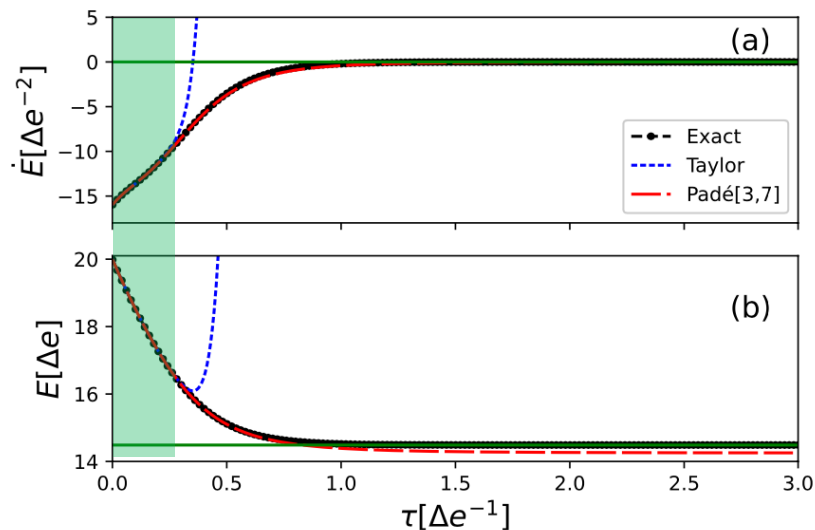
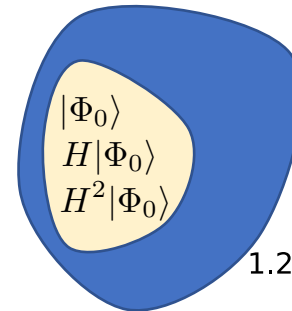


Illustration of the specific effort to make with qubits

Playing with binary numbers

Integer representation

$$j \xrightarrow{j = j_1 2^{n-1} + j_2 2^{n-2} \dots j_n 2^0} \{j_1, \dots, j_n\}$$

Binary representation

$$\frac{j}{2^n} \xrightarrow{\text{Binary fraction}} 0.j_1 \dots j_n = \frac{j_1}{2} + \frac{j_2}{2^2} + \dots + \frac{j_n}{2^n}$$

Binary fraction

Note that this is just the equivalent to our base 10 writing

$$0.j_1 \dots j_n = \frac{j_1}{10} + \frac{j_2}{10^2} + \dots + \frac{j_n}{10^n}$$

Example $0.1011 \rightarrow 1 \times 0.5 + 0 \times 0.25 + 1 \times 0.125 + 1 \times 0.0625$
 $\rightarrow 0.6875$

Approximation of 0.6875 can be obtained by removing some terms

$$0.101 \rightarrow 0.625$$

$$0.1 \rightarrow 0.5$$

2^{-1}	=	0.5
2^{-2}	=	0.25
2^{-3}	=	0.125
2^{-4}	=	0.0625
2^{-5}	=	0.03125
2^{-6}	=	0.015625
2^{-7}	=	0.0078125
...		

Illustration of the specific effort to make with qubits

A seminal example: the Quantum Fourier Transform (QFT)

Reminder: classical FT



Quantum FT



Qubit FT

Assume a set of values $\{x_0, \dots, x_{N-1}\}$, we can define a set of new numbers $\{f_0, \dots, f_{N-1}\}$ such that:

$$f_k = \frac{1}{\sqrt{N}} \sum_{j=0}^{N-1} x_j e^{2i\pi k(j/N)}$$

Then the inverse Fourier transform is given by¹:

$$x_j = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} f_k e^{-2i\pi j(k/N)}$$



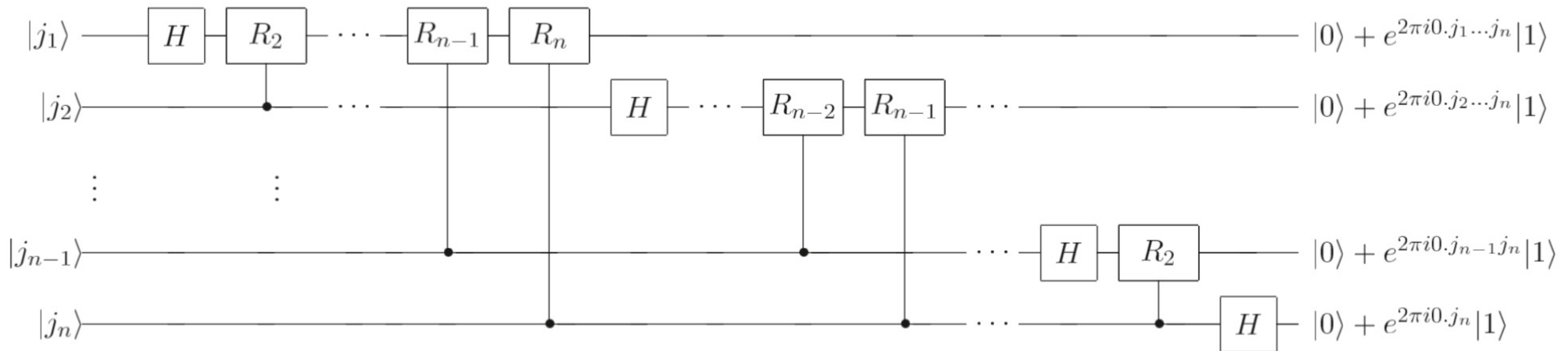
$$|j\rangle \rightarrow |\tilde{j}\rangle \equiv \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} e^{2i\pi \alpha(k/N)} |k\rangle$$



$$\begin{aligned} |j\rangle = |j_{n-1} \dots j_0\rangle &\rightarrow |\tilde{j}\rangle = \frac{1}{2^{n/2}} \sum_{k=0}^{2^n-1} e^{2i\pi jk/2^n} |k\rangle \\ &= \frac{1}{2^{n/2}} \bigotimes_{l=1}^n \left[|0\rangle + e^{2i\pi j 2^{-l}} |1\rangle \right] \\ &= \frac{1}{2^{n/2}} \left(|0\rangle + e^{2i\pi i_0 j_0} |1\rangle \right) \left(|0\rangle + e^{2i\pi i_0 j_1 j_0} |1\rangle \right) \dots \left(|0\rangle + e^{2i\pi i_0 j_{n-1} \dots j_1 j_0} |1\rangle \right) \end{aligned}$$

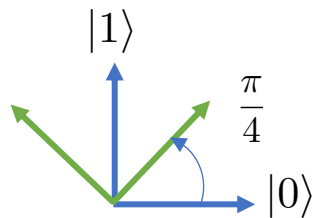
$$|j\rangle \rightarrow \frac{1}{2^n} \bigotimes_{l=1}^n \left[|0\rangle + e^{2\pi i j 2^{-l}} |1\rangle \right]$$

Quantum circuits



H-gate (unary gate)

Controlled -U gate (binary gate)



$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

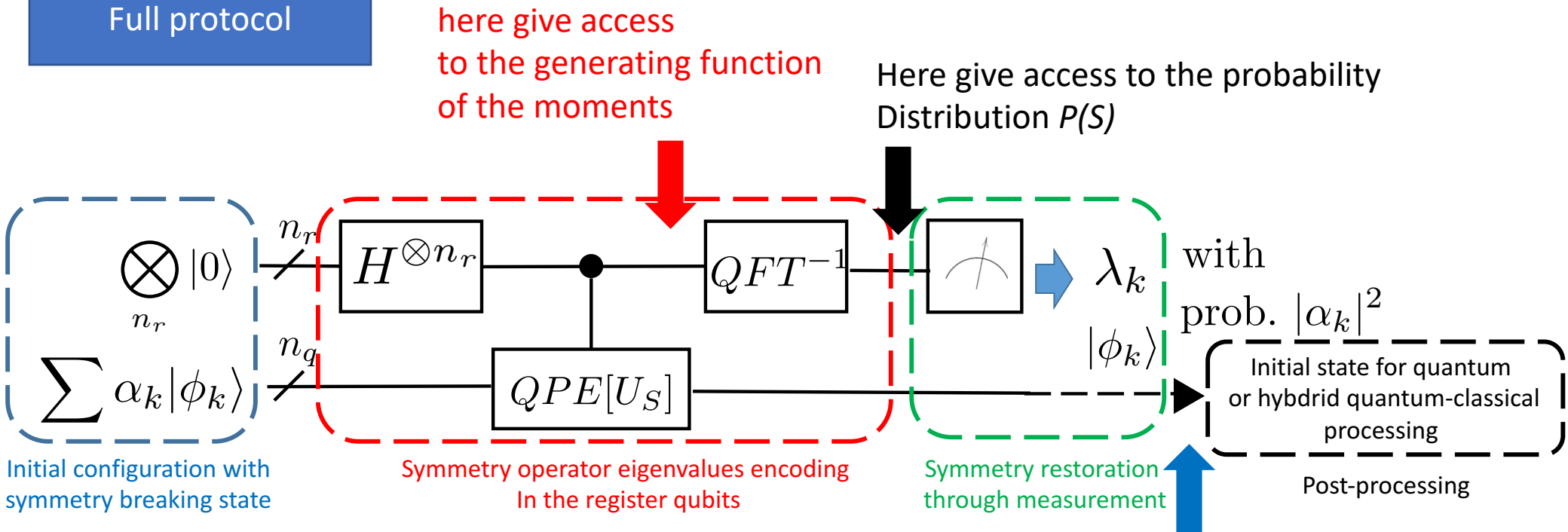
$$\text{Controlled -U} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & U_{11} & U_{12} \\ 0 & 0 & U_{21} & U_{22} \end{bmatrix} \begin{matrix} |00\rangle \\ |01\rangle \\ |10\rangle \\ |11\rangle \end{matrix}$$

$$R_k = \begin{bmatrix} 1 & 0 \\ 0 & e^{2\pi i/2^k} \end{bmatrix}$$

The quantum-Phase estimate (QPE) algorithm

For eigenvalue problems

Full protocol



Important remarks

In practice, for each measurement we obtain a binary string

$$\{0110 \dots\}_{n_r} \rightarrow \theta_k \rightarrow |2^{n_r} \theta_k\rangle \otimes |\phi_k\rangle \rightarrow |\phi_k\rangle \text{ for post-processing}$$

After measurement
The state is projected
on eigenstates or a set of
Eigenstates if degenerated

Use the QPE approach for operators with known eigenvalues to obtain entangled states

Hypothesis:

- ▶ Assume a hermitian operator S acting on nq qubits
- ▶ Assume that S has discrete eigenvalues $\{\lambda_0 \leq \dots \leq \lambda_M\}$ with $\lambda_k = am_k$
 $a = \text{cst}$
- ▶ Define the operator

$$U_S = \exp \left\{ 2\pi i \left[\frac{S - \gamma_0}{a2^{n_0}} \right] \right\}$$

- ▶ Eigenvalues of U_S are given by $e^{2\pi i \theta_k}$ with $\theta_k = (m_k - m_0)/2^{n_0}$

If $(m_k - m_0) < 2^{n_0}$ \Rightarrow $\theta_k < 1$
and θ_k is exactly written as a binary fraction

It is then optimal for the QPE use.
An optimal choice for the number of register qubits is $n_r = n_0$

and $n_r - 1 \leq \ln(m_k - m_0) / \ln 2 < n_r$.

Examples

- parity
- Part. number
- $J_z = \hbar m$
- $J^2 = \hbar^2 j(j+1)$

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- Part. number
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Illustration of the Jordan-Wigner use

Quantum Computing notations

$$\{\sigma_x, \sigma_y, \sigma_z\} \rightarrow \{X, Y, Z\}$$

Convention:

$$Q_j^+ = \frac{1}{2} (X_j - iY_j) \rightarrow Q_j^+ |0_j\rangle = |1_j\rangle$$

Example: particle number

$$\hat{N} = \sum_j a_j^\dagger a_j \xrightarrow[Z_j^2]{a_k^\dagger \rightarrow \prod_{k<j} (-Z_k) Q_j^+} N = \sum_j Q_j^+ Q_j = \frac{1}{2} \sum_j (I_j - Z_j) \rightarrow \begin{aligned} Q_j^+ Q_j |0_j\rangle &= 0 \\ Q_j^+ Q_j |1_j\rangle &= |1_j\rangle \end{aligned}$$

Particle number projection with the QPE

For n_q qubits, $n = 0, \dots, n_q$

$$U_N = \exp \{2\pi i N / 2^{n_r}\}$$

With the Condition $n_q < \sim 2^{n_r}$

Reminder

$$X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = |1\rangle\langle 0| + |0\rangle\langle 1|,$$

$$Y = i \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix} = i|1\rangle\langle 0| - i|0\rangle\langle 1|$$

$$Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} = |0\rangle\langle 0| - |1\rangle\langle 1|$$

N counts the number n of 1 in qubits

$n_q = 1$	\rightarrow	$n_r = 1$
$2 \leq n_q < 3$	\rightarrow	2
$4 \leq n_q < 7$	\rightarrow	3
$8 \leq n_q < 15$	\rightarrow	4
$16 \leq n_q < 31$	\rightarrow	5
$32 \leq n_q < 63$	\rightarrow	6
$64 \leq n_q < 127$	\rightarrow	7
$128 \leq n_q < 255$	\rightarrow	8
$256 \leq n_q < 511$	\rightarrow	9
$512 \leq n_q < 1023$	\rightarrow	10

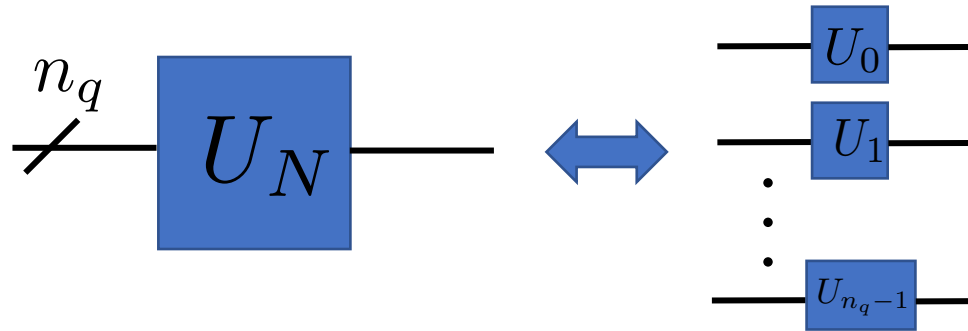
QPE applied to particle number projection

Practical details

$$U_N = \prod_j U_j$$

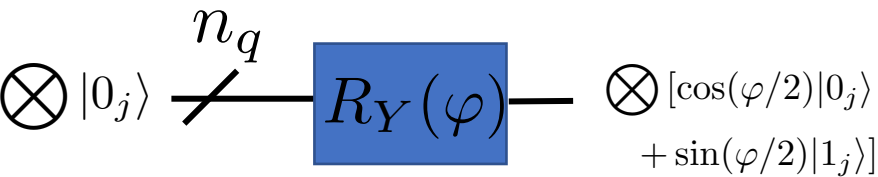
$$U_i = |0_i\rangle\langle 0_i| + \exp(i\pi/2^{n_0-1})|1_i\rangle\langle 1_i|$$

$$U_i = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi/2^{n_0-1}} \end{bmatrix}$$



Example: Qubit counting statistics

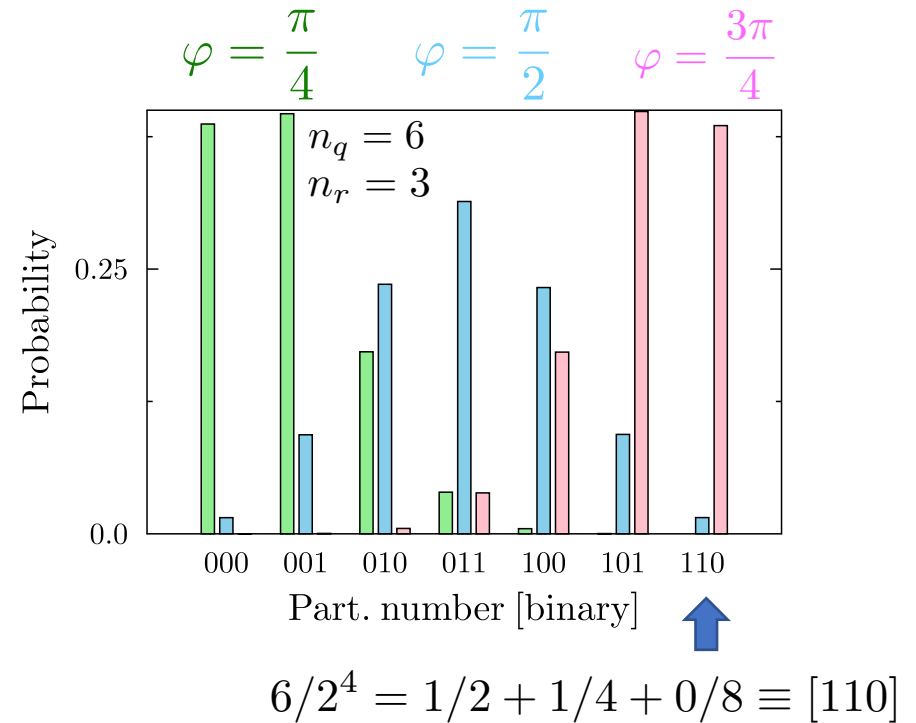
Initial state



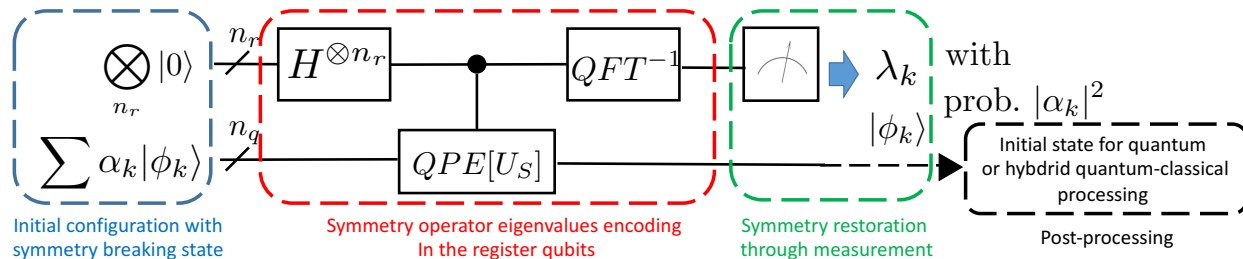
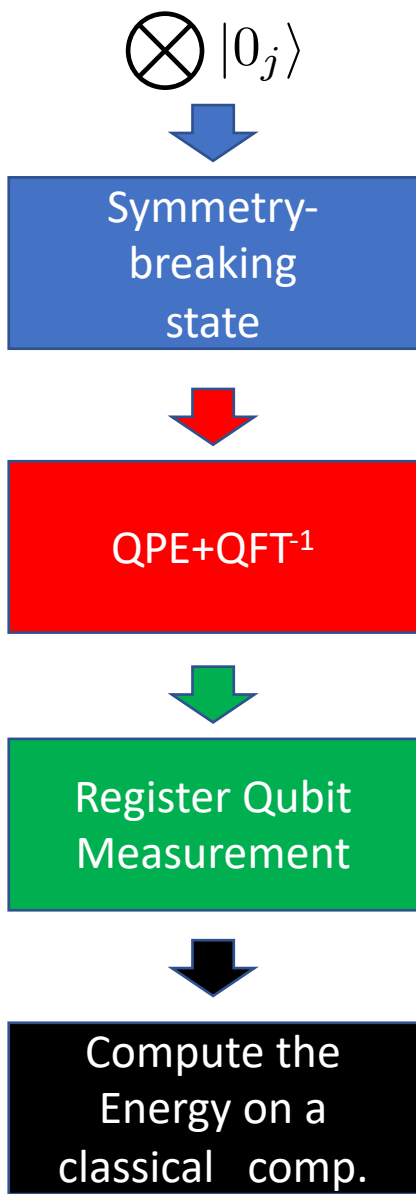
$$P(A) = C_{n_q}^A p^A (1-p)^{n_q-A}$$

$$p = \sin^2(\varphi/2)$$

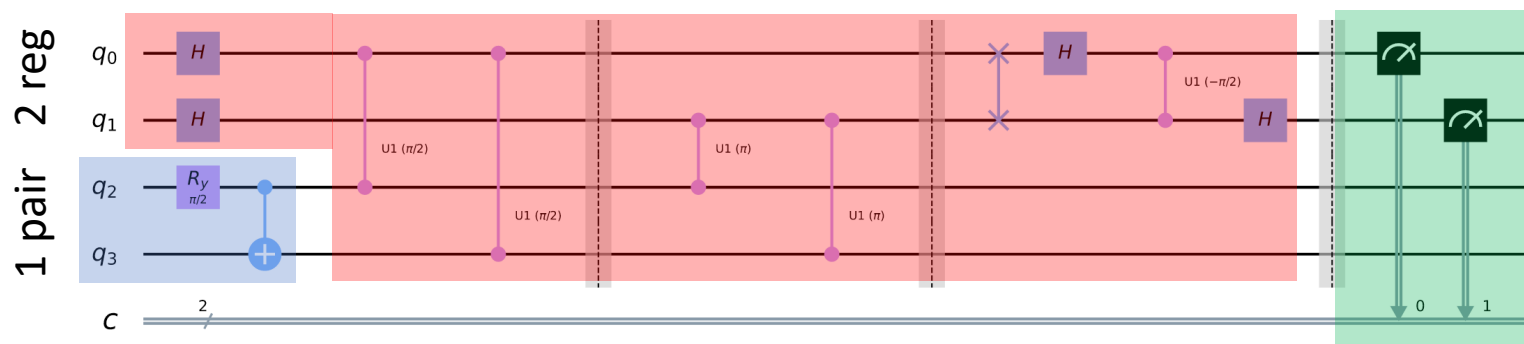
Calculation made with the IBM Qiskit python package



Applying the strategy to the pairing problem



Qiskit circuit for a single pair



Applying the strategy to the pairing problem

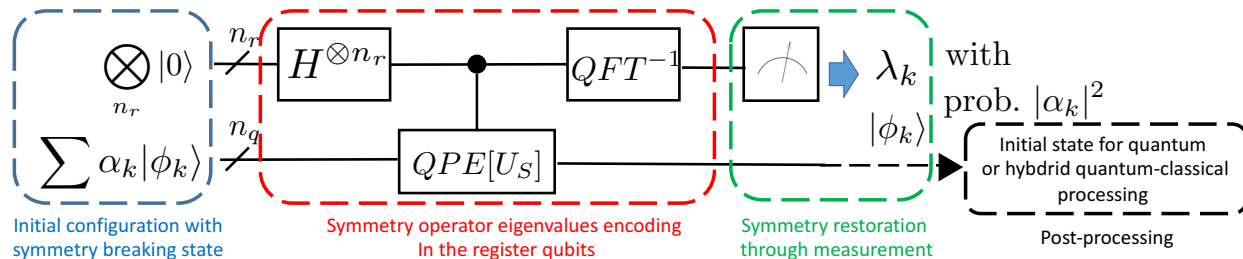
$$\bigotimes_{j=1}^n |0_j\rangle$$

Symmetry-breaking state

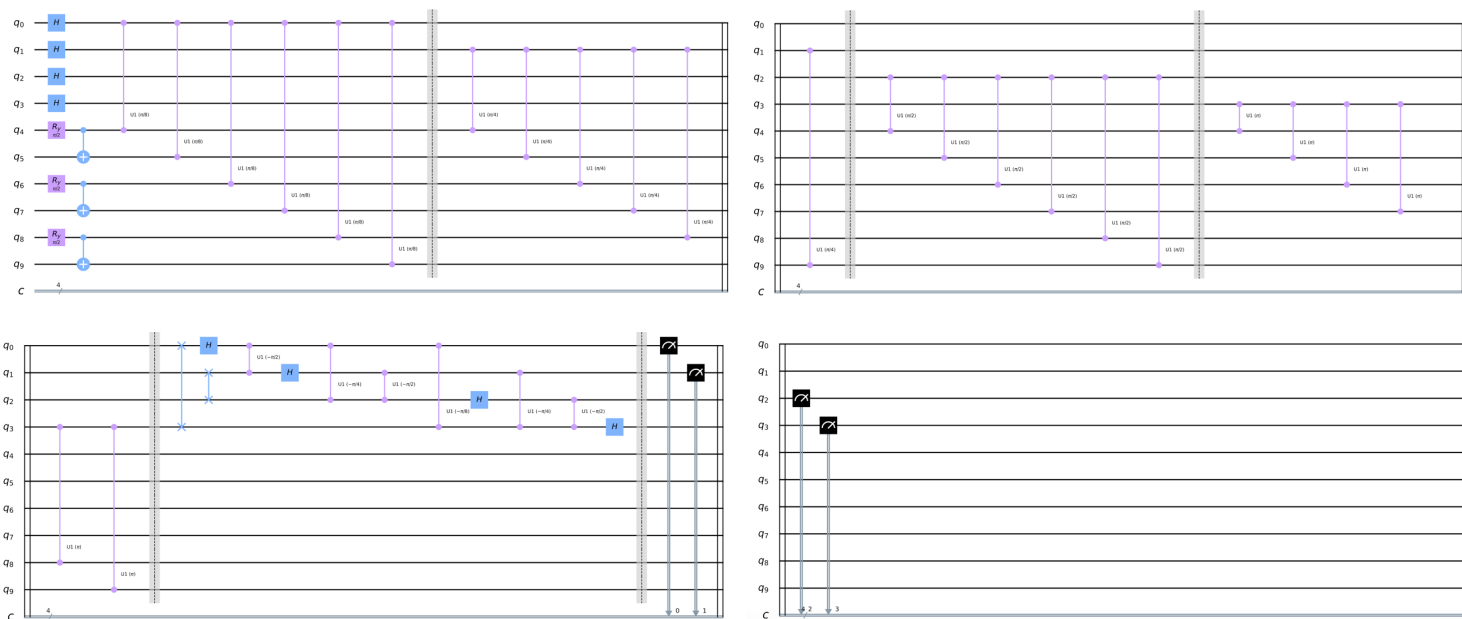
QPE+QFT⁻¹

Register Qubit Measurement

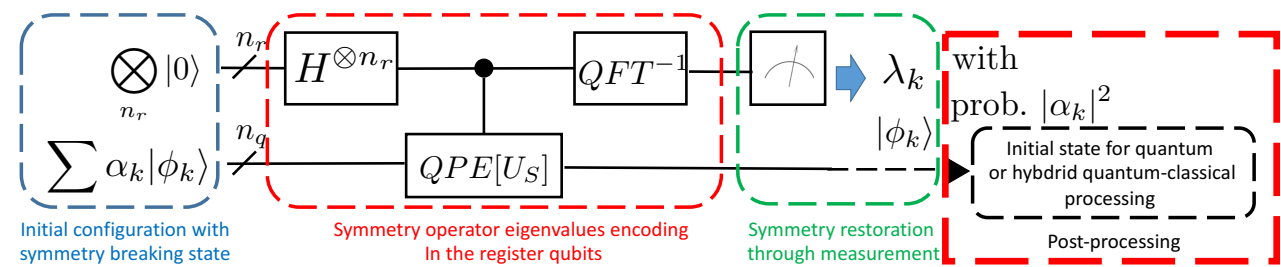
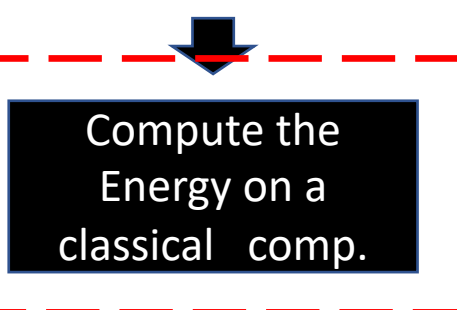
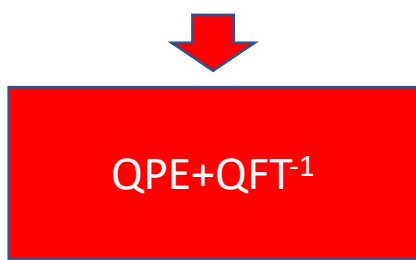
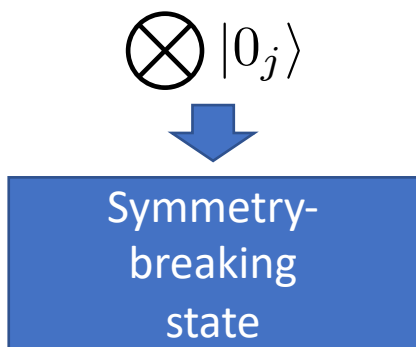
Compute the Energy on a classical comp.



3 pairs, 4 register



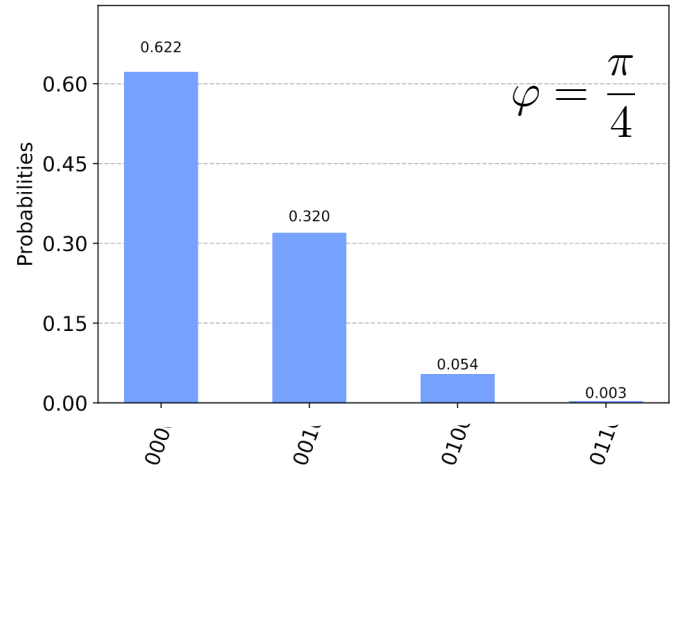
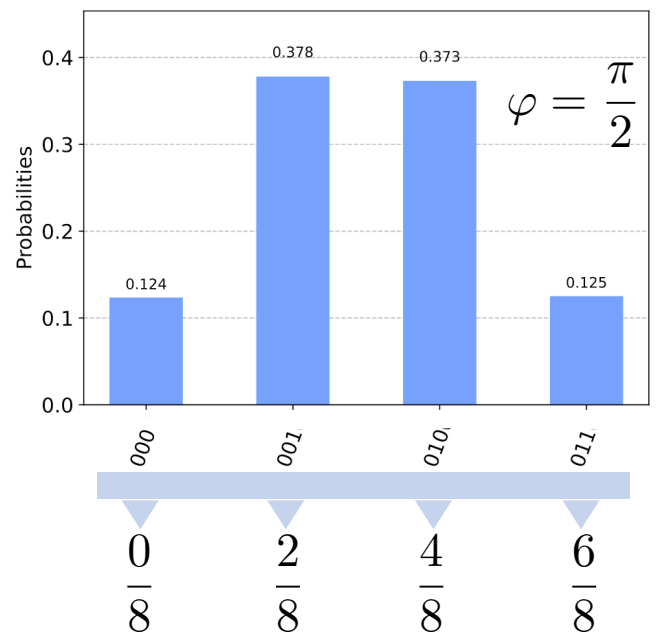
Applying the strategy to the pairing problem



Qubit counting statistics

Initial state $|\Psi\rangle = \prod_n \left[\cos\left(\frac{\varphi}{2}\right) I_n \otimes I_{n+1} + \sin\left(\frac{\varphi}{2}\right) Q_n^+ Q_{n+1}^+ \right] |-\rangle$

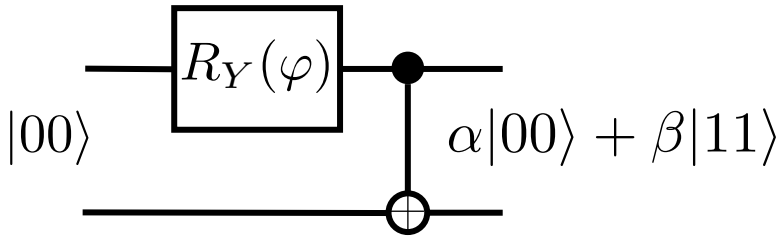
Example: 3 pairs, 3 registers



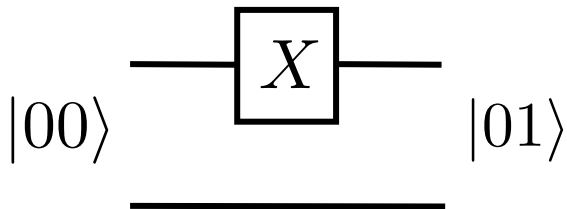
Exploring the possibilities of QC Using the Broken pair approximation

BCS/HFB state

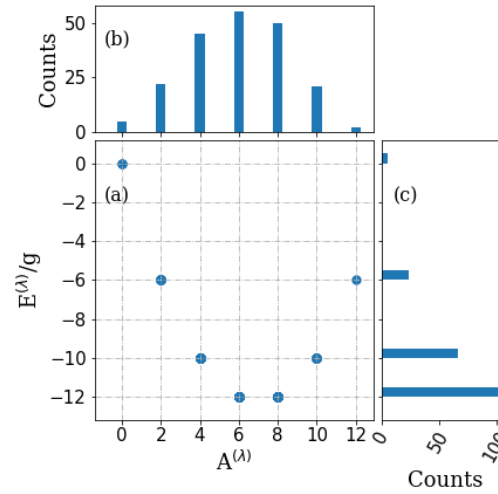
All pairs= generalized Bell state



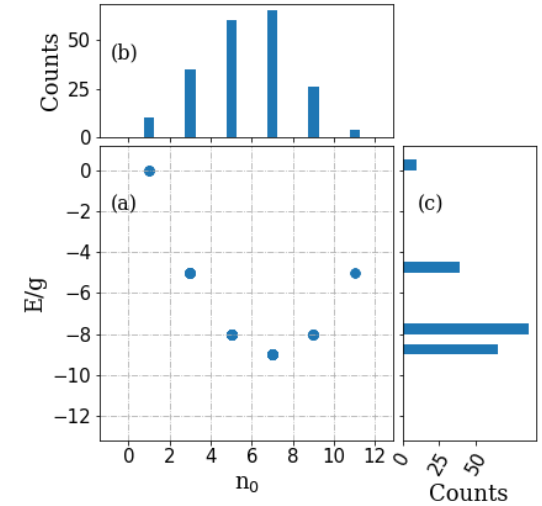
BCS/HFB state
+ some broken
pairs



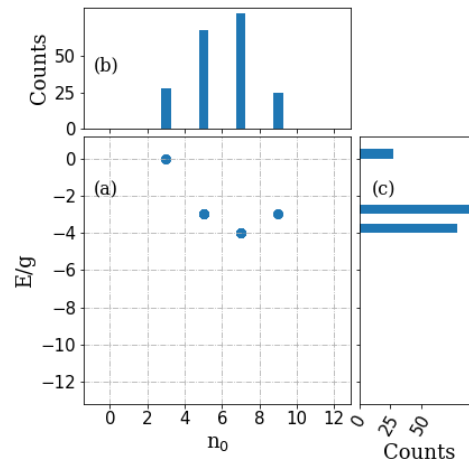
No broken pairs



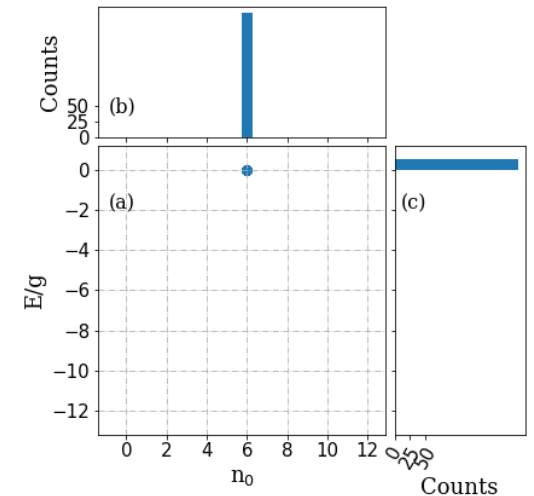
1 broken pairs



3 broken pairs

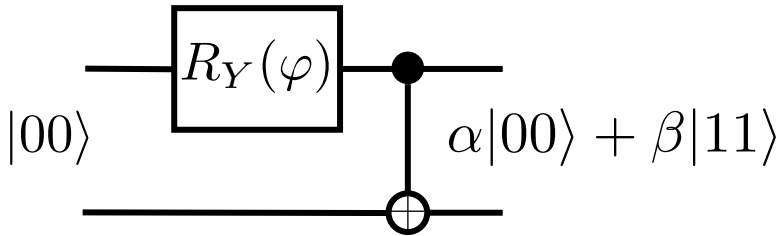


6 broken pairs

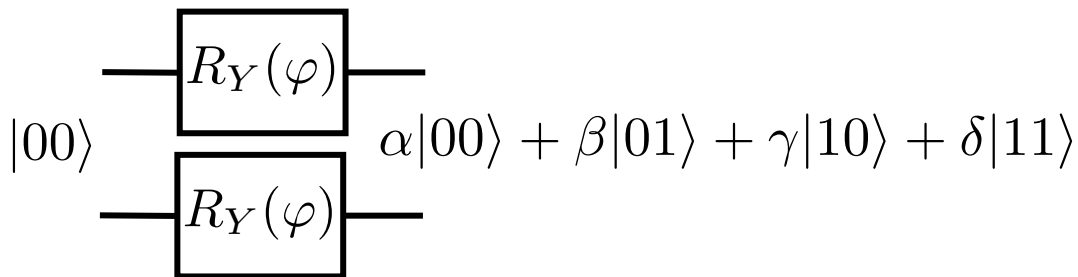


BCS/HFB state

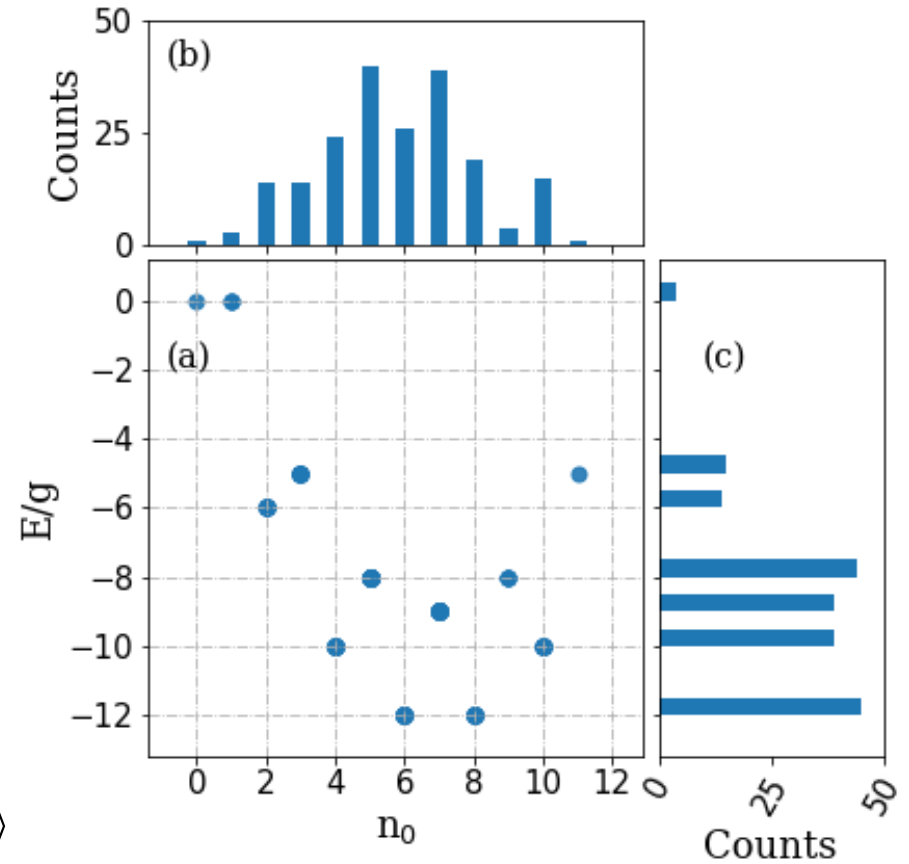
All pairs= generalized Bell state



Alternative circuits



Can give both odd and even simultaneously

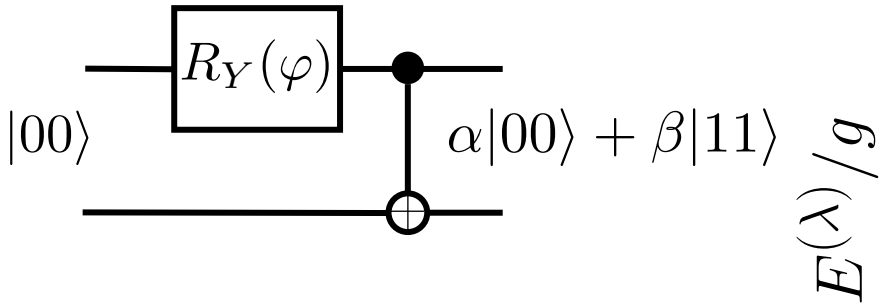


Exploring the possibilities of QC

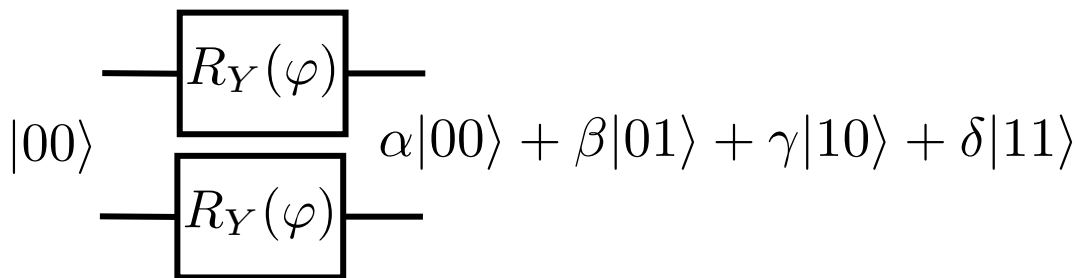
Some additional remarks

BCS/HFB state

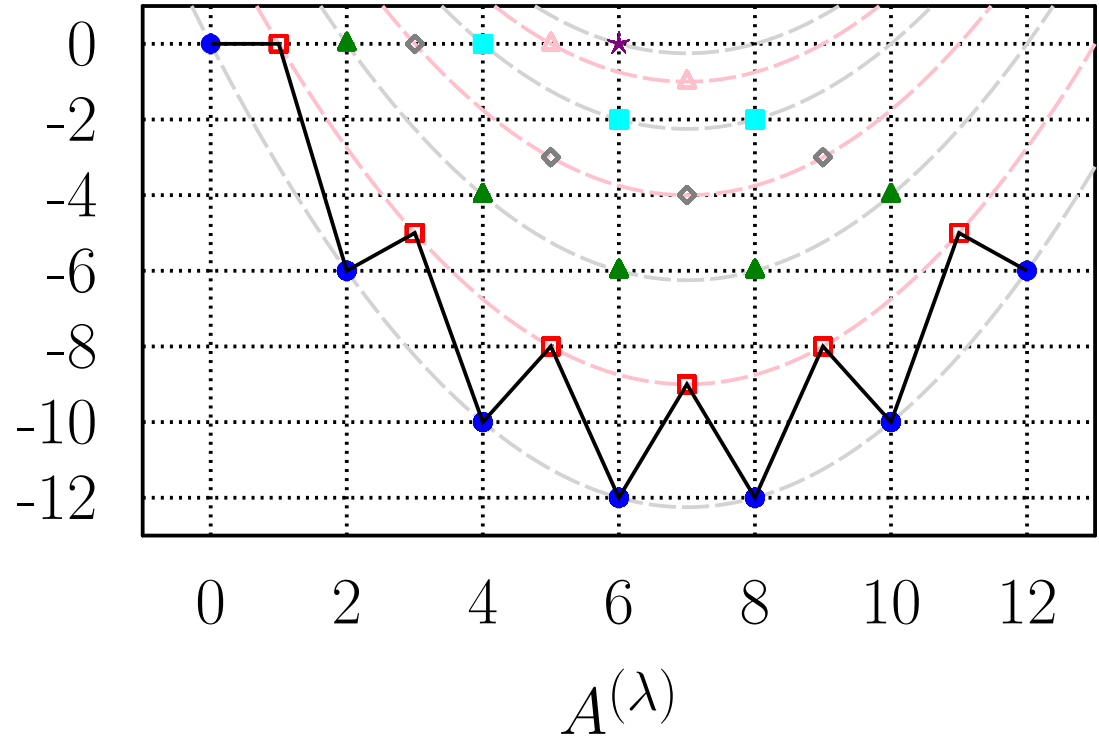
All pairs= generalized Bell state



Alternative circuits



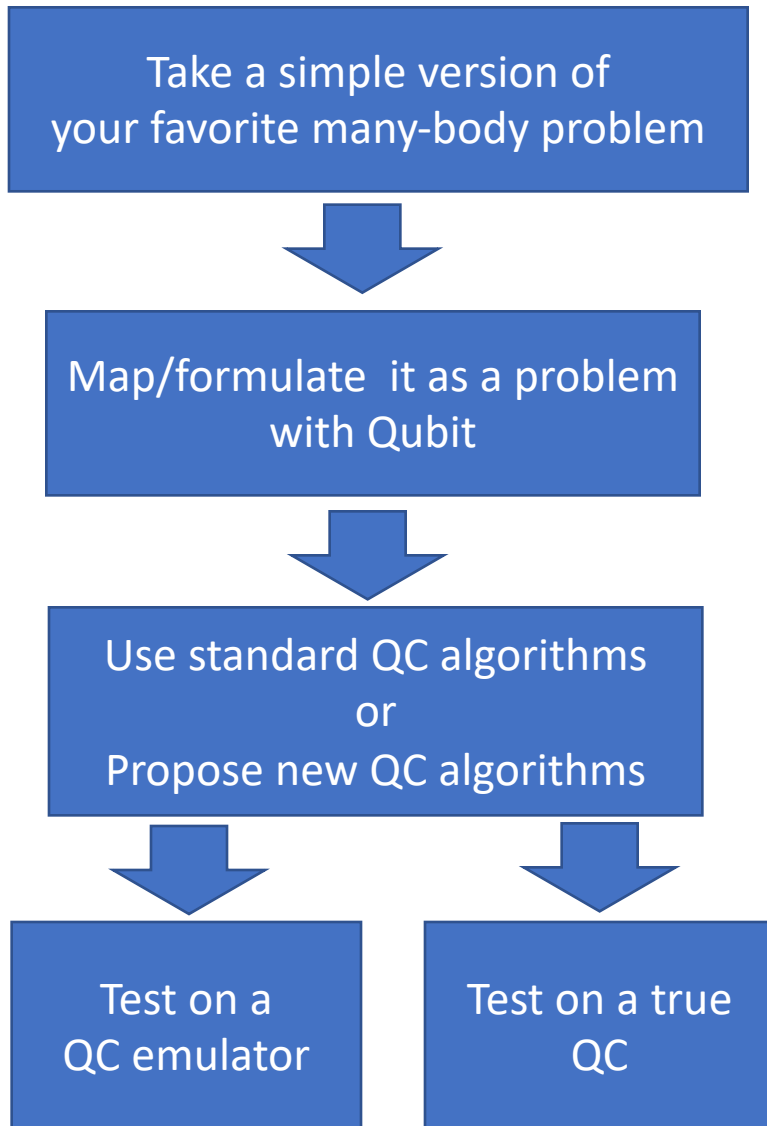
Altogether



Exact solution (lines)

$$E/g = -\frac{1}{4}(A - \nu)(2n_q - A - \nu - 2).$$

Strategy



σ Models on Quantum Computers

Andrei Alexandru,^{1,2,*} Paulo F. Bedaque,^{2,†} Henry Lamm^{2,‡} and Scott Lawrence^{2,§}

(NuQS Collaboration)

Start with the discretized σ model

$$\mathcal{H} = \sum_r \left(\frac{g^2}{2} \boldsymbol{\pi}(r)^2 + \frac{1}{2g^2 \Delta x^2} [\mathbf{n}(r+1) - \mathbf{n}(r)]^2 \right)$$

Map it to a Spin algebra (fuzzy sphere)

$$-\frac{g^2}{2} \nabla^2 \psi \rightarrow H^0 \Psi = \kappa \frac{g^2}{2} \sum_{k=1}^3 [\mathbb{J}_k, [\mathbb{J}_k, \Psi]],$$

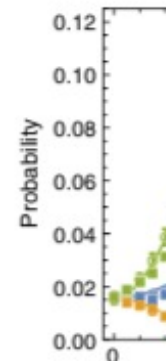
\mathbb{J}_k are generators of the SU(2) algebra

“only” $j=1/2$ was considered

$$j_1 = \mathbb{1} \otimes \sigma_2 / \sqrt{3}, \quad j_2 = \sigma_2 \otimes \sigma_3 / \sqrt{3}, \quad j_3 = \sigma_2 \otimes \sigma_1 / \sqrt{3},$$

This gives the link with Pauli matrices

Use the Suzuki + specific QC



GAUGE THEORIES FOR QUANTUM COMPUTING

Very limited studies exist, e.g.:

- 1+1 D QED (Schwinger model) on a few-qubit trapped-ion quantum computer [E. A. Martinez et al., Nature 534 (2016) 516, arXiv:1605.04570]
- Quantum-Classical calculation of Schwinger Model [N. Kico et al., Phys. Rev. A 98 (2018) 032331, arXiv:1803.03326]
- U(1) lattice gauge theory without matter in 2 & 3 spatial dimensions [D. Kaplan, J. Stryker, arXiv:1806.08797]
- Zeta-regularized vacuum expectation values [T. Hartung, K. Jansen, arXiv:1808.06784]
- O(3) nonlinear sigma model in 1+1 dimensions [F. Bruckmann, K. Jansen, S. Kuhn, arXiv:1812.00944]

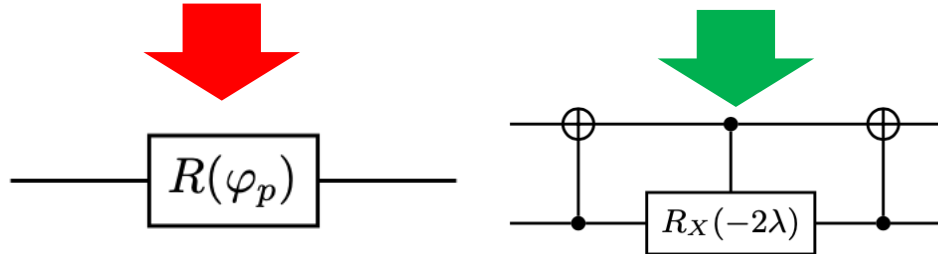
Extending the studies to 2+1 dimensions is extremely hard and has not been established yet on quantum devices

(from M. constantinou, Santa Fe)

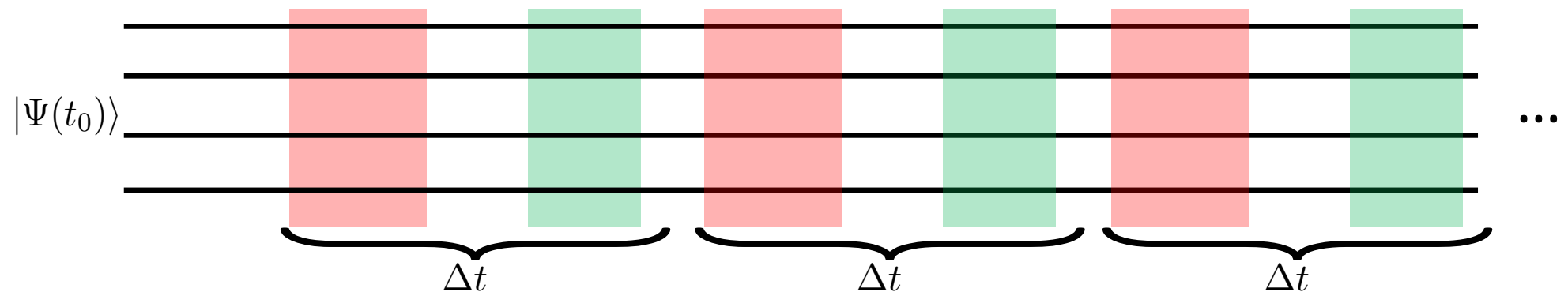
Pairing Hamiltonian

$$H_P = \sum_{i>0} \varepsilon_i (a_i^\dagger a_i + a_{\bar{i}}^\dagger a_{\bar{i}}) - g \sum_{i,j>0} a_i^\dagger a_{\bar{i}}^\dagger a_{\bar{j}} a_j$$

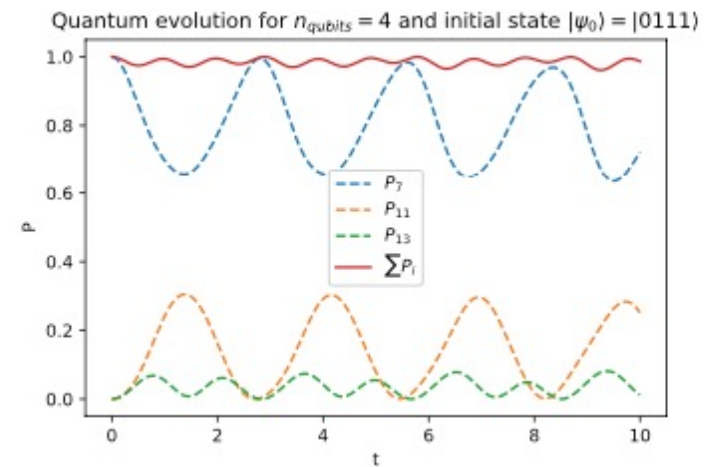
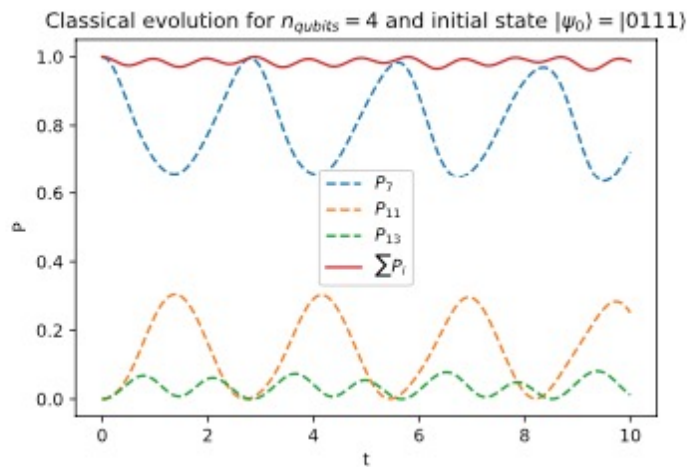
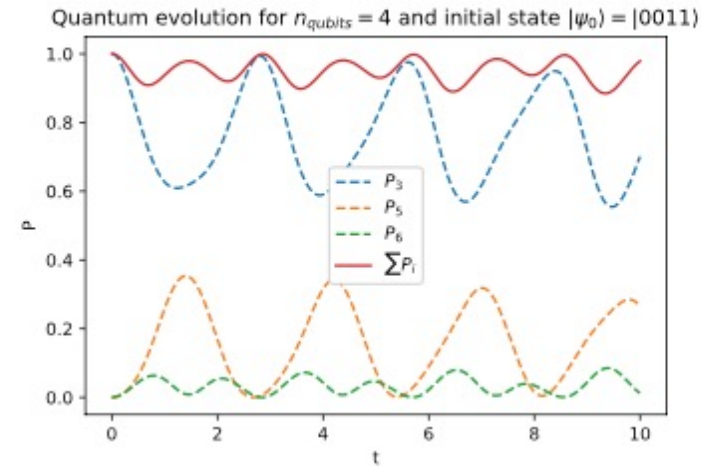
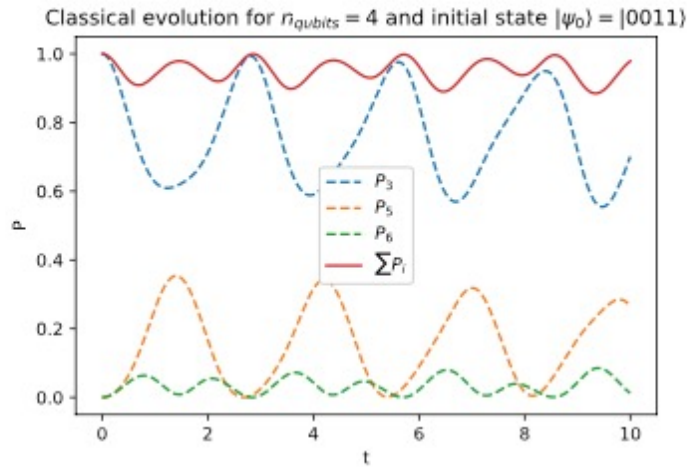
$$\prod_p \begin{pmatrix} 1 & 0 \\ 0 & \exp(-2i\tilde{\varepsilon}_p \Delta t) \end{pmatrix} \prod_{p>q} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & \cos(\lambda_{pq}) & i \sin(\lambda_{pq}) & 0 \\ 0 & i \sin(\lambda_{pq}) & \cos(\lambda_{pq}) & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$



The problem is that we can nowadays perform only few operations and with a limited *fidelity*


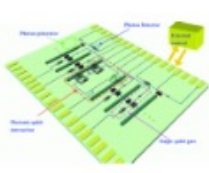





Results for 4 qubits



(From E. A. Ruiz Guzman, in PhD@IJCLab)

Building quantum computers: companies

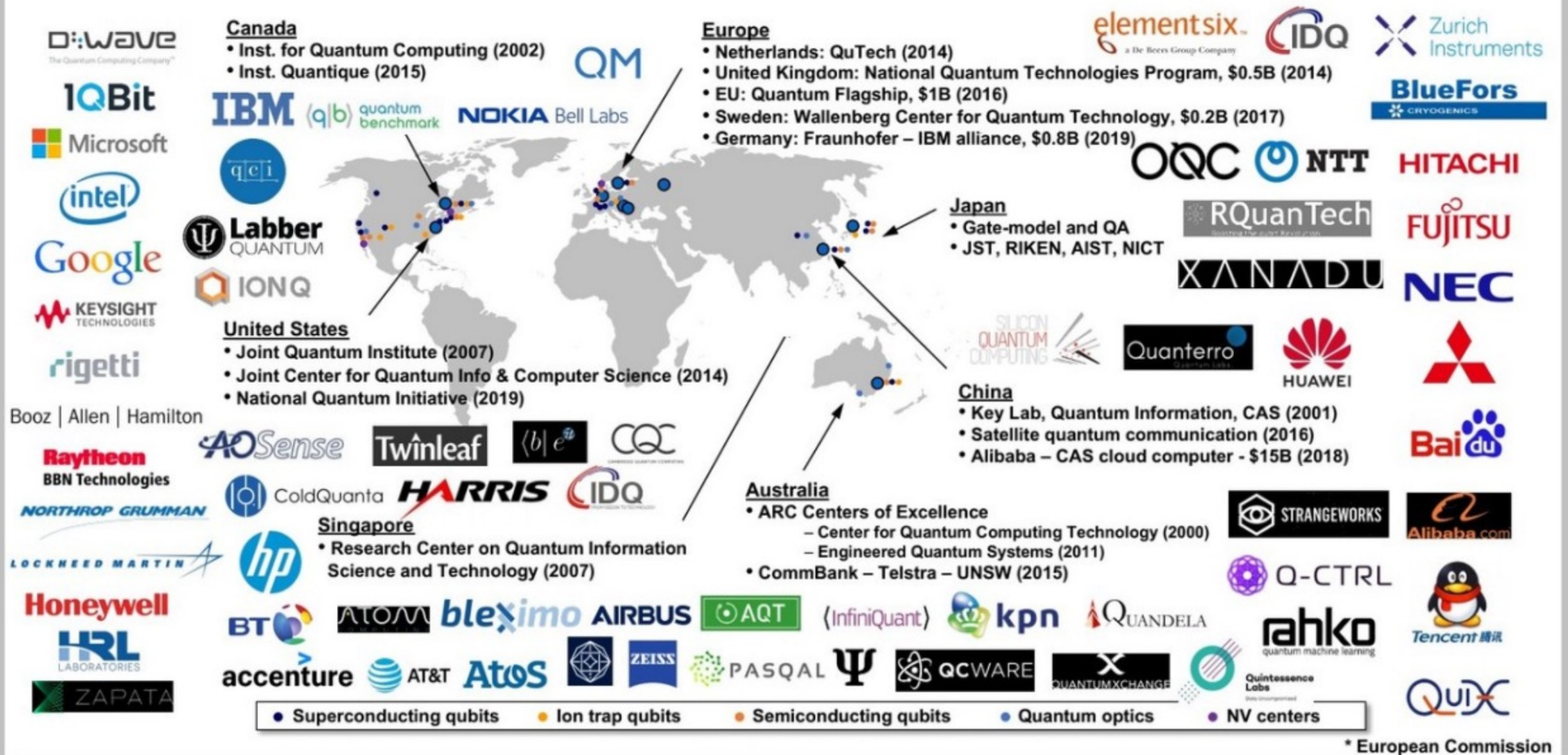
	atomes	électrons					photons	
	 <p>ions piégés</p>	 <p>atomes froids</p>	 <p>recuit quantique</p>	 <p>supra-conducteurs</p>	 <p>silicium</p>	 <p>centres NV (diamant)</p>	 <p>qubits topologiques</p>	 <p>photons</p>
entreprises et startups								
laboratoires (*)								

(cc) Olivier Ezratty, 5 septembre 2020

From article [O. Ezratty](#)

Quantum Investment Worldwide

(not exhaustive)



* European Commission